# Coding parking functions by pairs of permutations

Yurii Burman

Independent University of Moscow, 121002, 11, B.Vlassievsky per., Moscow, Russia burman@mccme.ru

Michael Shapiro

Department of Mathematics, Michigan State University East Lansing, MI 48824, Michigan mshapiro@math.msu.edu

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#### Abstract

We introduce a new class of admissible pairs of triangular sequences and prove a bijection between the set of admissible pairs of triangular sequences of length nand the set of parking functions of length n. For all u and v = 0, 1, 2, 3 and all  $n \leq 7$  we describe in terms of admissible pairs the dimensions of the bi-graded components  $h_{u,v}$  of diagonal harmonics  $\mathbb{C}[x_1, \ldots, x_n; y_1, \ldots, y_n]/S_n$ , i.e., polynomials in two groups of n variables modulo the diagonal action of symmetric group  $S_n$ .

### 1 Introduction

A sequence  $p = (p_0, \ldots, p_{n-1})$  is called a *parking function* if it is majorized by a permutation, that is, if there exists a permutation (one-to-one mapping)  $\sigma$  of the set  $\{0, 1, \ldots, n-1\}$  such that  $p_0 \leq \sigma(0), \ldots, p_{n-1} \leq \sigma(n-1)$ . The sequence  $p = (p_0, \ldots, p_{n-1})$  is a parking function if and only if for every  $s = 0, \ldots, n-1$  it contains at least s+1 terms  $p_i$  satisfying the inequality  $p_i \leq s$ . The set of parking functions with n terms will be denoted PF<sub>n</sub>.

Parking functions are a popular subject in combinatorics. Taking their name from a problem of car parking along a one-way street (see [1]), they attracted attention after the following theorem had been proved:

**Theorem 1.1 (Kreweras, [3], 1977).** For every  $k, 0 \le k \le n(n-1)/2$ , there exists a one-to-one correspondence  $\kappa_n$  between the set  $PF_n$  and the set of  $T_n$  trees with n+1numbered vertices such that if  $p_0 + \cdots + p_n = u$  then the tree  $D = \kappa_n(p)$  has exactly n(n-1)/2 - u inversions. In particular, the total number of parking functions is equal to the total number of trees, that is,  $(n+1)^{n-1}$ . A tree here means a connected graph without cycles, whose vertices are numbered  $0, 1, \ldots, n$ . We say that a pair of vertices (i, j) forms an inversion if i < j but the path joining the vertex i with the vertex 0 passes through j. The paper [3] contains an explicit construction of the correspondence involved. (Note that Kreweras's suites majeures differ formally from parking functions we have defined:  $(q_0, \ldots, q_{n-1})$  is a suite majeure if  $(n - q_0, \ldots, n - q_{n-1})$  is a parking function.)

The permutation group  $\Sigma_n$  acts in a natural way on the set  $PF_n$ . Consider a vector space of dimension  $(n + 1)^{n-1}$  with the basis  $e_p$  whose elements are numbered by the parking functions  $p \in PF_n$ . This space carries a natural linear representation of  $\Sigma_n$ ; denote this representation  $\mathcal{P}_n$ . Define a weight w(p) of the parking function p as  $w(p) = n(n-1)/2 - (p_0 + \cdots + p_{n-1})$ . The permutation group action preserves the weight, and therefore  $\mathcal{P}_n$  becomes a graded representation.

Another instance of parking functions (and the main inspiration of this paper) is the following theorem conjectured first in [1] and proved later in a series of works by the same author, see [2] and references therein. Consider a natural action of the permutation group  $\Sigma_n$  on the direct product  $V_n = (\mathbb{C}^2)^n$ . Let  $\mathbb{C}[V_n]$  be the ring of polynomials on  $V_n$ , and  $J_n \subset \mathbb{C}[V_n]$  be the ideal generated by  $\Sigma_n$ -invariant polynomials of positive degree. The factor  $R_n = \mathbb{C}[V_n]/J_n$  is called a module of diagonal harmonics. It is a doubly-graded module: if one denotes arguments of the polynomial  $f \in \mathbb{C}[V_n]$  as  $x_1, y_1, \ldots, x_n, y_n$   $(x_i, y_i)$  being coordinates in the *i*-th copy of  $\mathbb{C}^2$ ) then the gradings are the total degree of f with respect to all  $x_i$  and its total degree with respect to all  $y_i$ . Either grading makes  $R_n$  a graded representation of  $\Sigma_n$ .

**Theorem 1.2 (Haiman, [2], 2000).**  $R_n$  is isomorphic, as a graded representation of  $\Sigma_n$ , to the representation  $\mathcal{P}_n$  tensored by the sign representation  $\epsilon_n$ .

In particular, the dimension of the homogeneous component of  $R_n$  of the grading k is equal to the number of parking functions p with  $p_0 + \cdots + p_{n-1} = n(n-1)/2 - k$ , or to the number of trees with n+1 numbered vertex having exactly k inversions.

Note that in fact the representation  $R_n$  is bi-graded, but Theorem 1.2 ignores the second grading. There are explicit formulas for dimensions of the bihomogeneous components of  $R_n$  (see [2]) but they have nothing to do with trees and parking functions. Nevertheless, the theorem suggests that the representation  $\mathcal{P}_n$  also can be made doubly graded — that is, the sets of parking functions and trees should carry the second grading, yet unknown, different from the weight defined above.

The aim of our project was to find an elementary approach to the above bi-grading. Unfortunatelly we failed to define the second grading. This paper is a description of steps made in this direction, some of them being rigorously proved statements, and some, numerical observations. We start at sections 2 and 3 with a combinatorial construction encoding parking functions by pairs of permutations satisfying some admissibility condition. The last section contains some data shading light on the relation of this construction to Theorem 1.2 above.

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## 2 Main definitions

Throughout this paper we use the following conventions:

- 1. All the numbers mentioned are integers, unless otherwise stated.
- 2. All sequences are indexed starting from zero; giving numbers to elements of a finite set, we also start from zero.

### 2.1 Permutations and triangular sequences

We call a sequence  $a = (a_0, \ldots, a_{n-1})$  triangular if an inequality  $0 \le a_i \le i$  is satisfied for all  $i = 0, \ldots, n-1$ . Denote  $A_n$  the set of all triangular sequences of length n. Apparently, the cardinality of  $A_n$  is n!; this allows to put it into a one-to-one correspondence with the set  $\Sigma_n$  of all permutations of the set  $\{0, \ldots, n-1\}$ .

There are several explicit constructions for this correspondence. We will usually use the correspondence  $\alpha_n : A_n \to \Sigma_n$  defined inductively as follows. If n = 1, there is only one triangular sequence, only one permutation, and only one correspondence  $\alpha_1$  between them. Now let  $\alpha_{n-1}$  be defined and let  $\sigma' = (\sigma'(0), \ldots, \sigma'(n-2)) = \alpha_{n-1}(a_0, \ldots, a_{n-2})$ . Take the number n - 1 and insert it into the line  $(\sigma'(0), \ldots, \sigma'(n-2))$  between  $\sigma'(n - 2 - a_{n-1})$  and  $\sigma'(n - 1 - a_{n-1})$ ; the resulting sequence will represent the permutation  $\sigma = \alpha_n(a_0, \ldots, a_{n-1})$ . Formally,

$$\sigma(i) = \begin{cases} \sigma'(i), & \text{if } i \le n - 2 - a_{n-1}, \\ n - 1, & \text{if } i = n - 1 - a_{n-1}, \\ \sigma'(i - 1), & \text{if } i \ge n - a_{n-1}. \end{cases}$$

It is easy to see that  $\alpha_n$  is indeed a one-to-one correspondence. The inverse mapping can be described by the following rule: if  $a = \alpha_n^{-1}(\sigma)$  then  $a_i$  equals the number of j > isuch that  $\sigma(j) < \sigma(i)$ . A pair (i, j) such that j > i but  $\sigma(j) < \sigma(i)$  is called an *inversion* of the permutation  $\sigma$ ; the total number of inversions of the permutation  $\alpha_n(a)$  is thus equal to  $a_0 + \cdots + a_{n-1}$ .

### 2.2 Triangular sequences and parking functions

Consider a pair of triangular sequences  $k = (k_0, \ldots, k_{n-1}), l = (l_0, \ldots, l_{n-1}) \in A_n$  such that  $l_s \leq k_s$  for all  $s = 0, \ldots, n-1$ . Define a sequence  $\beta_n(k, l) = p = (p_0, \ldots, p_{n-1})$  by induction as follows. Let  $\beta_{n-1}$  be defined, and  $(p'_0, \ldots, p'_{n-1}) = \beta_{n-1}(k', l')$  where  $k' = (k_0, \ldots, k_{n-2}), l' = (l_0, \ldots, l_{n-2}) \in A_{n-2}$ . Now take the number  $k_{n-1}$  and insert it

into the line  $(p'_0, \ldots, p'_{n-1})$  between positions number  $k_{n-1} - l_{n-1} - 1$  and  $k_{n-1} - l_{n-1}$ ; the resulting sequence will be  $\beta_n(k, l)$ . Formally, if  $p = \beta_n(k, l)$  then

$$p_i = \begin{cases} p'_i, & \text{if } i \le k_{n-1} - l_{n-1} - 1, \\ k_{n-1}, & \text{if } i = k_{n-1} - l_{n-1}, \\ p'_{i-1}, & \text{if } i \ge k_{n-1} - l_{n-1} + 1. \end{cases}$$

It is easy to see that  $\beta_n(k, l)$  is a parking function for all k, l.

An equivalent description of this algorithm is as follows: let us have, first, n empty positions numbered from 0 to n-1, left to right. Take  $k_{n-1}$  and place it to the position number  $k_{n-1} - l_{n-1}$ . Then re-number empty positions using numbers from 0 to n-2, skipping the occupied position. Then take  $k_{n-2}$  and place it to the empty position whose (new) number is  $k_{n-2} - l_{n-2}$ . Again, re-number the empty positions using numbers from 0 to n-3, etc.

## 3 Admissible pairs of triangular sequences

Let, again,  $k = (k_0, \ldots, k_{n-1})$  and  $l = (l_0, \ldots, l_{n-1})$  be triangular sequences. We say that a pair of integers  $(i, j), 0 \le i < j \le n-1$  forms an *irregular position* for a pair  $k, l \in A_n$  if  $l_i > l_j$  and  $k_j \le i$ . The pair  $k, l \in A_n$  is called *admissible* if  $l_s \le k_s$  for all  $s = 0, \ldots, n-1$ , and no irregular positions exist. Denote Adm<sub>n</sub>  $\subset A_n \times A_n$  the set of all admissible pairs.

The next statement is the main proved result of the paper.

**Theorem 3.1.** The mapping  $\beta_n$  provides a one-to-one correspondence between the sets  $\operatorname{Adm}_n$  and  $\operatorname{PF}_n$ .

**Corollary 3.2.** There are  $(n+1)^{n-1}$  admissible pairs of triangular sequences.

To prove Theorem 3.1 we need two lemmas. Let  $p \in PF_n$  be a parking function, and r be a number such that  $p_r$  is the maximal term of the sequence p; if there are several such terms, take the smallest r possible. Let (k, l) be an admissible pair such that  $p = \beta_n(k, l)$ , and let s be a number such that  $k_s = p_r$  is sent to position r by the mapping  $\beta_n$  (in other words,  $s = \sigma_r$  where  $\sigma = \alpha_n(k - l)$ , cf. Section 2).

Consider the set U of all i,  $p_r \leq i \leq n-1$ , such that for all j,  $p_r \leq j \leq i$ , the inequality  $l_j \leq (n-1-r) + (p_r-i)$  takes place.

**Lemma 3.3.**  $s = \max U$ .

Proof. By the choice of s, we have  $k_j \leq k_s = p_r$  for every j > s. Now if  $l_j < l_s$  then  $k_j \leq k_s \leq s$ , and (s, j) is an irregular position. For an admissible pair (k, l) no such positions exist, and so  $l_j \geq l_s$  for all j > s. This inequality means that the mapping  $\beta_n$  sends every term  $k_j$ , j > s, to a position left of (less than) r, and therefore  $r = (n-1-s) + (k_s - l_s) = (n-1-s) + (p_r - l_s)$ . So,  $l_s = (n-1-r) + (p_r - s)$ .

Now let j be such that  $p_r \leq j \leq s - 1$ . Again, we have  $l_j \leq l_s$ , because (k, l) is an admissible pair. Therefore,  $l_j \leq (n - 1 - r) + (p_r - s)$  which means that  $s \in U$ .

Suppose there exists  $t \in U$  such that t > s. Then  $p_r \leq s \leq t - 1$ , and there holds the inequality  $l_s = (n - 1 - r) + p_r - s \leq (n - 1 - r) + p_r - t$  a contradiction.

Now delete the r-th element of the parking function p obtaining a sequence p'. Also, delete the s-th elements from sequences k and l resulting in k' and l', respectively.

#### Lemma 3.4.

- 1. The sequence p' is a parking function:  $p' \in PF_{n-1}$ .
- 2. The sequences k' and l' are triangular and form an admissible pair:  $(k', l') \in Adm_{n-1}$ .
- 3.  $\beta_n(k', l') = p'$ .

*Proof.* 1. Let  $\sigma$  be a permutation of the set  $\{0, \ldots, n-1\}$  majorizing the sequence p. Since  $p_r$  is the maximal term of p, then without loss of generality  $\sigma(r) = n - 1$ . Deleting the *r*-th term from  $\sigma$  one obtains a permutation  $\sigma'$  of the set  $\{0, \ldots, n-2\}$  majorizing p'.

2. If j < s then  $k'_j = k_j \leq j$ . As we noticed in the proof of Lemma 3.3,  $k_j \leq k_s$  for all  $j \geq s$ , and therefore  $k'_j = k_{j-1} \leq s \leq j$  for such j, too. Thus, the sequence k' is triangular. The inequalities  $l'_i \leq k'_i$  imply that the sequence l' is also triangular. Every irregular position (i, j) for (k', l') would be irregular for (k, l), too, and thus  $(k', l') \in \text{Adm}_{n-1}$ . 3. Evident.

Proof of Theorem 3.1. 1. Existence — prove that for every parking function  $p \in PF_n$ there exists an admissible pair  $(k, l) \in Adm_n$  such that  $p = \beta_n(k, l)$ . Use the induction by n, the base n = 1 being evident. To make the induction step, define the number rand the parking function  $p' \in PF_{n-1}$  as described in the beginning of this section. By induction hypothesis, there exists a pair  $(k', l') \in Adm_{n-1}$  such that  $p' = \beta_{n-1}(k', l')$ . Let U be the set of all  $i, 0 \leq i \leq n-1$ , such that for all  $j, p_r \leq j \leq i-1$ , the inequality  $l'_j \leq n-1-r+p_r-i$  takes place. Define the number s as the maximal element of U, assuming s = 0 if  $U = \emptyset$ . Now insert the terms  $k_s = p_r$  and  $l_s = (n-1-r)+p_r-s$  into k' and l' getting k and l, respectively. We are to prove now that (k, l) is an admissible pair of triangular sequences satisfying  $\beta_n(k, l) = p$ .

By the choice of s, the inequality  $k_s = p_r \leq s$  holds, which means that the sequence k is triangular. Prove that  $l_s \leq k_s$  (triangularity of l would follow). This inequality is equivalent to  $n - 1 - r \leq s$ , so if  $n - 1 - r \leq p_r$  then it follows from the previous one.

Suppose  $n-1-r > p_r$ . Then for every  $j, p_r \le j \le n-1-r$ , one has  $l'_j \le k'_j < p_r+1 = (n-1-r) + p_r - (n-2-r)$ . This means that the number n-1-r belongs to the set U, and therefore, again,  $n-1-r \le s$ . So,  $l_s \le k_s$  and l is triangular.

Prove now that  $\beta_n(k, l) = p$ . Show first that  $l_j \ge l_s$  for all j > s. Suppose that  $l_j < l_s$  for some j > s. By the induction hypothesis, (k', l') is a admissible pair, and therefore (s, j - 1) is not an irregular position for it. The inequality  $k'_{j-1} = k_j \ge s + 1 > p_r$  is impossible, so,  $l_{s+1} = l'_s \le l'_{j-1} = l_j$ , and therefore  $l_{s+1} < l_s = (n - 1 - r) + p_r - s$ , or  $l'_s \le (n - 1 - r) + p_r - (s + 1)$ . By the choice of s, the number s + 1 is not an element of

the set U. This means that there exists t,  $p_r \leq t < s$ , such that  $l_t > l_{s+1}$ . The inequality  $k'_s \geq t+1 > p_r$  is impossible, so in this case (t-1,s) is an irregular position for (k', l')— a contradiction.

So,  $l_j \ge l_s$  for all j > s. Since  $k_s = p_r$  is the maximal term of the sequence p, we have also  $k_j \le k_s$  all j > s. This implies that the mapping  $\beta_n$  sends all the  $k_j$  with j > s to positions left of r, and therefore  $k_s = p_r$  is sent to position r. It follows now from the induction hypothesis  $(p' = \beta_{n-1}(k', l'))$  that  $p = \beta_n(k, l)$ .

We proved already that the sequences k and l are triangular, and  $l_i \leq k_i$  for all  $i = 0, \ldots, n-1$ . Prove that there are no irregular positions for k, l and therefore  $(k, l) \in \text{Adm}_n$ . Let (u, v) be such a position; consider several cases:

Case 1. u < v < s. Then (u, v) is irregular for (k', l'), too — a contradiction. The same argument applies to cases u < s < v (the position (u, v - 1)) and s < u < v (the position (u - 1, v - 1)).

Case 2. u = s < v. This is impossible because, as we proved earlier,  $l_v \leq l_s$ .

Case 3. u < v = s. This means that  $p_r \leq u$  and thus  $l'_u = l_u > l_s = (n-1-r) + p_r - s$ , which is impossible by the definition of the set U.

2. Uniqueness. Again, use induction by n, the base n = 1 being evident. Let  $\beta_n(k^{(1)}, l^{(1)}) = \beta_n(k^{(2)}, l^{(2)}) = p \in \mathrm{PF}_n$  where  $(k^{(1)}, l^{(1)})$  and  $(k^{(2)}, l^{(2)})$  are admissible pairs. Choose the number r as above, and let  $s_1, s_2$  be numbers such that the mapping  $\beta_n$  applied to pairs  $(k^{(1)}, l^{(1)})$  and  $(k^{(2)}, l^{(2)})$  sends  $k_{s_1}^{(1)}$  and  $k_{s_2}^{(2)}$ , respectively, to position r. Let  $\tilde{k}^{(1)}$  and  $\tilde{l}^{(1)}$  be sequences obtained by deletion of the  $s_1$ -th term from  $k^{(1)}$  and  $l^{(1)}$ , and similarly  $\tilde{k}^{(2)}$  and  $\tilde{l}^{(2)}$ . By Lemma 3.4,  $\beta_{n-1}(k^{(1)}, l^{(1)}) = \beta_{n-1}(k^{(2)}, l^{(2)})$ , and by the induction hypothesis,  $\tilde{k}^{(1)} = \tilde{k}^{(2)}$ ,  $\tilde{l}^{(1)} = \tilde{l}^{(2)}$ .

The pair  $(k^{(1)}, l^{(1)})$  is admissible, and  $k_i^{(1)} \leq k_{s_1}^{(1)}$  for all *i*. It implies that  $l_j^{(1)} \geq l_{s_1}^{(1)}$  for all  $j > s_1$  and  $l_j^{(1)} \leq l_{s_1}^{(1)}$  for all  $j, p_r \leq j \leq s_1 - 1$ . The same is true for  $l^{(2)}$ . As we know, the sequences  $l^{(1)}$  and  $l^{(2)}$  become the same after deletion of the  $s_1$ -th and the  $s_2$ -th terms, respectively. Hence, if  $l_{s_1}^{(1)} > l_{s_2}^{(2)}$  then  $s_1 > s_2$ , and vice versa.

On the other hand, the mapping  $\beta_n$  sends all the  $k_j^{(1)}$  with  $j > s_1$  to positions left of r, and therefore  $r = (n - 1 - s_1) + (p_r - l_{s_1}^{(1)})$ . A similar equation is true for  $l^{(2)}$ , hence,  $l_{s_1}^{(1)} + s_1 = l_{s_2}^{(2)} + s_2$ . So,  $s_1 = s_2$  and  $l_{s_1}^{(1)} = l_{s_2}^{(2)}$  — uniqueness is proved.

### 4 Admissible pairs and diagonal harmonics

Here we present some relation between the construction of Theorem 3.1 and the module  $R_n$  of diagonal harmonics described in Section 1. Let  $H_{u,v}$  be the bihomogeneous component of the module  $R_n$  of bi-degree (u, v); denote  $h_{u,v}$  its dimension.

Define now the four sets  $Y_0, Y_1, Y_2, Y_3 \subset A_n$  of triangular sequences as follows.

- 0. The set  $Y_0$  consists of only one sequence, namely  $(0, 0, \ldots, 0)$ .
- 1. The set  $Y_1$  consists of sequences  $(l_0, \ldots, l_{n-1})$  such that  $l_0 = \cdots = l_{n-2} = 0$  and  $l_{n-1} \ge 1$ .

- 2. The set  $Y_2$  consists of sequences  $(l_0, \ldots, l_{n-1})$  such that  $l_0 = \cdots = l_{n-3} = 0$  and  $l_{n-1} \ge l_{n-2} \ge 1$ .
- 3.  $Y_3$  is a union of two sets,  $Y'_3$  and  $Y''_3$ . The set  $Y'_3$  consists of sequences  $(l_0, \ldots, l_{n-1})$  such that  $l_0 = \cdots = l_{n-3} = 0$  and  $1 \leq l_{n-2} > l_{n-1}$ . The set  $Y''_3$  consists of sequences  $(l_0, \ldots, l_{n-1})$  such that  $l_0 = \cdots = l_{n-4} = 0$  and  $1 \leq l_{n-3} \leq l_{n-2} \leq l_{n-1} \leq n-2$  (note an additional inequality at the end; triangularity requires only  $l_{n-1} \leq n-1$ ).

Numerical computations made for all  $n \leq 7$  (using tables of  $h_{u,v}$  taken from [1]) give the following observation:

For all u and v = 0, 1, 2, 3 and all  $n \leq 7$  the dimension  $h_{u,v}$  is equal to the number of admissible pairs (k, l) such that  $k_0 + \cdots + k_{n-1} = n(n-1)/2 - u$  and  $l \in Y_v$ .

Thus it can be conjectured that there exists a splitting of the set  $A_n$  into a disjoint union:  $A_n = \bigsqcup_{v=0}^{n(n-1)/2} Y_v$  such that the statement above holds for all u, v (and all n). The authors, though, know neither a construction of  $Y_v$  nor a proof of the conjecture above for small v.

Besides the numerical observations for  $n \leq 7$  there are some more facts supporting the conjecture.

1. Let  $S \subset \mathbb{C}[x_1, \ldots, x_n]$  be the ideal generated by symmetrical polynomials of positive degree. Apparently,  $\mathbb{C}[x_1, \ldots, x_n]/S$  is a graded module isomorphic to  $\bigoplus_u H_{u,0}$ . As it is well known, the dimension  $h_{u,0}$  of its component of gradung u is equal to the number of permutations having exactly u inversions. On the other hand, for v = 0 we have  $l = (0, 0, \ldots, 0)$ , and the admissibility condition does not impose any limitations on k. The results of Section 2 now imply that the observation above is true for v = 0 and all n.

2. Apparently,  $h_{u,v} = h_{v,u}$ , and therefore the dimension  $h_{0,v}$  is equal to the number of permutations with v inversions. The conjecture above implies that  $h_{0,v}$  equals to the number of admissible pairs (k, l) such that  $k_0 + \cdots + k_{n-1} = n(n-1)/2$  and  $l \in Y_v$ . The first equation holds only if  $k_i = i$  for all  $i = 0, \ldots, n-1$ . For such k and the pair (k, l)is admissible for every  $l \in A_n$ . This implies that the number of elements in  $Y_v$  should be equal to the number of permutations with v inversions. It is easy to check that this is true for v = 0, 1, 2, 3 (and all n).

3. It can be easily checked that the answer for  $h_{u,v}$  given by the conjecture satisfies the condition  $h_{u,v} = h_{v,u}$  for all n and all  $u, v \leq 3$  (that is, in all cases when the conjectural values of both  $h_{u,v}$  and  $h_{v,u}$  are known).

4. From Theorem 1.2 we know that  $\sum_{u} h_{u,v}$  equals to the number  $T_v$  of trees with v inversions. The conjecture implies then that the total number of admissible pairs (k, l) with  $l \in Y_v$  should be equal to  $T_v$ , too. A direct computation (see [3] for a formula for  $T_v$ ) shows that this is true for v = 0, 1, 2, 3 (and all n).

The splitting  $A_n = \bigsqcup_{v=0}^{n(n-1)/2} Y_v$ , if known, would provide the second grading on the sets of admissible pairs. The one-to-one correspondences  $\beta_n$  and  $\kappa_n$  mentioned in Section 1 would allow then to define the grading on the set of parking functions and on the set of trees. (Recall that the first grading for an admissible pair (k, l) equals  $n(n-1)/2 - (k_0 + l)$ 



Figure 1: Such trees probably have grading 1:  $c \ge b$ 

 $\dots + k_{n-1}$ ), for a parking function  $p = \beta_n(k, l)$  it is  $w(p) = n(n-1)/2 - (p_0 + \dots + p_{n-1}) = n(n-1)/2 - (k_0 + \dots + k_{n-1})$ , and for a tree  $D = \kappa_n(p)$  it is equal, by Theorem 1.1, to the number of inversions). Thus, by now we are able to describe conjecturally the sets of parking functions and trees of grading 0, 1, 2 and 3 using explicit constructions of  $\beta_n$  (see Section 2) and  $\kappa_n$  (see [3]). To exemplify what happens we give here the answers for gradings 0 and 1:

- 0. Parking functions of grading zero are triangular sequences:  $p_s \leq s$  for all  $s = 0, 1, \ldots, n-1$ . Trees of grading zero are "linear" trees with one branch only.
- 1. A parking function of grading 1 has exactly one term  $p_s$  such that  $p_s > s$ . After deletion of this term the remaining sequence is triangular (i.e. a parking function of grading 0). A tree of grading 1 has exactly one "branching point" — a vertex A having two children, B and C (all the other non-terminal vertices have one child only). The vertex B is terminal and carries the number b which is less or equal to the length of path joining A with the other terminal vertex (see Figure 1).

# References

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