

Tiling Parity Results and the Holey Square Solution

Bridget Eileen Tenner

Department of Mathematics
Massachusetts Institute of Technology
Cambridge, MA 02139, USA
`bridget@math.mit.edu`

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Abstract

We prove combinatorially that the parity of the number of domino tilings of a region is equal to the parity of the number of domino tilings of a particular subregion. Using this result we can resolve the holey square conjecture. We furthermore give combinatorial proofs of several other tiling parity results, including that the number of domino tilings of a particular family of rectangles is always odd.

1 Introduction

The number of domino tilings of the $2n \times 2n$ square with a centered hole of size $2m \times 2m$, a figure known as the *holey square* and denoted $\mathcal{H}(m, n)$, was conjectured by Edward Early to have the form $2^{n-m}(2k_{m;n} + 1)^2$. Several people have worked on this problem and obtained partial results, but the general conjecture has remained open until now.

In the mid to late 1990s, Jim Propp suggested a variety of tiling enumeration questions, together with the current status of their solutions [4]. Early, a member of Propp's Tilings Research Group at MIT, posed the holey square conjecture in 1997, and Lior Pachter subsequently showed that the number of domino tilings of $\mathcal{H}(m, n)$ is 2^{n-m} times a perfect square [3]. This result follows from a factorization theorem of Mihai Ciucu [1], who was also aware of the implication but did not work specifically on Early's conjecture. In [5], Roberto Tauraso proves the result for $m = n - 2$ and discusses the sequence of odd factors in this case.

Ciucu's factorization theorem, concerning the number of perfect matchings $M(G)$ of a bipartite graph G with a particular symmetry property, furthermore describes the remaining squared factor as the number of domino tilings of a particular subregion of $\mathcal{H}(m, n)$. The result allows weighted edges, but for our purposes we assume that each edge has weight 1, and simplify the statement of the theorem accordingly. Suppose that a bipartite graph G is invariant under reflection across the straight line ℓ , and the set of

vertices of G lying on ℓ is a cut set. The number of vertices in $\ell \cap G$ must be even, and are labeled consecutively $a_1, b_1, \dots, a_{w(G)}, b_{w(G)}$. Color the bipartition classes of G black and white, and remove the edges of all white a_i and black b_i vertices on one side of ℓ , and the edges of all black a_i and white b_i vertices on the other side of ℓ . Let G^+ be the subgraph on one side of ℓ , and G^- the subgraph on the other side of ℓ .

Theorem 1 (Ciucu). *Let G be a bipartite symmetric graph separated by its axis of symmetry in the manner described above. Then*

$$M(G) = 2^{w(G)} M(G^+) M(G^-).$$

Writing $\#R$ for the number of domino tilings of the region R , we immediately see from Theorem 1 (or rather, its dual) that $\#\mathcal{H}(m, n) = 2^{n-m} (\#H(m, n))^2$ for a region $H(m, n)$ defined as follows. Consider the $2n \times 2n$ square, coordinatized so that the lower left corner is at the origin and the upper right corner has coordinates $(2n, 2n)$. Divide the square into two congruent halves by the two-unit segments

$$\begin{aligned} & \{[2t, 2t+2] \times \{2n - (2t+1)\} : t = 0, \dots, n-1\} \\ & \cup \{\{2n - 2t\} \times [2t-1, 2t+1] : t = 1, \dots, n-1\}. \end{aligned}$$

Now remove the center $2m \times 2m$ square from the region. This leaves two congruent regions, denote each by $H(m, n)$.

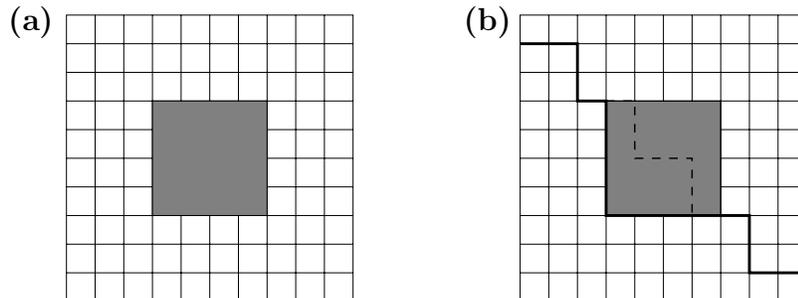


Figure 1: (a) The holey square $\mathcal{H}(2, 5)$. Throughout this paper, shading indicates a portion of the figure that is excluded from the region. (b) The region below the heavy line is $H(2, 5)$.

Thus the holey square conjecture reduces to determining the parity of the number of domino tilings of $H(m, n)$. This paper answers the holey square conjecture in the affirmative via a theorem about tiling parity with applications beyond the problem of the holey square. We demonstrate some of these other consequences.

As this paper solely concerns domino tilings, all tilings discussed can be assumed to be domino tilings. Following Pachter's notation in [3], we write $\#_2 R$ for the parity of the number of tilings of R . A *region* is the dual of a finite connected induced subgraph of \mathbb{Z}^2 . When only concerned with the configuration of part of a region, we may only draw this portion, indicating that the undrawn portion is arbitrary.

2 A tiling parity result

This section presents a theorem regarding the parity of the number of tilings of a region. The result depends only on a local property of the region, and makes no further restrictions. Before stating this property, a few definitions are necessary.

Definition 1. A region R has an $(\{s, t\}; 1)$ -corner if there is a convex corner in R where the segments bounding this corner have lengths s and t . For $p > 1$ and $\min\{s, t\} \geq 2$, an $(\{s, t\}; p)$ -corner is a $(\{1, s\}; 1)$ -corner, a $(\{1, t\}; 1)$ -corner, and $p - 2$ distinct $(\{1, 1\}; 1)$ -corners configured as in Figure 2.

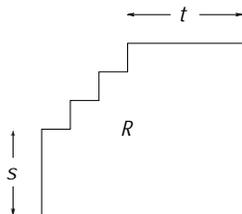


Figure 2: An $(\{s, t\}; 4)$ -corner.

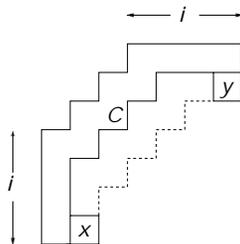
Definition 2. If the segment of length s in an $(\{s, t\}; p)$ -corner forms an $(\{s, t'\}; p')$ -corner at its other endpoint, then each of these corners is *walled* at s .

Definition 3. An $(\{i, j\}; p)$ -strip is a subregion of $i + j + 2p - 3$ squares that has an $(\{i, j\}; p)$ -corner.

The local property mentioned earlier requires that a particular subregion not have any “inconvenient holes.” We use a notion of completion to describe the avoidance of such holes, and precisely define it as follows.

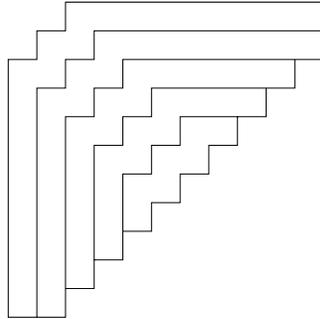
Definition 4. An $(\{s, t\}; p)$ -corner in a region R is *2-complete* if $2 \leq \min\{s, t\}$. For $2 < i \leq \min\{s, t\}$, the corner is *i-complete* if the following conditions are met:

- (a) Let C be the $(\{i, i\}; p)$ -strip in the $(\{s, t\}; p)$ -corner. Let x and y be the two squares adjacent to the ends of C but not along the edges forming the $(\{s, t\}; p)$ -corner. If either x or y is in R , then the $(\{i - 1, i - 1\}; p)$ -strip between x and y , inclusively, all of whose squares are adjacent to C , must also be a subregion of R .



- (b) Consider the $(\{s', t'\}; p)$ -corner formed by removing C from R . If $2 \leq i - 2 \leq \min\{s', t'\}$, then this corner must be j -complete for $j = 2, \dots, i - 2$.

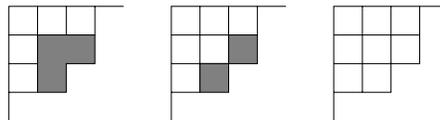
In the process of determining whether or not an $(\{s, t\}; p)$ -corner is k -complete, the largest potential subregion of R to be examined is



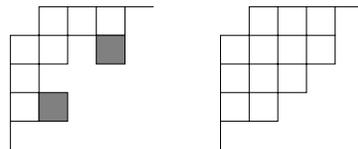
where there are $\lceil k/2 \rceil$ strips: one $(\{k, k\}; p)$ -strip, one $(\{k - 1, k - 1\}; p)$ -strip, one $(\{k - 3, k - 3\}; p)$ -strip, one $(\{k - 5, k - 5\}; p)$ -strip, \dots , concluding with a $(\{3, 3\}; p)$ -strip if k is even, or a $(\{2, 2\}; p)$ -strip if k is odd.

Definition 5. For $2 \leq k \leq \min\{s, t\}$, an $(\{s, t\}; p)$ -corner is *complete up to k* if that corner is i -complete for $i = 2, \dots, k$.

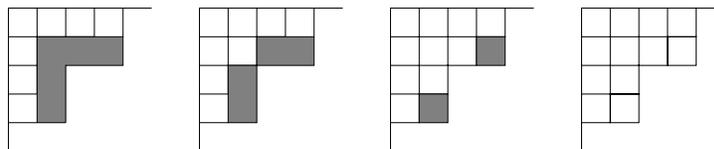
For example, if R has an $(\{s, t\}; 1)$ -corner that is complete up to 3, then this corner must have one of the following forms.



Similarly, an $(\{s, t\}; 2)$ -corner that is complete up to 3 is as follows.



If an $(\{s, t\}; 1)$ -corner is complete up to 4, then it must also be complete up to 3 so the above drawings together with the definition of 4-completeness require that all such corners are depicted below.



Proof. Induct on $n - m$. The base case $n = m + 1$ is trivial, as $\#_2H(m, m + 1) = 1$ and $\mathcal{H}(m, m + 1)$ can be tiled in two ways. Now assume that $\#_2H(m, n) = 1$. Consider the region $H'(m, n + 1)$ which has a $(\{2n, 2n + 1\}; 1)$ -corner walled at $2n$ and complete up to $2n$. Apply the parity theorem to this corner, specifically (2). The subregion indicated by the right-hand side of (2) is actually $H(m, n)$ reflected across the line $y = x$. Therefore $\#_2H'(m, n + 1) = \#_2H(m, n) = 1$. This completes the proof since $\#_2H'(m, n + 1) = \#_2H(m, n + 1)$, answering affirmatively the question posed by Early, and giving a combinatorial meaning to the odd factor in $\#\mathcal{H}(m, n)$. □

Analogous to $\mathcal{H}(m, n)$, let $\mathcal{H}^o(m, n)$ be the $(2n + 1) \times (2n + 1)$ square with a centered hole of size $(2m + 1) \times (2m + 1)$.

Corollary 2. *For all m and $n > m$, $\#\mathcal{H}^o(m, n) = 2^{n-m}(2k'_{m,n} + 1)^2$.*

Proof. The proof is analogous to the proof of the previous corollary, and once again the odd factor $2k'_{m,n} + 1$ is the number of domino tilings of a particular region. □

4 Further applications of the parity theorem

In addition to determining the number of domino tilings of the holey square, the parity theorem can be applied to other regions. One easy consequence is the following.

Corollary 3. *If a region R has an $(\{s, s\}; p)$ -corner that is complete up to s and walled at s along both sides, then $\#R$ is even.*

Proof. Both of the tilings depicted in (1) are impossible, so both terms on the right side of the equation are zero. □

P. W. Kasteleyn [2] gave an exact formula for the number of tilings of an $a \times b$ rectangle, denoted $N(a, b)$. When a and b are both even, this is:

$$\#N(a, b) = \prod_{i=1}^{a-2} \prod_{j=1}^{b-2} \left(4 \cos^2 \frac{i\pi}{a+1} + 4 \cos^2 \frac{j\pi}{b+1} \right). \quad (5)$$

Kasteleyn matrices provide a more general procedure for determining the number of tilings of particular regions, but, unfortunately, many consequences that can be proved using Kasteleyn matrices do not have combinatorial interpretations. Combinatorialists desire to remedy such situations, and perhaps also use combinatorial methods to obtain results that were not apparent by solely employing Kasteleyn matrices. For example, the parity of $\#N(a, b)$ is not at all obvious from (5). However, each corner of $N(a, b)$ is complete up to $\min\{a, b\}$, so we can apply the parity theorem to obtain the following result.

Corollary 4. For all positive integers k and n , $\#N(kn, (k + 1)n)$ is odd.

Proof. To describe the process more precisely, suppose that $N(kn, (k + 1)n)$ is oriented so that the sides of length kn are vertical. Apply the parity theorem to the upper left corner to remove the $(\{kn, kn + 1\}; 1)$ -strip. Similarly, if the values of k and n are sufficiently large, remove the $(\{n - 1, n\}; 1)$ -strip from the upper right corner, the $(\{(k - 1)n, (k - 1)n + 1\}; 1)$ -strip from the lower right corner, and the $(\{2n - 2, 2n - 1\}; 1)$ -strip from the lower left corner (where each of these corners are of subsequent subregions of $N(kn, (k + 1)n)$).

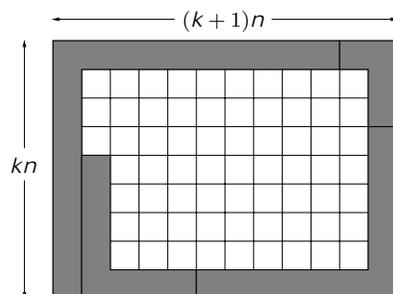


Figure 3: $N(kn, (k + 1)n)$ with the first four removed strips marked, as implied by the parity theorem.

The process of removing each strip can be summarized in the following table describing how many squares are removed from the sides of the region, starting with $N(kn, (k + 1)n)$, after each application of the parity theorem.

App.	Left	Top	Right	Bottom
1	kn	$kn + 1$	—	—
2	—	$n - 1$	n	—
3	—	—	$(k - 1)n$	$(k - 1)n + 1$
4	$2n - 1$	—	—	$2n - 2$
\vdots	\vdots	\vdots	\vdots	\vdots
$4j - 3$	$(k - 2j + 2)n$	$(k - 2j + 2)n + 1$	—	—
$4j - 2$	—	$(2j - 1)(n - 1)$	$(2j - 1)(n - 1) + 1$	—
$4j - 1$	—	—	$(k - 2j + 1)n$	$(k - 2j + 1)n + 1$
$4j$	$2j(n - 1) + 1$	—	—	$2j(n - 1)$

Continue to apply the parity theorem until the $(\{i, i + 1\}; 1)$ -strip removed is from an $(\{i, i + 1\}; 1)$ -corner. To determine when this might happen, we need to solve any of the following equations, where the term subtracted from the left side of each refers to the squares occupied by previously removed corners.

$$kn - (2j - 1) = 2j(n - 1) + 1 \tag{6}$$

$$(k + 1)n - (2j - 2) = (k - 2j + 2)n + 1 \tag{7}$$

$$kn - (2j - 2) = (2j - 1)(n - 1) + 1 \quad (8)$$

$$(k + 1)n - (2j - 1) = (k - 2j + 1)n + 1 \quad (9)$$

These equations correspond to the final $(\{i, i + 1\}; 1)$ -strip being removed from the lower left corner, the upper left corner, the upper right corner, and the lower right corner, respectively. The solution to (6) is $n = 0$ or $k = 2j$, the solution to (7) is $n = 1$ or $2j - 1 = 0$, the solution to (8) is $n = 0$ or $k = 2j - 1$, and the solution to (9) is $n = 1$ or $2j = 0$. Since n , k , and j are positive integers, the applications of the parity theorem cease in the manner described above when $n = 1$ or when $2j = k$ or $2j - 1 = k$, depending on the parity of k .

Suppose that $n > 1$, and let k be even, setting j to $k/2$. After removing the $(2k)^{\text{th}}$ strip, we are left with the subregion of $N(kn, (k + 1)n)$ formed by removing the top j rows, the bottom j rows, the left $j + 1$ columns, and the right j columns. This is $N(k(n - 1), (k + 1)(n - 1))$, so

$$\#_2N(kn, (k + 1)n) = \#_2N(k(n - 1), (k + 1)(n - 1)). \quad (10)$$

If, on the other hand, k is odd, let $j = (k + 1)/2$. After removing the $(2k)^{\text{th}}$ strip from $N(kn, (k + 1)n)$, the resulting subregion is $N(kn, (k + 1)n)$ with the top j rows, the bottom $j - 1$ rows, the left j columns, and the right j columns removed. Since $k = 2j - 1$, this once again gives (10).

Therefore, $\#_2N(kn, (k + 1)n) = \#_2N(k, k + 1)$ for all positive integers k and n . Applying the parity theorem once to any corner of $N(k, k + 1)$ indicates that

$$\#_2N(k, k + 1) = \#_2N(k - 1, k)$$

for all $k > 1$. Therefore $\#_2N(kn, (k + 1)n) = \#_2N(1, 2) = 1$ for all positive integers k and n .

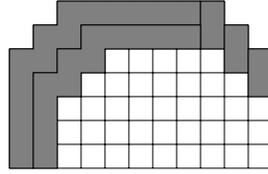
□

As mentioned previously, there are other ways to obtain this result, for example using Kasteleyn's formula. However, these methods tend to be much more analytic, and thus somewhat less intuitively clear, than the combinatorial proof presented here.

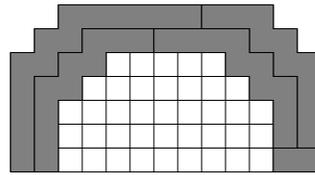
Corollary 4 studied $(\{s, t\}; 1)$ -corners, and we conclude this section by considering more general types of corners. Let $T(i, j, p)$ be the region with $i + p - 1$ rows, whose rows from top to bottom consist of the following number of squares: $j, j + 2, \dots, j + 2(p - 1), \dots, j + 2(p - 1)$, where there are i rows of $j + 2(p - 1)$ boxes. Similarly, let $D(i, j, p)$ be the region with $i + 2(p - 1)$ rows, whose rows from top to bottom consist of the following number of squares: $j, j + 2, \dots, j + 2(p - 1), \dots, j + 2(p - 1), j + 2(p - 1) - 2, \dots, j + 2, j$, where there are i rows of $j + 2(p - 1)$ boxes. In each of these regions, the centers of the rows are aligned vertically. For $T(i, j, p)$ to have an even number of squares (and hence be possibly tilable), either j or $i + p - 1$ must be even. Similarly, for $D(i, j, p)$ to have an even number of squares, either j or i must be even.

The parity theorem applies to $T(i, j, p)$ and $D(i, j, p)$, and examples of such applications are given below.

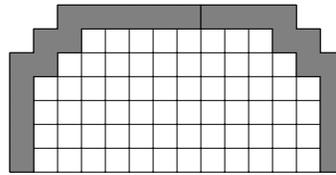
- (c) Apply the parity theorem twice to see that $\#_2T(k, k+2, p) = \#_2T(k-2, k, p)$, and $\#_2T(1, 3, p) = 0$ for all p , while $\#_2T(2, 4, p) = \#_2T(1, 2, p) = 1$ for all p .



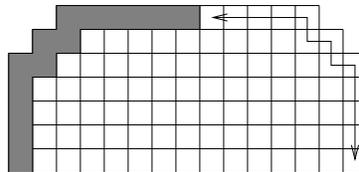
- (d) Several applications of the parity theorem imply that $\#_2T(k, 2k-1, p) = \#_2T(k-2, 2(k-2)-1, p)$, and this is 0 for both $k=1$ and $k=2$.



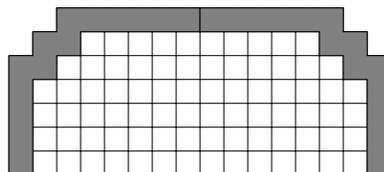
- (e) $\#_2T(k, 2k, p) = \#_2T(k-1, 2(k-1), p)$ and $\#_2T(1, 2, p) = 1$.



- (f) After applying the parity theorem once, apply Corollary 3.



- (g) Two applications of the parity theorem imply that $\#_2T(k, 2k+2, p) = \#_2T(k-1, 2k, p)$, and $\#_2T(1, 4, p) = 1$ for all p .



□

Corollary 6. (a) $\#_2D(k, k, p) = 0$.

