Distance domination and distance irredundance in graphs

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Abstract

A set $D \subseteq V$ of vertices is said to be a (connected) distance k-dominating set of G if the distance between each vertex $u \in V - D$ and D is at most k (and D induces a connected graph in G). The minimum cardinality of a (connected) distance k-dominating set in G is the (connected) distance k-domination number of G, denoted by $\gamma_k(G)$ ($\gamma_k^c(G)$, respectively). The set D is defined to be a total k-dominating set of G if every vertex in V is within distance k from some vertex of D other than itself. The minimum cardinality among all total k-dominating sets of G is called the total k-domination number of G and is denoted by $\gamma_k^t(G)$. For $x \in X \subseteq V$, if $N^k[x] - N^k[X - x] \neq \emptyset$, the vertex x is said to be k-irredundant in X. A set X containing only k-irredundant vertices is called k-irredundant. The k-irredundance number of G, denoted by $ir_k(G)$, is the minimum cardinality taken over all maximal k-irredundant sets of vertices of G. In this paper we establish lower bounds for the distance k-irredundance number of graphs and trees. More precisely, we prove that $\frac{5k+1}{2}ir_k(G) \geq \gamma_k^c(G) + 2k$ for each connected graph G and $(2k+1)ir_k(T) \geq \gamma_k^c(T) + 2k \geq |V| + 2k - kn_1(T)$ for each tree T = (V,E) with $n_1(T)$ leaves. A class of examples shows that the latter bound is sharp. The second inequality generalizes a result of Meierling and Volkmann [9] and Cyman, Lemańska and Raczek [2] regarding γ_k and the first generalizes a result of Favaron and Kratsch [4] regarding ir_1 . Furthermore, we shall show that $\gamma_k^c(G) \leq \frac{3k+1}{2}\gamma_k^t(G) - 2k$ for each connected graph G, thereby generalizing a result of Favaron and Kratsch [4] regarding k = 1.

Keywords: domination, irredundance, distance domination number, total domination number, connected domination number, distance irredundance number, tree 2000 Mathematics Subject Classification: 05C69

1 Terminology and introduction

In this paper we consider finite, undirected, simple and connected graphs G = (V, E) with vertex set V and edge set E. The number of vertices |V| is called the order of G and is denoted by n(G). For two distinct vertices u and v the $distance\ d(u,v)$ between u and v is the length of a shortest path between u and v. If X and Y are two disjoint subsets of V, then the distance between X and Y is defined as $d(X,Y) = \min\{d(x,y) \mid x \in X, y \in Y\}$. The $open\ k$ -neighborhood $N^k(X)$ of a subset $X \subseteq V$ is the set of vertices in $V \setminus X$ of distance at most k from X and the $closed\ k$ -neighborhood is defined by $N^k[X] = N^k(X) \cup X$. If $X = \{v\}$ is a single vertex, then we denote the (closed) k-neighborhood of v by $N^k(v)$ ($N^k[v]$, respectively). The (closed) 1-neighborhood of a vertex v or a set X of vertices is usually denoted by N(v) or N(X), respectively (N[v] or N[X], respectively). Now let U be an arbitrary subset of V and $u \in U$. We say that v is a private k-neighborhoof of v with respect to v if $v \in N^k[u] - N^k[v] - \{u\}$. The private v-neighborhood of v with respect to v will be denoted by v-neighborhood of v-neighborhood

For a vertex $v \in V$ we define the *degree* of v as d(v) = |N(v)|. A vertex of degree one is called a *leaf* and the number of leaves of G will be denoted by $n_1(G)$.

A set $D \subseteq V$ of vertices is said to be a *(connected) distance k-dominating set* of G if the distance between each vertex $u \in V - D$ and D is at most k (and D induces a connected graph in G). The minimum cardinality of a (connected) distance k-dominating set in G is the *(connected) distance k-domination number* of G, denoted by $\gamma_k(G)$ ($\gamma_k^c(G)$, respectively). The distance 1-domination number $\gamma_1(G)$ is the usual domination number $\gamma(G)$. A set $D \subseteq V$ of vertices is defined to be a total k-dominating set of G if every vertex in V is within distance k from some vertex of D other than itself. The minimum cardinality among all total k-dominating sets of G is called the total k-domination number of G and is denoted by $\gamma_k^t(G)$. We note that the parameters $\gamma_k^c(G)$ and $\gamma_k^t(G)$ are only defined for connected graphs and for graphs without isolated vertices, respectively.

For $x \in X \subseteq V$, if $PN^k[x] \neq \emptyset$, the vertex x is said to be k-irredundant in X. A set X containing only k-irredundant vertices is called k-irredundant. The k-irredundance number of G, denoted by $ir_k(G)$, is the minimum cardinality taken over all maximal k-irredundant sets of vertices of G.

In 1975, Meir and Moon [10] introduced the concept of a k-dominating set (called a 'k-covering' in [10]) in a graph, and established an upper bound for the k-domination number of a tree. More precisely, they proved that $\gamma_k(T) \leq |V(T)|/(k+1)$ for every tree T. This leads immediately to $\gamma_k(G) \leq |V(G)|/(k+1)$ for an arbitrary graph G. In 1991, Topp and Volkmann [11] gave a complete characterization of the class of graphs G that fulfill the equality $\gamma_k(G) = |V(G)|/(k+1)$.

The concept of k-irredundance was introduced by Hattingh and Henning [5] in 1995. With k = 1, the definition of an k-irredundant set coincides with the notion of an irredundant set, introduced by Cockayne, Hedetniemi and Miller [1] in 1978. Since then a lot of research has been done in this field and results have been presented by many authors (see [5]).

In 1991, Henning, Oellermann and Swart [8] motivated the concept of total distance domination in graphs which finds applications in many situations and structures which give rise to graphs.

For a comprehensive treatment of domination in graphs, see the monographs by Haynes, Hedetniemi and Slater [6], [7].

In this paper we establish lower bounds for the distance k-irredundance number of graphs and trees. More precisely, we prove that $\frac{5k+1}{2}ir_k(G) \geq \gamma_k^c(G) + 2k$ for each connected graph G and $(2k+1)ir_k(T) \geq \gamma_k(T) + 2k \geq |V| + 2k - kn_1(T)$ for each tree T = (V, E) with $n_1(T)$ leaves. A class of examples shows that the latter bound is sharp. Since $\gamma_k(G) \geq ir_k(G)$ for each connected graph G, the latter generalizes a result of Meierling and Volkmann [9] and Cyman, Lemanska and Raczek [2] regarding γ_k and the former generalizes a result of Favaron and Kratsch [4] regarding ir_1 . In addition, we show that if G is a connected graph, then $\gamma_k^c(G) \leq (2k+1)\gamma_k(G) - 2k$ and $\gamma_k^c(G) \leq \frac{3k-1}{2}\gamma_k^t(G) - 2k$ thereby generalizing results of Duchet and Meyniel [3] for k=1 and Favaron and Kratsch [4] for k=1, respectively.

2 Results

First we show the inequality $\gamma_k^c \leq (2k+1)\gamma_k - 2k$ for connected graphs.

Theorem 2.1. If G is a connected graph, then

$$\gamma_k^c(G) \le (2k+1)\gamma_k(G) - 2k.$$

Proof. Let G be a connected graph and let D be a distance k-dominating set. Then G[D] has at most |D| components. Since D is a distance k-dominating set, we can connect two of these components to one component by adding at most 2k vertices to D. Hence, we can construct a connected k-dominating set $D' \supseteq D$ in at most |D| - 1 steps by adding at most (|D| - 1)2k vertices to D. Consequently,

$$\gamma_k^c(G) \le |D'| \le |D| + (|D| - 1)2k = (2k + 1)|D| - 2k$$

and if we choose D such that $|D| = \gamma_k(G)$, the proof of this theorem is complete.

The results given below follow directly from Theorem 2.1.

Corollary 2.2 (Duchet & Meyniel [3] 1982). If G is a connected graph, then

$$\gamma^c(G) \le 3\gamma(G) - 2.$$

Corollary 2.3 (Meierling & Volkmann [9] 2005; Cyman, Lemańska & Raczek [2] 2006). If T is a tree with n_1 leaves, then

$$\gamma_k(T) \ge \frac{|V(T)| - kn_1 + 2k}{2k + 1}.$$

Proof. Since $\gamma_k^c(T) \geq |V(T)| - kn_1$ for each tree T, the proposition is immediate.

The following lemma is a preparatory result for Theorems 2.5 and 2.7.

Lemma 2.4. Let G be a connected graph and let I be a maximal k-irredundant set such that $ir_k(G) = |I|$. If $I_1 = \{v \in I \mid v \in PN^k[v]\}$ is the set of vertices that have no k-neighbor in I, then

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$$\gamma_k^c(G) \le (2k+1)ir_k(G) - 2k + (k-1)\frac{|I-I_1|}{2}.$$

Proof. Let G be a connected graph and let $I \subseteq V$ be a maximal k-irredundant set. Let

$$I_1 := \{ v \in I \mid v \in PN^k[v] \}$$

be the set of vertices in I that have no k-neighbors in I and let

$$I_2 := I - I_1$$

be the complement of I_2 in I. For each vertex $v \in I_2$ let $u_v \in PN^k[v]$ be a k-neighbor of v such that the distance between v and u_v is minimal and let

$$B := \{u_v \mid v \in I_2\}$$

be the set of these k-neighbors. Note that $|B| = |I_2|$. If w is a vertex such that $w \notin N^k[I \cup B]$, then $I \cup \{w\}$ is a k-irredundant set of G that strictly contains I, a contradiction. Hence $I \cup B$ is a k-dominating set of G.

Note that $G[I \cup B]$ has at most $|I \cup B| = |I_1| + 2|I_2|$ components. From $I \cup B$ we shall construct a connected k-dominating set $D \supseteq I \cup B$ by adding at most

$$|I_2|(k-1) + (|I_1| + \left\lfloor \frac{|I_2|}{2} \right\rfloor - 1)2k + \left\lfloor \frac{|I_2|}{2} \right\rfloor (k-1)$$

vertices to $I \cup B$.

We can connect each vertex $v \in I_2$ with its corresponding k-neighbor $u_v \in B$ by adding at most k-1 vertices to $I \cup B$.

Recall that each vertex $v \in I_2$ has a k-neighbor $w \neq v$ in I_2 . Therefore we need to add at most k-1 vertices to $I \cup B$ to connect such a pair of vertices.

By combining the two observations above, we can construct a k-dominating set $D' \supseteq I \cup B$ from $I \cup B$ with at most $|I_1| + \lfloor |I_2|/2 \rfloor$ components by adding at most $(k-1)|I_2| + (k-1)\lceil |I_2|/2 \rceil$ vertices to $I \cup B$. Since D' is a k-dominating set of G, these components can be joined to a connected k-dominating set D by adding at most $(|I_1| + \lfloor |I_2|/2 \rfloor - 1)2k$ vertices to D'.

All in all we have shown that there exists a connected k-dominating set D of G such that

$$|D| \le |I_1| + 2|I_2| + (k-1)|I_2| + (k-1)\left\lceil \frac{|I_2|}{2} \right\rceil + 2k(|I_1| + \left\lfloor \frac{|I_2|}{2} \right\rfloor - 1)$$

$$\le (2k+1)|I| - 2k + (k-1)\frac{|I_2|}{2}.$$

Hence, if we choose the set I such that $|I| = ir_k(G)$, the proof of this lemma is complete.

Since $|I_2| \leq |I|$ for each k-irredundant set I, we derive the following theorem.

Theorem 2.5. If G is a connected graph, then

$$\gamma_k^c(G) \le \frac{5k+1}{2} i r_k(G) - 2k.$$

The next result follows directly from Theorem 2.5.

Corollary 2.6 (Favaron & Kratsch [4] 1991). If G is a connected graph, then

$$\gamma^c(G) \le 3ir(G) - 2.$$

For acyclic graphs Lemma 2.4 can be improved as follows.

Theorem 2.7. If T is a tree, then

$$\gamma_k^c(T) \le (2k+1)ir_k(T) - 2k.$$

Proof. Let T be a tree and let $I \subseteq V$ be a maximal k-irredundant set. Let

$$I_1 := \{ v \in I \mid v \in PN^k[v] \}$$

be the set of vertices in I that have no k-neighbors in I and let

$$I_2 := I - I_1$$

be the complement of I_2 in I. For each vertex $v \in I_2$ let $u_v \in PN^k[v]$ be a k-neighbor of v such that the distance between v and u_v is minimal and let

$$B := \{ u_v \mid v \in I_2 \}$$

be the set of these k-neighbors. Note that $|B| = |I_2|$. If w is a vertex such that $w \notin N^k[I \cup B]$, then $I \cup \{w\}$ is a k-irredundant set of G that strictly contains I, a contradiction. Hence $I \cup B$ is a k-dominating set of G.

Note that $T[I \cup B]$ has at most $|I \cup B| = |I_1| + 2|I_2|$ components. From $I \cup B$ we shall construct a connected k-dominating set $D \supseteq I \cup B$ by adding at most

$$(2k-1)|I_2| + 2k(|I_1|-1)$$

vertices to $I \cup B$. To do this we need the following definitions. For each vertex $v \in I_2$ let P_v be the (unique) path between v and u_v and let x_v be the predecessor of u_v on P_v . Let $I_2 = S \cup L_1 \cup L_2$ be a partition of I_2 such that

$$S = \{ v \in I_2 \mid d(v, u_v) = 1 \}$$

is the set of vertices of I_2 that are connected by a 'short' path with u_v ,

$$L_1 = \{ v \in I_2 \mid N^k(x_v) \cap I_1 \neq \emptyset \}$$

is the set of vertices of I_2 that are connected by a 'long' path with u_v and the vertex x_v has a k-neighbor in I_1 and

$$L_2 = I_2 - (S \cup L_1)$$

is the complement of $S \cup L_1$ in I_2 . In addition, let $L = L_1 \cup L_2$. We construct D following the procedure given below.

Step 0: Set $\mathcal{I} := I_2$, $\mathcal{S} := S$ and $\mathcal{L} := L$.

Step 1: We consider the vertices in S.

Step 1.1: If there exists a vertex $v \in \mathcal{S}$ such that $d(v, w) \leq k$ for a vertex $w \in \mathcal{L}$, we can connect the vertices v, u_v, w and u_w to one component by adding at most 2(k-1) vertices to $I \cup B$.

Set $\mathcal{I} := \mathcal{I} - \{v, w\}, \, \mathcal{S} := \mathcal{S} - \{v\} \text{ and } \mathcal{L} := \mathcal{L} - \{w\} \text{ and repeat Step 1.1.}$

Step 1.2: If there exists a vertex $v \in \mathcal{S}$ such that $d(v, w) \leq k$ for a vertex $w \in \mathcal{S}$ with $v \neq w$, we can connect the vertices v, u_v , w and u_w to one component by adding at most k-1 vertices to $I \cup B$.

Set $\mathcal{I} := \mathcal{I} - \{v, w\}$ and $\mathcal{S} := \mathcal{S} - \{v, w\}$ and repeat Step 1.2.

Step 1.3: If there exists a vertex $v \in \mathcal{S}$ such that $d(v, w) \leq k$ for a vertex $w \in I_2 - (\mathcal{S} \cup \mathcal{L})$, we can connect the vertices v and u_v to w by adding at most k-1 vertices to $I \cup B$.

Set $\mathcal{I} := \mathcal{I} - \{v\}$ and $\mathcal{S} := \mathcal{S} - \{v\}$ and repeat Step 1.3.

Note that after completing Step 1 the set S is empty and there are at most $|I_1| + 2|I_2| - 3(r_1 + r_2) - 2r_3$ components left, where r_i denotes the number of times Step 1.i was repeated for i = 1, 2, 3. Furthermore, we have added at most $(k-1)(2r_1 + r_2 + r_3)$ vertices to $I \cup B$.

Step 2: We consider the vertices in L_1 .

If there exists a vertex $v \in L_1 \cap \mathcal{L}$, let $w \in I_1$ be a k-neighbor of x_v . We can connect the vertices v, u_v and w to one component by adding at most 2(k-1) vertices to $I \cup B$.

Set $\mathcal{I} := \mathcal{I} - \{v\}$ and $\mathcal{L} := \mathcal{L} - \{v\}$ and repeat Step 2.

Note that after completing Step 2 we have $\mathcal{L} \subseteq L_2$ and there are at most $|I_1| + 2|I_2| - 3(r_1 + r_2) - 2r_3 - 2s$ components left, where s denotes the number of times Step 2 was repeated and the numbers r_i are defined as above. Furthermore, we have added at most $(k-1)(2r_1 + r_2 + r_3 + 2s)$ vertices to $I \cup B$.

Step 3: We consider the vertices in L_2 . Recall that for each vertex $v \in L_2$ the vertex x_v has a k-neighbor $w \in I_2$ besides v.

Let v be a vertex in $L_2 \cap \mathcal{L}$ such that x_v has a k-neighbor $w \in I_2 - \mathcal{I}$. We can connect the vertices v, u_v and w by adding at most 2(k-1) vertices to $I \cup B$. Set $\mathcal{I} := \mathcal{I} - \{v\}$ and $\mathcal{L} := \mathcal{L} - \{v\}$ and repeat Step 3.

Note that after completing Step 3 the sets \mathcal{I} and \mathcal{L} are empty and there are at most $|I_1| + 2|I_2| - 3(r_1 + r_2) - 2r_3 - 2s - 2t$ components left, where t denotes the number of times Step 3 was repeated and the numbers r_i and s are defined as above. Furthermore, we have added at most $(k-1)(2r_1 + r_2 + r_3 + 2s + 2t)$ vertices to $I \cup B$.

Step 4: We connect the remaining components to one component.

Let D' be the set of vertices that consists of $I \cup B$ and all vertices added in Steps 1 to 3. Since D' is a k-dominating set of G, the remaining at most $|I_1| + 2|I_2| - 3(r_1 + r_2) - 2r_3 - 2s - 2t$ components can be connected to one component by adding at most $(|I_1| + 2|I_2| - 3(r_1 + r_2) - 2r_3 - 2s - 2t - 1)2k$ vertices to D'.

After completing Step 4 we have constructed a connected k-dominating set $D \supseteq I \cup B$ by adding at most

vertices to $I \cup B$.

We shall show now that the number of vertices we have have added is less or equal than $(2k-1)|I_2| + 2k(|I_1|-1)$. Note that $|I_2| = 2r_1 + 2r_2 + r_3 + s + t$. Then

$$(k-1)(2r_1+r_2+r_3+2s+2t) + (|I_1|+2|I_2|-3(r_1+r_2)-2r_3-2s-2t-1)2k$$

$$-(2k-1)|I_2|-2k(|I_1|-1)$$

$$=(2k+1)|I_2|-3k(2r_1+2r_2+r_3+s+t)-k(r_3+s+t)$$

$$+(k-1)(2r_1+r_2+r_3+2s+2t)$$

$$=-(k-1)(2r_1+2r_2+r_3+s+t)-k(r_3+s+t)+(k-1)(2r_1+r_2+r_3+2s+2t)$$

$$=-(k-1)r_2-kr_3-s-t$$

$$<0.$$

If we choose |I| such that $|I| = ir_k(T)$, it follows that

$$\gamma_k^c(T) \le |D| \le |I_1| + 2|I_2| + 2k|I_1| + (2k-1)|I_2| - 2k$$

= $(2k+1)|I| - 2k$
= $(2k+1)ir_k(T) - 2k$

which completes the proof of this theorem.

As an immediate consequence we get the following corollary.

Corollary 2.8. If T is a tree with n_1 leaves, then

$$ir_k(G) \ge \frac{|V(T)| - kn_1 + 2k}{2k + 1}.$$

Proof. Since $\gamma_k^c(T) \geq |V(T)| - kn_1$ for each tree T, the result follows directly from Theorem 2.7.

Note that, since $\gamma_k(G) \geq ir_k(G)$ for each graph G, Corollary 2.8 is also a generalization of Corollary 2.3. The following theorem provides a class of examples that shows that the bound presented in Theorem 2.7 is sharp.

Theorem 2.9 (Meierling & Volkmann [9] 2005; Cyman, Lemanska & Raczek [2] 2006). Let \mathcal{R} denote the family of trees in which the distance between each pair of distinct leaves is congruent $2k \mod (2k+1)$. If T is a tree with n_1 leaves, then

$$\gamma_k(T) = \frac{|V(T)| - kn_1 + 2k}{2k + 1}$$

if and only if T belongs to the family \mathcal{R} .

Remark 2.10. The graph in Figure 1 shows that the construction presented in the proof of Theorem 2.7 does not work if we allow the graph to contain cycles. It is easy to see that $I = \{v_1, v_2\}$ is an ir_2 -set of G and that $D = \{u_1, u_2, x_1, x_2, x_3\}$ is a γ_2^c -set of G. Following the construction in the proof of Theorem 2.7, we have $I_1 = \emptyset$, $I_2 = \{v_1, v_2\}$ and $B = \{u_1, u_2\}$ and consequently, $D' = I_2 \cup B \cup \{x_1, x_2, x_3\}$. But $|D'| = 7 \not\leq 6 = (2 \cdot 2 + 1)|I| - 2 \cdot 2$ and D contains none of the vertices of I.

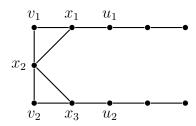


Figure 1.

Nevertheless, we think that the following conjecture is valid.

Conjecture 2.11. If G is a connected graph, then

$$\gamma_k^c(G) \le (2k+1)ir_k(G) - 2k.$$

Now we analyze the relation between the connected distance domination number and the total distance domination number of a graph.

Theorem 2.12. If G is a connected graph, then

$$\gamma_k^c(G) \le \frac{3k+1}{2} \gamma_k^t(G) - 2k.$$

Proof. Let G be a connected graph and let D be a total k-dominating set of G of size $\gamma_k^t(G)$. Each vertex $x \in D$ is in distance at most k of a vertex $y \in D - \{x\}$. Thus we get a dominating set of G with at most $\lfloor |D|/2 \rfloor$ components by adding at most $\lceil |D|/2 \rceil (k-1)$ vertices to D. As in the proof of Lemma 2.4, the resulting components can be joined to a connected k-dominating set |D'| by adding at most $(\lfloor |D|/2 \rfloor - 1)2k$ vertices. Consequently,

$$\gamma_k^c(G) \leq |D'| \leq |D| + \left\lceil \frac{|D|}{2} \right\rceil (k-1) + \left(\left| \frac{|D|}{2} \right| - 1 \right) 2k \leq \frac{3k+1}{2} |D| - 2k = \frac{3k+1}{2} \gamma_k^t(G) - 2k$$

and the proof is complete.

For distance k = 1 we obtain the following result.

Corollary 2.13 (Favaron & Kratsch [4] 1991). If G is a connected graph, then

$$\gamma^c(G) \le 2\gamma^t(G) - 2.$$

The following example shows that the bound presented in Theorem 2.12 is sharp.

Example 2.14. Let P be the path on n=(3k+1)r vertices with $r \in \mathbb{N}$. Then $\gamma_k^c(P)=n-2k,\ \gamma_k^t(P)=2r$ and thus, $\gamma_k^c(P)=\frac{3k+1}{2}\gamma_k^t(P)-2k$.

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