

The n -card problem, stochastic matrices, and the Extreme Principle

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Abstract

The n -card problem is to determine the minimal intervals $[u, v]$ such that for every $n \times n$ stochastic matrix A there is an $n \times n$ permutation matrix P (depending on A) such that $\text{tr}(PA) \in [u, v]$. This problem is closely related to classical mathematical problems from industry and management, including the linear assignment problem and the travelling salesman problem. The minimal intervals for the n -card problem are known only for $n \leq 4$.

We introduce a new method of analysis for the n -card problem that makes repeated use of the Extreme Principle. We use this method to answer a question posed by Sands (2011), by showing that $[1, 2]$ is a solution to the n -card problem for all $n \geq 2$. We also show that each closed interval of length $\frac{n}{n-1}$ contained in $[0, 2)$ is a solution to the n -card problem for all $n \geq 2$.

Keywords: stochastic matrix; permutation matrix; transversal sum; trace; Extreme Principle; n -card problem

1 Introduction

Let $n \geq 2$ be an integer. An $n \times n$ *stochastic matrix* is an $n \times n$ matrix (a_{ij}) of non-negative real numbers, each of whose row sums is 1. A *transversal sum* of (a_{ij}) is a sum of the form $\sum_{i=1}^n a_{\sigma(i), i}$, for some permutation σ of $\{1, 2, \dots, n\}$. A *solution to the n -card problem* is an interval $[u, v]$ such that every $n \times n$ stochastic matrix contains at least one transversal sum in $[u, v]$. Equivalently, a solution to the n -card problem is an interval $[u, v]$ such that

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for every $n \times n$ stochastic matrix A there is an $n \times n$ permutation matrix P such that $\text{tr}(PA) \in [u, v]$. We wish to determine the minimal solutions to the n -card problem for each n , namely those solutions $[u, v]$ for which no proper subinterval of $[u, v]$ is a solution.

The n -card problem is closely related to well-known mathematical problems with industrial and management applications, involving the possible values of $\text{tr}(PA)$ for a permutation matrix P and a fixed square matrix A . In particular, the linear assignment problem, “one of the most famous problems in linear programming and in combinatorial optimization [BDM09],” is to minimize $\text{tr}(PA)$ over all permutation matrices P (see [BDM09, Chapter 4] for a detailed historical account of the development of algorithms for its solution, and an equivalent formulation in terms of weighted bipartite matchings); the travelling salesman problem is the special case in which the permutation corresponding to P is cyclic [Flo56].

The terminology of the n -card problem reflects its original formulation [San01], [LS05] involving a set of n cards, each containing n non-negative real numbers written in a row and summing to 1, with the transversal sum representing the diagonal sum formed when the cards are placed one below the other according to some permutation. In 2001, Sands [San01] asked for a proof that $[\frac{1}{2}, \frac{3}{2}]$ is a solution to the 3-card problem. Lenza and Sands [LS05] introduced the generalization to the n -card problem in 2005.

The interval $[0, 1]$ is a minimal solution to the n -card problem for all $n \geq 2$ [San11, Lemma 3], but it is not the only minimal solution. Indeed, it is easily checked by hand that the minimal solutions to the 2-card problem are

$$[0, 1], \quad [1, 2]. \quad (1)$$

The minimal solutions to the 3-card problem are [LS05]

$$[0, 1], \quad [\frac{1}{2}, \frac{3}{2}], \quad [1, 2], \quad (2)$$

and the minimal solutions to the 4-card problem are [LS05], [San11]

$$[0, 1], \quad [\frac{1}{3}, \frac{4}{3}], \quad [\frac{2}{3}, \frac{5}{3}], \quad [1, 2]. \quad (3)$$

The method of [LS05] and [San11] is to establish a particular interval $[u, v]$ as a solution to the n -card problem (for $n = 3$ or 4), by using intersecting permutations to show that the number of transversal sums greater than v , plus the number of transversal sums less than u , is always less than $n!$. The method has two drawbacks: it relies on a laborious case analysis for $n = 4$, and does not extend to $n \geq 5$ [LS05, p.6].

In this paper we introduce a new method for analysing the n -card problem that makes repeated use of the Extreme Principle [Zei07]. We believe that this method could shed light on other problems involving $\text{tr}(PA)$, where P is a permutation matrix and A is a fixed square matrix. The Extreme Principle directs attention to the “largest” and “smallest” elements of a problem. In the present context, we assume for a contradiction that no transversal sum of an $n \times n$ stochastic matrix lies in some interval $[u, v]$, and then consider the smallest transversal sum d greater than v . Then, if a transversal sum is less than d , it must be less than u . We seek such transversal sums, involving exactly $n - 2$ of

the original summands of d , in order to reach a contradiction. We thereby obtain strong new restrictions for all $n \geq 5$. In particular, we solve Problem 5 of [San11] as follows.

Theorem 1. *For all $n \geq 2$, the interval $[1, 2]$ is a solution to the n -card problem.*

We also prove the following result.

Theorem 2. *For all $n \geq 2$, each closed interval of length $\frac{n}{n-1}$ contained in $[0, 2)$ is a solution to the n -card problem.*

The length $\frac{n}{n-1}$ in Theorem 2 is the smallest possible for a general interval, because the $n \times n$ stochastic matrix having one row $(0, \frac{1}{n-1}, \frac{1}{n-1}, \dots, \frac{1}{n-1})$ and $n-1$ rows $(1, 0, 0, \dots, 0)$ has transversal sums of 0 and $\frac{n}{n-1}$ only.

On the other hand, the known complete set of minimal solutions (2) for $n = 3$ and (3) for $n = 4$ shows that the interval length in Theorem 2 can be reduced to 1 for specific intervals. By reference to particular $n \times n$ stochastic matrices, Sands [San11] showed that every solution to the n -card problem for $n \geq 2$ must contain a length 1 interval $[\frac{k}{n-1}, 1 + \frac{k}{n-1}]$ for some $k \in \{0, 1, \dots, n-1\}$. Problem 3 of [San11] asks whether each such length 1 interval is itself a solution to the n -card problem, which would imply that the complete set of minimal solutions to the n -card problem comprises these n intervals. This question remains open for $n > 4$.

2 The interval $[1, 2]$

In this section we prove Theorem 1. We firstly establish some preliminary lemmas.

Lemma 3. *Let (a_{ij}) be an $n \times n$ stochastic matrix, all of whose transversal sums lie outside an interval $[u, v]$ containing 1. Then (a_{ij}) has at least one transversal sum less than u , and at least one transversal sum greater than v .*

Proof. Each entry of (a_{ij}) is contained in exactly $(n-1)!$ transversal sums, so the mean of all transversal sums is $(n-1)!(\sum_{i,j} a_{ij})/n! = 1$. Therefore at least one transversal sum is at most 1, and so by assumption less than u . Similarly, at least one transversal sum is greater than v . $\square$

An immediate consequence of Lemma 3 is that, as noted earlier, $[0, 1]$ is a solution to the n -card problem for all $n \geq 2$.

The rows and columns of an $n \times n$ stochastic matrix can be permuted without changing the set of its $n!$ transversal sums. Our method relies on examining the effect of transposing two rows of an $n \times n$ stochastic matrix, and thereby bounding the matrix entries. We now show that if an $n \times n$ stochastic matrix has a sufficiently large diagonal sum then there must be a transposition of two rows that decreases this diagonal sum. We prove this result for the following slightly more general case of an $n \times n$ substochastic matrix (each of whose row sums is at most 1).

Lemma 4. Let (a_{ij}) be an $n \times n$ substochastic matrix. Suppose (a_{ij}) has diagonal sum greater than 1. Then

$$a_{kk} + a_{\ell\ell} > a_{k\ell} + a_{\ell k} \quad \text{for some } k, \ell.$$

Proof. Suppose, for a contradiction, that $a_{ii} + a_{jj} \leq a_{ij} + a_{ji}$ for all i, j . Sum this inequality over all i, j to obtain $2n \sum_i a_{ii} \leq 2 \sum_{i,j} a_{ij} \leq 2n$, since by assumption the row sums of (a_{ij}) are each at most 1. This implies that the diagonal sum satisfies $\sum_i a_{ii} \leq 1$, giving the required contradiction. $\square$

We next give conditions under which the sum of two diagonal entries of an $n \times n$ stochastic matrix can be bounded from below.

Lemma 5. Let (a_{ij}) be an $n \times n$ stochastic matrix with diagonal sum d , and suppose all transversal sums of (a_{ij}) lie outside the interval $[u, d)$. Then, for all i, j ,

$$a_{ii} + a_{jj} > a_{ij} + a_{ji} \quad \text{implies} \quad a_{ii} + a_{jj} > d - u.$$

Proof. Suppose $a_{ii} + a_{jj} > a_{ij} + a_{ji}$. Then the positive quantity $a_{ii} + a_{jj} - a_{ij} - a_{ji}$ is the decrease in the diagonal sum caused by transposing rows i and j of the matrix, and so by assumption is greater than $d - u$. We therefore have $a_{ii} + a_{jj} \geq a_{ii} + a_{jj} - a_{ij} - a_{ji} > d - u$. $\square$

Define an $n \times n$ stochastic matrix (a_{ij}) to be *diagonally ordered* if its diagonal entries are in non-increasing order:

$$a_{11} \geq a_{22} \geq \dots \geq a_{nn}.$$

We are now ready to prove Theorem 1.

Proof of Theorem 1. We know from (1) and (2) that the result holds for $n = 2$ and 3, so we may take $n \geq 4$. Suppose, for a contradiction, that (a_{ij}) is an $n \times n$ stochastic matrix whose transversal sums all lie outside the interval $[1, 2]$. Then by Lemma 3, (a_{ij}) has a transversal sum greater than 2 and a transversal sum less than 1. Let $2 + \epsilon$ be the smallest transversal sum greater than 2, and reorder the rows and columns of (a_{ij}) so that the summands of this transversal sum occur on the matrix diagonal and so that the matrix is diagonally ordered. By Lemma 5 with $d = 2 + \epsilon$ and $u = 1$,

$$a_{ii} + a_{jj} > a_{ij} + a_{ji} \quad \text{implies} \quad a_{ii} + a_{jj} > 1 + \epsilon. \quad (4)$$

Now the $(n - 1) \times (n - 1)$ submatrix of (a_{ij}) formed by deleting the first row and column has diagonal sum $2 + \epsilon - a_{11} > 1$. Apply Lemma 4 to this submatrix to show that, for some distinct $k > 1$ and $\ell > 1$,

$$a_{kk} + a_{\ell\ell} > a_{k\ell} + a_{\ell k}.$$

We therefore have $a_{kk} + a_{\ell\ell} > 1 + \epsilon$ by (4), and so

$$a_{22} + a_{33} > 1 + \epsilon \quad (5)$$

since the matrix is diagonally ordered and k, ℓ are distinct. Since the diagonal sum of (a_{ij}) is $2 + \epsilon$, we have

$$a_{ii} + a_{11} \leq 2 + \epsilon - a_{22} - a_{33} \quad \text{for all } i > 3,$$

and therefore $a_{ii} + a_{11} < 1$ for all $i > 3$, by (5). Then, since the matrix is diagonally ordered,

$$a_{ii} + a_{jj} < 1 \quad \text{for all } i, j \text{ with } i > 3,$$

which in turn implies by (4) that

$$a_{ii} + a_{jj} \leq a_{ij} + a_{ji} \quad \text{for all } i, j \text{ with } i > 3. \quad (6)$$

We complete the proof by showing that (5) and (6) force the sum of the entries of (a_{ij}) to be too large. We have

$$\begin{aligned} \sum_{i,j} a_{ij} &\geq \sum_{i \leq 3} a_{ii} + \sum_{i > 3} a_{ii} + \sum_{i > 3, j \leq 3} (a_{ij} + a_{ji}) \\ &\geq (n-2) \sum_{i \leq 3} a_{ii} + 4 \sum_{i > 3} a_{ii} \end{aligned}$$

by substitution from (6). Therefore

$$\begin{aligned} \sum_{i,j} a_{ij} &\geq (n-2) \sum_{i \leq 3} a_{ii} + 2 \sum_{i > 3} a_{ii} \\ &= (n-4) \sum_{i \leq 3} a_{ii} + 2 \sum_i a_{ii} \\ &\geq (n-4)(1+\epsilon) + 2(2+\epsilon) \end{aligned}$$

by (5), using $n \geq 4$. Therefore $\sum_{i,j} a_{ij} > n$, which is a contradiction because each row sum of (a_{ij}) is 1. $\square$

3 Intervals of length $\frac{n}{n-1}$

In this section we prove Theorem 2.

Proposition 6. *Let $n \geq 4$ and let (a_{ij}) be a diagonally ordered $n \times n$ stochastic matrix. Suppose the diagonal sum d of (a_{ij}) satisfies $d \in (1, 2]$. Then (a_{ij}) has a transversal sum lying in the interval $[d - \frac{n}{n-1}, d)$.*

Proof. Suppose, for a contradiction, that no transversal sum of (a_{ij}) lies in the interval $[d - \frac{n}{n-1}, d)$. Then by Lemma 5 with $u = d - \frac{n}{n-1}$,

$$a_{ii} + a_{jj} > a_{ij} + a_{ji} \quad \text{implies} \quad a_{ii} + a_{jj} > \frac{n}{n-1}. \quad (7)$$

Since $d > 1$, Lemma 4 gives

$$a_{kk} + a_{\ell\ell} > a_{k\ell} + a_{\ell k} \quad \text{for some distinct } k, \ell,$$

and it follows from (7) that $a_{kk} + a_{\ell\ell} > \frac{n}{n-1}$. Since the matrix is diagonally ordered and k, ℓ are distinct, this implies

$$a_{11} + a_{22} > \frac{n}{n-1} \quad (8)$$

and

$$a_{11} > \frac{1}{2} \cdot \frac{n}{n-1}. \quad (9)$$

We now claim that

$$a_{ii} + a_{jj} \leq \frac{n}{n-1} \quad \text{for all distinct } i > 1 \text{ and } j > 1. \quad (10)$$

Suppose otherwise, for a contradiction, so that $a_{rr} + a_{ss} > \frac{n}{n-1}$ for some distinct $r > 1$ and $s > 1$. Since the matrix is diagonally ordered, this gives

$$a_{22} + a_{33} > \frac{n}{n-1}. \quad (11)$$

Therefore, for all $i > 3$, we have $a_{ii} + a_{11} \leq d - a_{22} - a_{33} < \frac{n-2}{n-1}$ because $d \leq 2$. Since the matrix is diagonally ordered, we then have

$$a_{ii} + a_{jj} < \frac{n-2}{n-1} \quad \text{for all } i, j \text{ with } i > 3,$$

which by (7) implies

$$a_{ii} + a_{jj} \leq a_{ij} + a_{ji} \quad \text{for all } i, j \text{ with } i > 3. \quad (12)$$

Write

$$\begin{aligned} \sum_{i,j} a_{ij} &\geq \sum_{i \leq 3} a_{ii} + \sum_{i > 3, j \leq 3} (a_{ij} + a_{ji}) \\ &\geq (n-2) \sum_{i \leq 3} a_{ii} + 3 \sum_{i > 3} a_{ii} \end{aligned}$$

by substitution from (12), so that

$$\begin{aligned} \sum_{i,j} a_{ij} &\geq (n-2) \sum_{i \leq 3} a_{ii} \\ &> (n-2) \cdot \frac{3}{2} \cdot \frac{n}{n-1} \end{aligned}$$

from (9) and (11). Since $\sum_{i,j} a_{ij} = n$ and $n \geq 4$, this is a contradiction and proves the claim (10).

It then follows from (7) that

$$a_{ii} + a_{jj} \leq a_{ij} + a_{ji} \quad \text{for all distinct } i > 1 \text{ and } j > 1.$$

Summing over all i, j satisfying $1 < i < j$, we find that

$$(n-2) \sum_{i>1} a_{ii} \leq \sum_{1<i<j} (a_{ij} + a_{ji}). \quad (13)$$

Now let m be the largest integer i such that $a_{11} + a_{ii} > \frac{n}{n-1}$. Note that $m \geq 2$, by (8). By (7) we have $a_{11} + a_{ii} \leq a_{1i} + a_{i1}$ for $i > m$, so that

$$a_{11} \leq a_{1i} + a_{i1} \quad \text{for } i > m. \quad (14)$$

We now show that (13) and (14) force the entries of (a_{ij}) to be too large. We have

$$\begin{aligned} \sum_{i,j} a_{ij} &\geq a_{11} + \sum_{i>1} a_{ii} + \sum_{i>m} (a_{1i} + a_{i1}) + \sum_{1<i<j} (a_{ij} + a_{ji}) \\ &\geq (n-m+1)a_{11} + (n-1) \sum_{i>1} a_{ii} \end{aligned}$$

by substitution from (13) and (14). Therefore

$$\begin{aligned} \sum_{i,j} a_{ij} &\geq (n-m+1)a_{11} + (n-1) \sum_{1<i\leq m} a_{ii} \\ &> (n-m+1)a_{11} + (n-1)(m-1) \left(\frac{n}{n-1} - a_{11} \right) \end{aligned}$$

by definition of m and the diagonal ordering of (a_{ij}) , and so

$$\sum_{i,j} a_{ij} > n(m-1 - (m-2)a_{11}).$$

Since $m \geq 2$ and $a_{11} \leq 1$, this implies the contradiction $\sum_{i,j} a_{ij} > n$ and so completes the proof. $\square$

We now combine Proposition 6 with Theorem 1 to prove Theorem 2.

Proof of Theorem 2. We know from (1) and (2) that the result holds for $n = 2$ and 3 , so we may take $n \geq 4$. Suppose, for a contradiction, that (a_{ij}) is an $n \times n$ stochastic matrix whose transversal sums all lie outside the interval $[u, u + \frac{n}{n-1}]$ for some $u \in [0, \frac{n-2}{n-1}]$. Since this interval contains 1 , by Theorem 1 the matrix (a_{ij}) therefore has a transversal sum in the interval $(u + \frac{n}{n-1}, 2]$. Let d be the smallest such transversal sum. Reorder the rows and columns of (a_{ij}) so that the summands of this transversal sum occur on the matrix diagonal and so that the matrix is diagonally ordered. Then by Proposition 6, (a_{ij}) has a transversal sum lying in the interval $[d - \frac{n}{n-1}, d)$. By choice of d , this gives the required contradiction. $\square$

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