Γ -species and the enumeration of k-trees

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Submitted: Oct 10, 2012; Accepted: Dec 4, 2012; Published: Dec 13, 2012 Mathematics Subject Classifications: 05C30, 05E18

Abstract

We study the class of graphs known as k-trees through the lens of Joyal's theory of combinatorial species (and an extension known as ' Γ -species' which incorporates data about 'structural' group actions). This culminates in a system of recursive functional equations giving the generating function for unlabeled k-trees which allows for fast, efficient computation of their numbers. Enumerations up to k = 10 and n = 30 (for a k-tree with n + k - 1 vertices) are included in tables, and Sage code for the general computation is included in an appendix.

1 Introduction

1.1 k-trees

Trees and their generalizations have played an important role in the literature of combinatorial graph theory throughout its history. The multi-dimensional generalization to so-called 'k-trees' has proved to be particularly fertile ground for both research problems and applications.

The class \mathfrak{a}_k of k-trees (for $k \in \mathbb{N}$) may be defined recursively:

Definition 1.1. The complete graph on k vertices (K_k) is a k-tree, and any graph formed by adding a single vertex to a k-tree and connecting that vertex by edges to some existing k-clique (that is, induced k-complete subgraph) of that k-tree is a k-tree.

The graph-theoretic notion of k-trees was first introduced in 1968 in [7]; vertex-labeled k-trees were quickly enumerated in the following year in both [10] and [2]. The special case k = 2 has been especially thoroughly studied; enumerations are available in the literature for edge- and triangle-labeled 2-trees in [11], for plane 2-trees in [12], and for unlabeled 2-trees in [7] and [6]. In 2001, the theory of species was brought to bear on 2-trees in [4],

resulting in more explicit formulas for the enumeration of unlabeled 2-trees. An extensive literature on other properties of k-trees and their applications has also emerged; Beineke and Pippert claim in [1] that "[t]here are now over 100 papers on various aspects of k-trees". However, no general enumeration of unlabeled k-trees appears in the literature to date.

Although we do not derive a closed form for the number of k-trees, the work in this paper does permit efficient recursive computation of their generating function. A formula for this generating function is given in Corollary 6.5 using components defined recursively in Corollary 6.2.

To begin, we establish two definitions for substructures of k-trees which we will use extensively in our analysis.

Definition 1.2. A hedron of a k-tree is a (k + 1)-clique and a front is a k-clique.

We will frequently describe k-trees as assemblages of hedra attached along their fronts rather than using explicit graph-theoretic descriptions in terms of edges and vertices, keeping in mind that the structure of interest is graph-theoretic and not geometric. The recursive addition of a single vertex and its connection by edges to an existing k-clique in Definition 1.1 is then interpreted as the attachment of a hedron to an existing one along some front, identifying the k vertices they have in common. The analogy to the recursive definition of conventional trees is clear, and in fact the class \mathfrak{a} of trees may be recovered by setting k = 1. For higher k, the structures formed are still distinctively tree-like; for example, 2-trees are formed by gluing triangles together along their edges without forming loops of triangles (see Fig. 1), while 3-trees are formed by gluing tetrahedra together along their triangular faces without forming loops of tetrahedra.



Figure 1: A (vertex-labeled) 2-tree

In graph-theoretic contexts, it is conventional to label graphs on their vertices and possibly their edges. However, for our purposes, it will be more convenient to label hedra and fronts. Throughout, we will treat the species \mathfrak{a}_k of k-trees as a two-sort species, with X-labels on the hedra and Y-labels on their fronts; in diagrams, we will generally use capital letters for the hedron-labels and positive integers for the front-labels (see Fig. 2). A formula for the cycle index of the species \mathfrak{a}_k is given in Theorem 5.10 using components defined recursively in Theorems 5.5 and 5.6. (Readers unfamiliar with the theory of species

and its applications to graph enumeration may find a full exposition of the subject in [3], which also will serve as a reference for any unexplained notation.)



Figure 2: A (hedron-and-front-labeled) 2-tree

2 The dissymmetry theorem for k-trees

Studies of tree-like structures—especially those explicitly informed by the theory of species, as ours will be—often feature decompositions based on *dissymmetry*, which allow enumerations of unrooted structures to be recharacterized in terms of rooted structures. For example, as seen in [3, §4.1], the species \mathfrak{a} of trees and $\mathcal{A} = \mathfrak{a}^{\bullet}$ of rooted trees are related by the equation

$$\mathcal{A} + \mathcal{E}_2(\mathcal{A}) = \mathfrak{a} + \mathcal{A}^2$$

where the proof hinges on a recursive structural decomposition of trees. In this case, the species \mathcal{A} is relatively easy to characterize explicitly, so this equation serves to characterize the species \mathfrak{a} , which would be difficult to do directly.

A similar theorem holds for k-trees.

Theorem 2.1. The species \mathfrak{a}_k^X and \mathfrak{a}_k^Y of k-trees rooted at hedra and fronts respectively, \mathfrak{a}_k^{XY} of k-trees rooted at a hedron with a designated front, and \mathfrak{a}_k of unrooted k-trees are related by the equation

$$\mathfrak{a}_k^X + \mathfrak{a}_k^Y = \mathfrak{a}_k + \mathfrak{a}_k^{XY} \tag{1}$$

as an isomorphism of species.

Proof. We give a bijective, natural (*i.e.* label-equivariant) map from $(\mathfrak{a}_k^X + \mathfrak{a}_k^Y)$ -structures on the left side to $(\mathfrak{a}_k + \mathfrak{a}_k^{XY})$ -structures on the right side. Define a *k*-path in a *k*-tree to be a non-self-intersecting sequence of consecutively adjacent hedra and fronts, and define the *length* of a *k*-path to be the total number of hedra and fronts along it. Note that the ends of every maximal k-path in a k-tree are fronts. It is easily verified, as in [9], that every k-tree has a unique *center* clique (either a hedron or a front) which is the midpoint of every longest k-path (or, equivalently, has the greatest k-eccentricity, defined appropriately).

An $(\mathfrak{a}_k^X + \mathfrak{a}_k^Y)$ -structure on the left-hand side of the equation is a k-tree T rooted at some clique c, which is either a hedron or a front. Suppose that c is the center of T. We then map T to its unrooted equivalent in \mathfrak{a}_k on the right-hand side. This map is a natural bijection from its preimage, the set of k-trees rooted at their centers, to \mathfrak{a}_k , the set of unrooted k-trees.

Now suppose that the root clique c of the k-tree T is not the center, which we denote C. Identify the clique c' which is adjacent to c along the k-path from c to C. We then map the k-tree T rooted at the clique c to the same tree T rooted at both c and its neighbor c'. This map is also a natural bijection, in this case from the set of k-trees rooted at vertices which are not their centers to the set \mathfrak{a}_k^{XY} of k-trees rooted at an adjacent hedron-front pair.

Since these maps are label-equivariant bijections, they induce an isomorphism of species

$$\mathfrak{a}_k^X + \mathfrak{a}_k^Y = \mathfrak{a}_k + \mathfrak{a}_k^{XY}$$

as desired, completing the proof.

In general we will reformulate the dissymmetry theorem as follows:

Corollary 2.2. For the various forms of the species \mathfrak{a}_k as above, we have

$$\mathbf{a}_k = \mathbf{a}_k^X + \mathbf{a}_k^Y - \mathbf{a}_k^{XY}.$$
 (2)

as an isomorphism of species.

This species subtraction is well-defined in the sense that since the species \mathfrak{a}_k^{XY} embeds in the species $\mathfrak{a}_k^X + \mathfrak{a}_k^Y$ by the centering map described in the proof of Theorem 2.1. Essentially, Eq. (2) identifies each unrooted k-tree with itself rooted at its center simplex. This may be understood formally either in the sense of virtual species as in [3, §2.5] or in the sense of species maps as in [5, Def. 1.3.3].

Theorem 2.1 and the consequent Eq. (2) allow us to reframe enumerative questions about generic k-trees in terms of questions about k-trees rooted in various ways. However, the rich internal symmetries of large cliques obstruct direct analysis of these rooted structures. We need to break these symmetries to proceed.

3 Coherently-oriented *k*-trees

3.1 Symmetry-breaking

In the case of the species $\mathcal{A} = \mathfrak{a}_1^{\bullet}$ of rooted trees, we may obtain a simple recursive functional equation [3, §1, eq. (9)]:

$$\mathcal{A} = X \cdot \mathcal{E}(\mathcal{A}). \tag{3}$$

This completely characterizes the combinatorial structure of the class of trees.

However, in the more general case of k-trees, no such simple relationship obtains; attached to a given hedron is a collection of sets of hedra (one such set per front), but simply specifying which fronts to attach to which does not fully specify the attachings, and the structure of that collection of sets is complex. We will break this symmetry by adding additional structure which we can later remove using the theory of quotient species.

Definition 3.1. Let h_1 and h_2 be two hedra joined at a front f, hereafter said to be *adjacent*. Each other front of one of the hedra shares k - 1 vertices with f; we say that two fronts f_1 of h_1 and f_2 of h_2 are *mirror with respect to* f if these shared vertices are the same, or equivalently if $f_1 \cap f = f_2 \cap f$.

Observation 3.2. Let T be a k-tree with two hedra h_1 and h_2 joined at a front f. Then there is exactly one front of h_2 mirror to each front of h_1 with respect to their shared front f.

Definition 3.3. Define an *orientation* of a hedron to be a cyclic ordering of the set of its fronts and an *orientation* of a k-tree to be a choice of orientation for each of its hedra. If two oriented hedra share a front, their orientations are *compatible* if they correspond under the mirror bijection. Then an orientation of a k-tree is *coherent* if every pair of adjacent hedra is compatibly-oriented.

See Fig. 3 for an example. Note that every k-tree admits many coherent orientations any one hedron of the k-tree may be oriented freely, and a unique orientation of the whole k-tree will result from each choice of such an orientation of one hedron. We will denote by \vec{a}_k the species of coherently-oriented k-trees.

By shifting from the general k-tree setting to that of coherently-oriented k-trees, we break the symmetry described above. If we can now establish a group action on \vec{a}_k whose orbits are generic k-trees we can use the theory of quotient species to extract the generic species a_k . First, however, we describe an encoding procedure which will make future work more convenient.

3.2 Bicolored tree encoding

Although k-trees are graphs (and hence made up simply of edges and vertices), their structure is more conveniently described in terms of their simplicial structure of hedra and fronts. Indeed, if each hedron has an orientation of its faces and we choose in advance which hedra to attach to which by what fronts, the requirement that the resulting k-tree be coherently oriented is strong enough to characterize the attaching completely. We thus pass from coherently-oriented k-trees to a surrogate structure which exposes the salient features of this attaching structure more clearly—structured bicolored trees in the spirit of the R, S-enriched bicolored trees of [3, §3.2].

A $(\mathcal{C}_{k+1}, \mathcal{E})$ -enriched bicolored tree is a bicolored tree each black vertex of which carries a \mathcal{C}_{k+1} -structure (that is, a cyclic ordering on k+1 elements) on its white neighbors. (The \mathcal{E} -structure on the black neighbors of each white vertex is already implicit in the bicolored



Figure 3: A coherently-oriented 2-tree

tree itself.) For later convenience, we will sometimes call such objects k-coding trees, and we will denote by \mathfrak{CT}_k the species of such k-coding trees.

We now define a map $\beta : \vec{\mathfrak{a}}_k[n] \to \mathfrak{CT}_k[n]$. For a given coherently-oriented k-tree T with n hedra:

- For every hedron of T construct a black vertex and for every front a white vertex, assigning labels appropriately.
- For every black-white vertex pair, construct a connecting edge if the white vertex represents a front of the hedron represented by the black vertex.
- Finally, enrich the collection of neighbors of each black vertex with a C_{k+1} -structure inherited directly from the orientation of the k-tree T.

The resulting object $\beta(T)$ is clearly a k-coding tree with n black vertices.

We can recover a T from $\beta(T)$ by following the reverse procedure. For an example, see Fig. 4, which shows the 2-coding tree associated to the coherently-oriented 2-tree of Fig. 3. Note that, for clarity, we have rendered the black vertices (corresponding to hedra) with squares.

Theorem 3.4. The map β induces an isomorphism of species $\vec{\mathfrak{a}}_k \simeq \mathfrak{CT}_k$.

Proof. It is clear that β sends each coherently-oriented k-tree to a unique k-coding tree, and that this map commutes with permutations on the label sets (and thus is categorically natural). To show that β induces a species isomorphism, then, we need only show that β is a surjection onto $\mathfrak{CT}_k[n]$ for each n. Throughout, we will say 'F and G have contact of order n' when the restrictions $F_{\leq n}$ and $G_{\leq n}$ of the species F and G to label sets of cardinality at most n are naturally isomorphic.

First, we note that there are exactly k! coherently-oriented k-trees with one hedron one for each cyclic ordering of the k + 1 front labels. There are also k! coding trees with



Figure 4: A $(\mathcal{C}_{k+1}, \mathcal{E})$ -enriched bicolored tree encoding a coherently-oriented 2-tree

one black vertex, and the encoding β is clearly a natural bijection between these two sets. Thus, the species $\vec{\mathfrak{a}}_k$ of coherently-oriented k-trees and \mathfrak{CT}_k of k-coding trees have contact of order 1.

Now, by way of induction, suppose $\vec{\mathfrak{a}}_k$ and \mathfrak{CT}_k have contact of order $n \ge 1$. Let C be a k-coding tree with n + 1 black vertices. Then let C_1 and C_2 be two distinct sub-k-coding trees of C, each obtained from C by removing one black node which has only one white neighbor which is not a leaf. Then, by hypothesis, there exist coherently-oriented k-trees T_1 and T_2 with n hedra such that $\beta(T_1) = C_1$ and $\beta(T_2) = C_2$. Moreover, $\beta(T_1 \cap T_2) = \beta(T_1) \cap \beta(T_2)$, and this k-coding tree has n - 1 black vertices, so $T_1 \cap T_2$ has n - 1 hedra. Thus, $T = T_1 \cup T_2$ is a coherently-oriented k-tree with n + 1 black hedra, and $\beta(T) = C$ as desired. Thus, $\beta^{-1}(\beta(T_1) \cup \beta(T_2)) = T_1 \cup T_2 = T$, and hence $\vec{\mathfrak{a}}_k$ and \mathfrak{CT}_k have contact of order n + 1.

Thus, $\vec{\mathfrak{a}}_k$ and \mathfrak{CT}_k are isomorphic as species; however, k-coding trees are much simpler than coherently-oriented k-trees as graphs. Moreover, k-coding trees are doubly-enriched bicolored trees as in [3, §3.2], for which the authors of that text develop a system of functional equations which fully characterizes the cycle index of such a species. We thus will proceed in the following sections with a study of the species \mathfrak{CT}_k , then lift our results to the k-tree context.

3.3 Functional decomposition of k-coding trees

With the encoding $\beta : \vec{\mathfrak{a}}_k \to \mathfrak{CT}_k$, we now have direct graph-theoretic access to the attaching structure of coherently-oriented k-trees. We therefore turn our attention to the

k-coding trees themselves to produce a recursive decomposition. As with k-trees, we will study rooted versions of the species \mathfrak{CT}_k of k-coding trees first, then use dissymmetry to apply the results to unrooted enumeration.

Theorem 3.5. The species \mathfrak{CT}_k^X of X-rooted k-coding trees, \mathfrak{CT}_k^Y of Y-rooted k-coding trees, and \mathfrak{CT}_k^{XY} of edge-rooted k-coding trees satisfy the functional equations

$$\mathfrak{CT}_{k}^{X} = X \cdot \mathfrak{C}_{k+1}(\mathfrak{CT}_{k}^{Y}) \tag{4a}$$

$$\mathfrak{CT}_{k}^{Y} = Y \cdot \mathcal{E}\left(X \cdot \mathcal{L}_{k}(\mathfrak{CT}_{k}^{Y})\right)$$
(4b)

$$\mathfrak{CT}_{k}^{XY} = \mathfrak{CT}_{k}^{Y} \cdot X \cdot \mathcal{L}_{k} \big(\mathfrak{CT}_{k}^{Y} \big) = X \cdot \mathcal{L}_{k+1} \big(\mathfrak{CT}_{k}^{Y} \big)$$
(4c)

as isomorphisms of species.

Proof. By construction, a \mathfrak{CT}_k^X -structure consists of a single X-label and a cyclicallyordered (k + 1)-set of \mathfrak{CT}_k^Y -structures. This gives Eq. (4a). See Fig. 5 for an example of this construction.



Figure 5: An example $\mathfrak{CT}_4^X\text{-structure, rooted at the X-vertex.}$

Similarly, a \mathfrak{CT}_k^Y -structure consists of a single Y-label and a (possibly empty) set of structures which are a slight variant of the \mathfrak{CT}_k^X -structures discussed above. Every white neighbor of the black root of a \mathfrak{CT}_k^X -structure is labeled in the construction above, but the white parent of a \mathfrak{CT}_k^X -structure in this recursive decomposition is already labeled. Thus,

the structure around a black vertex which is a child of a white vertex consists of an X label and a linearly-ordered k-set of \mathfrak{CT}_k^Y -structures. Thus, a \mathfrak{CT}_k^Y -structure consists of a Y-label and a set of pairs of an X label and an \mathcal{L}_k -structure of \mathfrak{CT}_k^Y -structures. This gives Eq. (4b). We note here for conceptual consistency that in fact $\mathcal{L}_k = \mathfrak{C}_{k+1}'$ for \mathcal{L} the species of linear orders and \mathfrak{C} the species of cyclic orders and that $\mathcal{E}' = \mathcal{E}$ for \mathcal{E} the species of sets; readers familiar with the R, S-enriched bicolored trees of [3, §3.2] will recognize echoes of their decomposition in these facts.

Finally, a \mathfrak{CT}_k^{XY} -structure is simply an $X \cdot \mathcal{L}_k(\mathfrak{CT}_k^Y)$ -structure as described above (corresponding to the black vertex) together with a \mathfrak{CT}_k^Y -structure (corresponding to the white vertex). For reasons that will become clear later, we note that we can incorporate the root white vertex into the linear order by making it last, thus representing a \mathfrak{CT}_k^{XY} structure instead as an $X \cdot \mathcal{L}_{k+1}(\mathfrak{CT}_k^Y)$ -structure. This gives Eq. (4c). See Fig. 6 for an example of this construction.



Figure 6: An example \mathfrak{CT}_4^{XY} -structure, rooted at the X-vertex and the thick edge adjoining it.

However, a recursive characterization of the various species of k-coding trees is insufficient to characterize the species of k-trees itself, since k-coding trees represent k-trees with coherent orientations.

4 Generic *k*-trees

In [4], the orientation-reversing action of \mathfrak{S}_2 on $\operatorname{Cyc}_{[3]}$ is exploited to study 2-trees speciestheoretically. We might hope to develop an analogous group action under which general *k*-trees are naturally identified with orbits of coherently-oriented *k*-trees under an action of \mathfrak{S}_k . Unfortunately:

Proposition 4.1. For $k \ge 3$, no transitive action of any group on the set $\operatorname{Cyc}_{[k+1]}$ of cyclic orders on [k+1] commutes with the action of \mathfrak{S}_{k+1} that permutes labels.

Proof. We represent the elements of $\operatorname{Cyc}_{[k+1]}$ as cyclic permutations on the alphabet [k+1]; then the action of \mathfrak{S}_{k+1} that permutes labels is exactly the conjugation action on these permutations. Consider an action of a group G on $\operatorname{Cyc}_{[k+1]}$ that commutes with this conjugation action. Then, for any $g \in G$ and any $c \in \operatorname{Cyc}_{[k+1]}$, we have that

$$g \cdot c = g \cdot ccc^{-1} = c(g \cdot c)c^{-1} \tag{5}$$

and so c and $g \cdot c$ commute. Thus, c commutes with every element of its orbit under the action of G. But, for $k \ge 3$, not all elements of $\operatorname{Cyc}_{[k+1]}$ commute, so the action is not transitive.

We thus cannot hope to attack the coherent orientations of k-trees by acting directly on the cyclic orderings of fronts. Instead, we will use the additional structure on *rooted* coherently-oriented k-trees; with rooting, the cyclic orders around black vertices are converted into linear orders, for which there is a natural action of \mathfrak{S}_{k+1} .

4.1 Group actions on k-coding trees

We have noted previously that every labeled k-tree admits exactly k! coherent orientations. Thus, there are k! distinct k-coding trees associated to each labeled k-tree, which differ only in the \mathcal{C}_{k+1} -structures on their black vertices. Consider a rooted k-coding tree T and a black vertex v which is not the root vertex. Then one white neighbor of v is the 'parent' of v (in the sense that it lies on the path from v to the root). We thus can convert the cyclic order on the k + 1 white neighbors of v to a linear order by choosing the parent white neighbor to be last. There is a natural, transitive, label-independent action of \mathfrak{S}_{k+1} on the set of such linear orders which induces an action on the cyclic orders from which the linear orders are derived. However, only elements of \mathfrak{S}_{k+1} which fix k + 1 will respect the structure around the black vertex we have chosen, since its parent white vertex must remain last.

In addition, if we simply apply the action of some $\sigma \in \mathfrak{S}_{k+1}$ to the order on white neighbors of v, we change the coherently-oriented k-tree $\beta^{-1}(T)$ to which T is associated in such a way that it no longer corresponds to the same unoriented k-tree. Let t denote the unoriented k-tree associated to $\beta^{-1}(T)$; then there exists a coherent orientation of t which agrees with orientation around v induced by σ . The k-coding tree T' corresponding to this new coherent orientation has the same underlying bicolored tree as T but possibly different orders around its black vertices. If we think of the k-coding tree T' as the image of T under a global action of σ , orbits under all of \mathfrak{S} will be precisely the classes of k-coding trees corresponding to all coherent orientations of specified k-trees, allowing us to study unoriented k-trees as quotients. The orientation of T' will be that obtained by applying σ at v and then recursively adjusting the other cyclic orders so that fronts which were mirror are made mirror again. This will ensure that the combinatorial structure of the underlying k-tree t is preserved.

Therefore, when we apply some permutation $\sigma \in \mathfrak{S}_{k+1}$ to the white neighbors of a black vertex v, we must also permute the cyclic orders of the descendant black vertices of v. In particular, the permutation σ' which must be applied to some immediate black descendant v' of v is precisely the permutation on the linear order of white neighbors of v' induced by passing over the mirror bijection from v' to v, applying σ , and then passing back. We can express this procedure in formulaic terms:

Theorem 4.2. If a permutation $\sigma \in \mathfrak{S}_{k+1}$ is applied to a linearized orientation of a black vertex v in rooted k-coding tree, the permutation which must be applied to the linearized orientation a child black vertex v' which was attached to the ith white child of v (with respect to the linear ordering induced by the orientation) to preserve the mirror relation is $\rho_i(\sigma)$, where ρ_i is the map given by

$$\rho_i(\sigma): a \mapsto \sigma(i+a) - \sigma(i) \tag{6}$$

in which all sums and differences are reduced to their representatives modulo k + 1 in $\{1, 2, \ldots, k + 1\}$.

Proof. Let v' denote a black vertex which is attached to v by the white vertex 1, which we suppose to be in position i in the linear order induced by the original orientation of v. Let 2 denote the white neighbor of v' which is ath in the linear order induced by the original orientation around v'. It is mirror to the white neighbor 3 of v which is (i + a)th in the linear order induced by the original orientation around v. After the action of σ is applied, vertex 3 is $\sigma(i + a)$ th in the new linear order around v. We require that 2 is still mirror to 3, so we must move it to position $\sigma(i + a) - \sigma(i)$ when we create a new linear order around v'. This completes the proof.

This procedure is depicted in Fig. 7.

As an aside, we note that, although the construction ρ depends on k, the value of k will be fixed in any given context, so we suppress it in the notation.

Any σ which is to be applied to a non-root black vertex v must of course fix k + 1. We can think of \mathfrak{S}_k as the subgroup of \mathfrak{S}_{k+1} of permutations fixing k + 1, and so in what follows we will refer to this as an \mathfrak{S}_k -action where appropriate.

In light of Theorem 3.5, we now wish to adapt these ideas into explicit \mathfrak{S}_{k} - and \mathfrak{S}_{k+1} -actions on \mathfrak{CT}_{k}^{X} , \mathfrak{CT}_{k}^{Y} , and \mathfrak{CT}_{k}^{XY} whose orbits correspond to the various coherent orientations of single underlying rooted k-trees. In the case of a Y-rooted k-coding tree T, if we declare that $\sigma \in \mathfrak{S}_{k}$ acts on T by acting directly on each of the black vertices immediately adjacent to the root and then applying ρ -derived permutations recursively to



Figure 7: Application of a permutation σ to the orientation of a non-root black vertex v. The vertices 2 and 3 are mirror in the original orientation (lower set of edges), as shown by the arrows μ , so we must preserve this mirror relation when we apply σ . The permutation σ moves 3 from the (i + a)th place to the $\sigma(i + a)$ th, so $\rho_i(\sigma)$ must carry 2 from the *a*th place to the $(\sigma(i + a) - \sigma(i))$ th.

their descendants, orbits behave as expected. The same \mathfrak{S}_k -action serves equally well for edge-rooted k-coding trees, where (for purposes of applying the action of some σ) we can simply ignore the black vertex in the root.

However, if we begin with an X-rooted k-coding tree, the cyclic ordering of the white neighbors of the root black vertex has no canonical choice of linearization. If we make an arbitrary choice of one of the k + 1 available linearizations, and thus convert to an edge-rooted k-coding tree, the full \mathfrak{S}_{k+1} -action defined previously can be applied directly to the root vertex. The orbit under this action of some edge-rooted k-coding tree T with a choice of linearization at the root then includes all possible linearizations of the root orders of all possible X-rooted k-coding tree.

It follows that:

Lemma 4.3. The actions of \mathfrak{S}_k on \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} and the action of \mathfrak{S}_{k+1} on \mathfrak{CT}_k^{XY} are transitive in the sense that each orbit corresponds to the set of all coherent orientations of a single underlying rooted unoriented k-tree. Moreover, these actions all commute with the \mathfrak{S}_n -actions which permute labels.

4.2 k-trees as quotients

We have now equipped the various species of rooted k-coding trees with actions which commute with permutations of labels, making them 'species-compatible'. Moreover, the non-oriented rooted k-trees underlying the k-coding trees are naturally identified with orbits under these actions, suggesting that rooted k-trees are 'quotients' of rooted kcoding trees. The notion of a Γ -species as developed in [8, §3] formalizes this notion of compatibility and provides an enumerative toolset for dealing with quotients of this sort. In this language, we will treat \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} as \mathfrak{S}_k -species and \mathfrak{CT}_k^{XY} as an \mathfrak{S}_{k+1} -species¹ (indicating that they carry equivariant actions of the specified group). (Hereafter, when it is necessary to distinguish, a species which is *not* equipped with any Γ -species structure will be dubbed an 'actionless' species.)

As a consequence of Lemma 4.3, then, we can then relate the rooted Γ -species forms of \mathfrak{CT}_k to the various (actionless) species forms of generic rooted k-trees in Theorem 2.1:

Theorem 4.4. For the various rooted forms of the (actionless) species \mathfrak{a}_k as in Theorem 2.1 and the various rooted Γ -species forms of \mathfrak{CT}_k as in Theorem 3.5 (interpreted as \mathfrak{S}_k - and \mathfrak{S}_{k+1} -species), we have

$$\mathfrak{a}_{k}^{Y} = \frac{\mathfrak{CT}_{k}^{Y}}{\mathfrak{S}_{k}}$$
(7a)

$$\mathfrak{a}_{k}^{XY} = \frac{\mathfrak{CT}_{k}^{XY}}{\mathfrak{S}_{k}}$$
(7b)

$$\mathfrak{a}_{k}^{X} = \frac{\mathfrak{CT}_{k}^{XY}}{\mathfrak{S}_{k+1}}$$
(7c)

as isomorphisms of (actionless) species, where \mathfrak{CT}_k^{XY} is an \mathfrak{S}_k -species in Eq. (7b) and an \mathfrak{S}_{k+1} -species in Eq. (7c).

As a result, we have explicit characterizations of all the rooted components of the original dissymmetry theorem, Theorem 2.1. Thus, through the enumerative toolset of species theory, we can enumerate k-trees through a careful enumeration of each of \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} .

5 Automorphisms and cycle indices

Species theory associates to each species F an enumerative power series Z_F dubbed the 'cycle index' which keeps track of the number of structures with a given automorphism type. Γ -species theory provides a natural extension of the cycle index for a Γ -species F, denoted Z_F^{Γ} , which keeps track of the number of γ -invariant structures with a given automorphism type for each $\gamma \in \Gamma$, defined in [8, §3]. We reprint its definition here for convenience:

$$Z_F^{\Gamma}(\gamma) := \sum_{n \ge 0} \frac{1}{n!} \sum_{\nu \in \mathfrak{S}_n} \operatorname{fix}(\gamma \cdot F[\nu]) p_{\nu}, \tag{8}$$

where $p_{\nu} = p_1^{\nu_1} p_2^{\nu_2} \dots$ (for ν_i the number of *i*-cycles of ν) is the monomial recording the cycle structure of ν . (It is important to note that the action of \mathfrak{S}_n on labels and the actions of \mathfrak{S}_k and \mathfrak{S}_{k+1} on orientations of *k*-coding trees are distinct. We will refer to label permutations again in the proof of Theorem 5.5.)

¹The \mathfrak{S}_{k} - and \mathfrak{S}_{k+1} -actions on \mathfrak{CT}_{k}^{XY} are compatible, but we will make explicit reference to \mathfrak{CT}_{k}^{XY} as an \mathfrak{S}_{k} - or \mathfrak{S}_{k+1} -species whenever it is important and not completely clear from context which we mean.

Just as with classical cycle indices (or enumerative generating functions of non-speciestheoretic combinatorics), the algebra of Γ -cycle indices is closely associated to the combinatorial algebra of their species. We will now apply the results of the preceding sections to compute the cycle indices of the various Γ -species we have developed.

k-coding trees: \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} 5.1

Corollary 2.2 of the dissymmetry theorem for k-trees has a direct analogue in terms of cycle indices:

Theorem 5.1. For the various forms of the species \mathfrak{a}_k as in Section 2, we have

$$Z_{\mathfrak{a}_k} = Z_{\mathfrak{a}_k^X} + Z_{\mathfrak{a}_k^Y} - Z_{\mathfrak{a}_k^{XY}}.$$

Thus, we need to calculate the cycle indices of the three rooted forms of \mathfrak{a}_k . A straightforward application of Burnside's lemma allows us to pass from the cycle index of a Γ -species F to the ordinary cycle index of the quotient species Γ/F :

Lemma 5.2. For a Γ -species F, the ordinary cycle index of the quotient species F/Γ is given by

$$Z_{F/\Gamma} = \overline{Z_F^{\Gamma}} := \frac{1}{|\Gamma|} \sum_{\gamma \in \Gamma} Z_F^{\Gamma}(\gamma) = \frac{1}{|\Gamma|} \sum_{\substack{n \ge 0\\\nu \in \mathfrak{S}_n\\\gamma \in \Gamma}} \frac{1}{n!} (\gamma \cdot F[\nu]) p_{\nu}.$$

where we define $\overline{Z_F^{\Gamma}} = \frac{1}{|\Gamma|} \sum_{\gamma \in \Gamma} Z_F^{\Gamma}(\gamma)$ for future convenience.

From Theorem 4.4 and by Lemma 5.2 we obtain:

Theorem 5.3. For the various forms of the species \mathfrak{a}_k as in Section 2 and the various \mathfrak{S}_k -species and \mathfrak{S}_{k+1} -species forms of \mathfrak{CT}_k as in Section 4.1, we have

$$Z_{\mathfrak{a}_{k}^{Y}} = \overline{Z_{\mathfrak{CT}_{k}^{Y}}^{\mathfrak{S}_{k}}} = \frac{1}{k!} \cdot \sum_{\sigma \in \mathfrak{S}_{k}} Z_{\mathfrak{CT}_{k}^{Y}}^{\mathfrak{S}_{k}}(\sigma)$$
(9a)

$$Z_{\mathfrak{a}_{k}^{XY}} = \overline{Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k}}} = \frac{1}{k!} \cdot \sum_{\sigma \in \mathfrak{S}_{k}} Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k}}(\sigma)$$
(9b)

$$Z_{\mathfrak{a}_{k}^{X}} = \overline{Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k+1}}} = \frac{1}{(k+1)!} \cdot \sum_{\sigma \in \mathfrak{S}_{k+1}} Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k+1}}(\sigma)$$
(9c)

We thus need only calculate the various Γ -cycle indices for the \mathfrak{S}_k -species and \mathfrak{S}_{k+1} -species forms of \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} to complete our enumeration of general k-trees. In Theorem 3.5, the functional equations for the (actionless) species \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} both include terms of the form $\mathcal{L}_k \circ \mathfrak{CT}_k^Y$. The plethysm of (actionless) species does have a generalization to Γ -species, as given in [8, §3], but it does not correctly describe the manner in which \mathfrak{S}_k acts on linear orders of \mathfrak{CT}_k^Y -structures in these recursive decompositions.

Specifically, for two Γ -species F and G, an element $\gamma \in \Gamma$ acts on an $(F \circ G)$ -structure (colloquially, 'an F-structure of G-structures') by acting on the F-structure and on each of the G-structures independently. In our action of \mathfrak{S}_k , however, the actions of σ on the descendant \mathfrak{CT}_k^Y -structures are *not* independent—they depend on the position of the structure in the linear ordering around the parent black vertex. In particular, if σ acts on some non-root black vertex, then $\rho_i(\sigma)$ acts on the white vertex in the *i*th place, where in general $\rho_i(\sigma) \neq \sigma$.

Thus, we consider automorphisms of these \mathfrak{S}_k -structures directly. First, we consider the component species $X \cdot \mathcal{L}_k(\mathfrak{CT}_k^Y)$.

Lemma 5.4. Let B be a structure of the species $F_k = X \cdot \mathcal{L}_k(\mathfrak{CT}_k^Y)$. Let W_i be the \mathfrak{CT}_k^Y -structure in the *i*th position in the linear order. Then some $\sigma \in \mathfrak{S}_k$ acts as an automorphism of B if and only if, for each $i \in [k+1]$, we have $(\rho_i \sigma) W_i \cong W_{\sigma(i)}$.

Proof. Recall that the action of $\sigma \in \mathfrak{S}_k$ is in fact the action of the lift of σ as an element of \mathfrak{S}_{k+1} . The X-label on the black root of B is not affected by the action of σ , so no conditions on σ are necessary to accommodate it. However, the \mathcal{L}_k -structure on the white children of the root is permuted by σ , and we apply to each of the W_i 's the action of $(\rho_i \sigma)$. (See Fig. 8.) Thus, σ is an automorphism of B if and only if the combination of applying σ to the linear order and $\rho_i \sigma$ to each W_i is an automorphism. Since σ 'carries' each W_i onto $W_{\sigma(i)}$, we must have that $(\rho_i \sigma) W_i \cong W_{\sigma(i)}$, as claimed. That this suffices is clear. \Box

We hereafter treat F_k as a \mathfrak{S}_k -species with respect to this action, but note that $F_k = X \cdot \mathcal{L}_k(\mathfrak{CT}_k^Y)$ is *not* an isomorphism of \mathfrak{S}_k -species.

We now have the tools in hand to find recursive functional equations satisfied by $Z_{\mathcal{CT}_k^Y}^{\mathfrak{S}_k}$ and $Z_{\mathcal{CT}_k^Y}^{\mathfrak{S}_k}$.

Theorem 5.5 (Cycle index of \mathfrak{CT}_k^Y). The \mathfrak{S}_k -cycle index for the species \mathfrak{CT}_k^Y is characterized by the recursive functional equations

$$Z_{\mathfrak{CT}_{k}^{Y}}^{\mathfrak{S}_{k}} = Z_{Y}^{\mathfrak{S}_{k}} \cdot \left(Z_{\mathcal{E}}^{\mathfrak{S}_{k}} \circ Z_{F_{k}}^{\mathfrak{S}_{k}} \right)$$
(10a)

$$Z_{F_{k}}^{\mathfrak{S}_{k}}(\sigma) = p_{1}[x] \cdot \prod_{c \in C(\sigma)} Z_{\mathfrak{CT}_{k}^{Y}}^{\mathfrak{S}_{k}} \left(\prod_{i \in c} \rho_{i}(\sigma)\right) \left(p_{|c|}[x], p_{2|c|}[x], \dots; p_{|c|}[y], p_{2|c|}[y], \dots\right).$$
(10b)

In Eq. (10a), the plethysm \circ is that of \mathfrak{S}_k -species, \mathcal{E} is the \mathfrak{S}_k -species of sets with the trivial \mathfrak{S}_k -action, and Y is the \mathfrak{S}_k -species of Y-labeled singletons with the trivial \mathfrak{S}_k -action. In Eq. (10b), $C(\sigma)$ denotes the set of cycles of σ (as a k-permutation), and the product is taken in order with respect to any choice of linearization of the cyclic order of the elements of c.

Proof. Let T be a structure of the \mathfrak{S}_k -species \mathfrak{CT}_k^Y . It is clear that T may be decomposed into a singleton of type Y, corresponding to the label on the root white vertex, and an arbitrary set (\mathcal{E} -structure) of descendant X-rooted trees, all of the same "structure type" (that is, species), which we denote F_k . Furthermore, it is apparent that some $\sigma \in \mathfrak{S}_k$ acts



Figure 8: An example F_4 -structure.

on T by acting independently on Y (by fixing the root), on \mathcal{E} (by preserving the set of neighbors of that root), and on each F_k -structure (in some more complicated manner, to be discussed in what follows. Equation (10a) follows immediately from these observations. It remains only to analyze the \mathfrak{S}_k -species F_k of these X-rooted descendant structures.

By Lemma 5.4, if $\sigma \cdot CT_k^Y[\pi, \tau]$ fixes some F_k structure, the $\rho_i(\sigma)$ -image of the CT_k^Y structure at position *i* must be its $(CT_k^Y[\pi, \tau])^{-1}$ -image also. Furthermore, if *i* is in a cycle *c* of length |c|, then all the other CT_k^Y -structures along *c* are determined by the choice of the structure at position *i*, and that CT_k^Y -structure must be sent to *itself* by $\prod_{j\in c} \rho_j(\sigma)$. The action of this permutation on that structure must then be identical to that of $(CT_k^Y[\pi, \tau])^{-|c|}$, which we observe must then restrict to an automorphism of that CT_k^Y -structure, each *l*-cycle of which corresponds to an $(l \cdot |c|)$ -cycle of $CT_k^Y[\pi, \tau]$. Now we consider the \mathfrak{S}_k -cycle index $Z_{F_k}^{\mathfrak{S}_k}$ of this \mathfrak{S}_k -species F_k . Fix partitions π

Now we consider the \mathfrak{S}_k -cycle index $Z_{F_k}^{\mathfrak{S}_k}$ of this \mathfrak{S}_k -species F_k . Fix partitions π (representing the cycle type of a permutation acting on black (x-) vertex labels) and τ (similarly for white (y-) labels) and some $\sigma \in \mathfrak{S}_k$. (In what follows, we let $\hat{\lambda}$ denote some arbitrarily-chosen permutation of cycle type λ when needed.) Then the coefficient on the $p_{\pi}[x]p_{\tau}[y]$ -term of $Z_{F_k}^{\mathfrak{S}_k}(\sigma)$ is $\frac{1}{|\pi|!|\tau|!}$ times the number of F_k -structures for which $\sigma \cdot \mathfrak{CT}_k^Y[\hat{\pi}, \hat{\tau}]$ is an automorphism. Let $f(\pi, \tau)$ denote this number. Suppose that $f(\pi, \tau)$ is nonzero. By the above, there must exist decompositions² $\pi - \{1\} = \bigsqcup_{i} \pi_{i}$ and $\tau = \bigsqcup_{i} \tau_{i}$, each having $|C(\sigma)|$ components, such that all the parts of π_{i} and τ_{i} are multiples of the length of the *i*th cycle of σ . Let $\frac{\pi_{i}}{|c_{i}|}$ and $\frac{\tau_{i}}{|c_{i}|}$ denote the partitions resulting from dividing each part of π_{i} and τ_{i} respectively by $|c_{i}|$. Then, for each such decomposition, an F_{k} -structure may be assembled by choosing, for each cycle c_{i} , a \mathfrak{CT}_{k}^{Y} -structure for which $\left(\prod_{j \in c_{i}} \rho_{j}(\sigma)\right) \cdot \mathfrak{CT}_{k}^{Y} \left[\frac{\hat{\pi}_{i}}{|c_{i}|}, \frac{\hat{\tau}_{i}}{|c_{i}|}\right]$ is an automorphism, then distributing its ρ -images around the $|c_{i}|$ positions of the cycle. The number of such choices is exactly $\left|\frac{\pi_{i}}{|c_{i}|}\right|! \left|\frac{\tau_{i}}{|c_{i}|}\right|!$ times the coefficient of the $p_{\frac{\hat{\pi}_{i}}{|c_{i}|}}[x]p_{\frac{\hat{\tau}_{i}}{|c_{i}|}}[y]$ -term of $Z_{\mathfrak{CT}_{k}^{\mathfrak{S}_{k}}\left(\prod_{j \in c_{i}} \rho_{j}(\sigma)\right)$. Accordingly, the total number of F_{k} -structures for which $\sigma \cdot \mathfrak{CT}_{k}^{Y}[\hat{\pi}, \hat{\tau}]$ is an automorphism is exactly $|\pi|! |\tau|!$ times the coefficient of $p_{\pi-[1]}[x]p_{\tau}[y]$ in $\prod_{c \in C[\sigma]} Z_{\mathfrak{CT}_{k}^{\mathfrak{S}_{k}}}\left(\prod_{i \in c} \rho_{i}(\sigma)\right)\left(p_{|c|}[x], p_{2|c|}[x], \ldots; p_{|c|[y], p_{2|c|}[y], \ldots}\right)$. Equation (10b) follows immediately.

Theorem 5.6 (Cycle index of \mathfrak{CT}_k^{XY}). The \mathfrak{S}_{k+1} -cycle index for the species \mathfrak{CT}_k^{XY} is given by

$$Z_{\mathcal{CT}_{k}^{XY}}^{\mathfrak{S}_{k+1}}(\sigma) = p_{1}[x] \cdot \prod_{c \in C(\sigma)} Z_{\mathcal{CT}_{k}^{Y}}^{\mathfrak{S}_{k}} \left(\prod_{i \in c} \rho_{i}[\sigma]\right) \left(p_{|c|}[x], p_{2|c|}[x], \dots, p_{|c|}[y], p_{2|c|}[y], \dots\right).$$
(11)

under the same conditions as Theorem 5.5.

Proof. The proof is essentially identical to the part of that of Theorem 5.5 which concerns Eq. (10b).

Terms of the form $\prod_{i \in c} \rho_i(\sigma)$ appear in Eqs. (10) and (11). For the simplification of calculations, we note here two useful results about these products.

First, we observe that certain ρ -maps preserve cycle structure:

Lemma 5.7. Let $\sigma \in \mathfrak{S}_k$ be a permutation of which $i \in [k]$ is a fixed point. Then $\rho_i(\sigma)$ has the same cycle type as σ .

Proof. Let $\theta_i \in \mathfrak{S}_k$ denote the permutation given by $\theta_i(a) = a + i$ reduced modulo k + 1 as in the proof of Theorem 4.2. Then, since i is a fixed point of σ , we have that, for each $a \in \{1, 2, \ldots, k + 1\}$,

$$\rho_i(\sigma)(a) = \sigma(i+a) - i = (\theta_i \sigma \theta_i^{-1})(a),$$

so $\rho_i(\sigma)$ is a conjugate of σ and thus has the same cycle structure as σ , as desired. \Box

But then we note that the products in the above theorems are in fact permutations obtained by applying such ρ -maps:

Lemma 5.8. Let $\sigma \in \mathfrak{S}_k$ be a permutation with a cycle c. Then the cycle type of $(\prod_{i \in c} \rho_i(\sigma))$ is the same as that of $\sigma^{|c|}$. In other words, $\lambda(\prod_{i \in c} \rho_i(\sigma)) = \lambda(\sigma^{|c|})$.

 $^{^{2}}$ By 'decomposition' of a partition we mean a partition of it as a multiset—i.e. a choice of sub-multisets whose pairwise intersections are trivial and whose union is the original partition.

Proof. Let $c = (c_1, c_2, \ldots, c_{|c|})$. First, we calculate:

$$\begin{split} \prod_{i=1}^{|c|} \rho_{c_i}(\sigma) =& \rho_{c_{|c|}}(\sigma) \circ \dots \circ \rho_{c_2}(\sigma) \circ \rho_{c_1}(\sigma) \\ &= \rho_{c_{|c|}}(\sigma) \circ \dots \circ \rho_{c_2}(\sigma)(a \mapsto \sigma(c_1 + a) - \sigma(c_1)) \\ &= \rho_{c_{|c|}}(\sigma) \circ \dots \circ \rho_{c_3}(\sigma)(a \mapsto \sigma(c_2 + \sigma(c_1 + a) - \sigma(c_1)) - \sigma(c_2)) \\ &= \rho_{c_{|c|}}(\sigma) \circ \dots \circ \rho_{c_3}(\sigma)(a \mapsto \sigma^2(c_1 + a) - \sigma^2(c_1)) \\ &\vdots \\ &= a \mapsto \sigma^{|c|}(c_1 + a) - \sigma^{|c|}(c_1) \\ &= \rho_{c_1}(\sigma^{|c|}). \end{split}$$

But c_1 is a fixed point of $\sigma^{|c|}$, so by the result of Lemma 5.7, this has the same cycle structure as $\sigma^{|c|}$.

As a result, we can bypass all calculations of $\rho_i(\sigma)$ in the computation of terms of the cycle indices we have developed.

We also note an important general fact about the cycle indices of Γ -species:

Lemma 5.9. Let F be a Γ -species. Then $Z_F^{\Gamma}(\gamma)$ is a class function of γ .

This will simplify computational enumeration of k-trees significantly when k is large, since the number of elements of \mathfrak{S}_k is factorial in k while the number of conjugacy classes (indexed by partitions) is exponential in k.

5.2 k-trees: \mathfrak{a}_k

We now have all the pieces in hand to apply Theorem 5.1 to compute the cycle index of the species \mathfrak{a}_k of general k-trees. Theorem 5.1 characterizes the cycle index of the generic k-tree species \mathfrak{a}_k in terms of the cycle indices of the rooted species \mathfrak{a}_k^X , \mathfrak{a}_k^Y , and \mathfrak{a}_k^{XY} ; Theorem 4.4 gives the cycle indices of these three rooted species in terms of the Γ -cycle indices $Z_{\mathfrak{CT}_k^Y}^{\mathfrak{S}_k}$, $Z_{\mathfrak{CT}_k^{XY}}^{\mathfrak{S}_k}$, and $Z_{\mathfrak{CT}_k^{XY}}^{\mathfrak{S}_{k+1}}$; and, finally, Theorems 5.5 and 5.6 give these Γ -cycle indices explicitly. By tracing the formulas in Eqs. (10) and (11) back through this sequence of functional relationships, we can conclude:

Theorem 5.10 (Cycle index for the species of k-trees). For \mathfrak{a}_k the species of general k-trees, $Z_{\mathfrak{CT}_k^Y}^{\mathfrak{S}_k}$ as in Eq. (10), and $Z_{\mathfrak{CT}_k^{XY}}^{\mathfrak{S}_{k+1}}$ as in Eq. (11) we have:

$$Z_{\mathfrak{a}_{k}} = \frac{1}{(k+1)!} \cdot \sum_{\sigma \in \mathfrak{S}_{k+1}} Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k+1}}(\sigma) + \frac{1}{k!} \cdot \sum_{\sigma \in \mathfrak{S}_{k}} Z_{\mathfrak{CT}_{k}^{Y}}^{\mathfrak{S}_{k}}(\sigma) - \frac{1}{k!} \cdot \sum_{\sigma \in \mathfrak{S}_{k}} Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k}}(\sigma)$$
(12a)

$$=\overline{Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k+1}}}+\overline{Z_{\mathfrak{CT}_{k}^{Y}}^{\mathfrak{S}_{k}}}-\overline{Z_{\mathfrak{CT}_{k}^{XY}}^{\mathfrak{S}_{k}}}.$$
(12b)

Equation (12) in fact represents a recursive system of functional equations, since the formulas for the Γ -cycle indices of \mathfrak{CT}_k^Y and \mathfrak{CT}_k^{XY} are recursive. Computational methods can yield explicit enumerative results. However, a bit of care will allow us to reduce the computational complexity of this problem significantly.

6 Unlabeled enumeration and the generating function $\tilde{\mathfrak{a}}_k(x)$

Equation (12) in Theorem 5.10 gives a recursive formula for the cycle index of the (actionless) species \mathfrak{a}_k of k-trees. The number of unlabeled k-trees with n hedra is historically an open problem, but it is straightforward to extract their ordinary generating function from the cycle index $Z_{\mathfrak{a}_k}$ once it is computed. Actually computing terms of the cycle index in order to derive the coefficients of the generating function is, however, a computationally expensive process, since the cycle index is by construction a power series in two infinite sets of variables. The computational process can be simplified significantly by taking advantage of the relatively straightforward combinatorial structure of the structural decomposition used to derive the recursive formulas for the cycle index.

For a Γ -species F, the ordinary generating function $F_{\gamma}(x)$ counting unlabeled γ -invariant F-structures is given by

$$\tilde{F}(\gamma)(x) = Z_F^{\Gamma}(\gamma)(x, x^2, x^3, \ldots)$$

and the ordinary generating function for counting unlabeled F/Γ -structures is given by

$$\tilde{F}(x) = \frac{1}{|\Gamma|} \sum_{\gamma \in \Gamma} \tilde{F}(\gamma)(x).$$

These formula admits an obvious multisort extension, but we in fact wish to count k-trees with respect to just one sort of label (the X-labels on hedra), so we will not deal with multisort issues here. Each of our two-sort cycle indices can be converted to one-sort by substituting $p_i[y] = 1$ for all *i*. For the rest of this section, we will deal directly with these one-sort versions of the cycle indices.

We begin by considering the explicit recursive functional equations in Theorems 5.5 and 5.6. In each case, by the above, the ordinary generating function is exactly the result of substituting $p_i[x] = x^i$ and $p_i[y] = 1$ into the given formula. Thus, we have:

Theorem 6.1. For CT_k^Y the \mathfrak{S}_k -species of Y-rooted k-coding trees and CT_k^{XY} the \mathfrak{S}_{k+1} -species of edge-rooted k-coding trees, the corresponding single-variable Γ -ordinary generating functions are given by

$$\widetilde{\mathfrak{CT}_k^Y}(\sigma)(x) = \exp\left(\sum_{n \ge 1} \left(\frac{x^n}{n} \cdot \prod_{c \in C(\sigma^n)} \widetilde{\mathfrak{CT}_k^Y}\left(\prod_{i \in c} \rho_i(\sigma^n)\right)(x^{n|c|})\right)\right)$$
(13a)

and

$$\widetilde{\mathfrak{CT}_{k}^{XY}}(\sigma)(x) = x \cdot \prod_{c \in C(\sigma)} \widetilde{\mathfrak{CT}_{k}^{Y}} \Big(\prod_{i \in c} \rho_{i}(\sigma) \Big) \big(x^{|c|} \big).$$
(13b)

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where $\widetilde{\operatorname{CT}_k^Y}$ is an \mathfrak{S}_k -generating function and $\widetilde{\operatorname{CT}_k^{XY}}$ is an \mathfrak{S}_{k+1} -generating function.

However, as a consequence of Lemmas 5.8 and 5.9, we can simplify these expressions significantly. Let λ be the function mapping each permutation to its cycle type interpreted as an integer partition. Then:

Corollary 6.2. For \mathfrak{CT}_k^Y the \mathfrak{S}_k -species of Y-rooted k-coding trees and \mathfrak{CT}_k^{XY} the \mathfrak{S}_{k+1} -species of edge-rooted k-coding trees, the corresponding single-variable Γ -ordinary generating functions are given by

$$\widetilde{\operatorname{CT}_k^Y}(\lambda)(x) = \exp\left(\sum_{n \ge 1} \left(\frac{x^n}{n} \cdot \prod_{i \in \lambda^n} \widetilde{\operatorname{CT}_k^Y}(\lambda^{ni})(x^{ni})\right)\right)$$
(14a)

and

$$\widetilde{\mathfrak{CT}_{k}^{XY}}(\lambda)(x) = x \cdot \prod_{i \in \lambda} \widetilde{\mathfrak{CT}_{k}^{Y}}(\lambda^{i} - \{1\})(x^{i})$$
(14b)

where $\prod_{i \in \lambda}$ denotes a product over the parts *i* of λ taken with multiplicity, where λ^i denotes the *i*th 'partition power' of λ — that is, if σ is any permutation of cycle type λ , then λ^i denotes the cycle type of σ^i — and where $f(\lambda)(x)$ denotes the value of $f(\sigma)(x)$ for any σ of cycle type λ .

To clarify the notation, we work out the case k = 2 more explicitly here:

Example 6.3 (k = 2). There are only two partitions of k = 2: $\{1, 1\}$ and $\{2\}$, corresponding to the two permutations *id* and (12) respectively. Since $\{1, 1\}^i = \{1, 1\}$ for any *i*, we have that

$$\widetilde{\mathfrak{CT}_2^Y}(\{1,1\})(x) = \exp\left(\sum_{n \ge 1} \frac{x^n}{n} \widetilde{\mathfrak{CT}_2^Y}(\{1,1\})(x^n)^2\right).$$

The situation for $\{2\}$ is slightly more complex, since $\{2\}^i = \{1, 1\}$ if *i* is even but $\{2\}$ if *i* is odd. Thus, we have

$$\widetilde{\mathfrak{CT}_{2}^{Y}}(\{2\})(x) = \exp\left(\sum_{n \ge 1} \left(\frac{x^{2n-1}}{2n-1}\widetilde{\mathfrak{CT}_{2}^{Y}}(\{2\})(x^{2n-1})^{2}\right) + \left(\frac{x^{2n}}{2n}\widetilde{\mathfrak{CT}_{2}^{Y}}(\{1,1\})(x^{2n})^{2}\right)\right).$$

Conventional computational techniques (such as those demonstrated in Appendix B) then suffice to compute that

$$\widetilde{CT}_{2}^{Y}(\{1,1\})(x) = 1 + x + 3x^{2} + 10x^{3} + 39x^{4} + 160x^{5} + \dots$$

and

$$\widetilde{\operatorname{CT}}_{2}^{Y}(\{2\})(x) = 1 + x + x^{2} + 2x^{3} + 3x^{4} + 6x^{5} + \dots$$

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Corollary 6.2 characterizes the ordinary generating functions $\widetilde{\mathbb{CT}_k^Y}$ and $\widetilde{\mathbb{CT}_k^{XY}}$. The cycle index of the species \mathfrak{a}_k , as seen in Eq. (12), is given simply in terms of quotients of the cycle indices of the two Γ -species \mathbb{CT}_k^Y and \mathbb{CT}_k^{XY} , and this result can pass to the generating-function level. Thus, we also have:

Theorem 6.4. For \mathfrak{a}_k the species of k-trees and $\widetilde{\mathfrak{CT}_k^Y}$ and $\widetilde{\mathfrak{CT}_k^{XY}}$ as in Theorem 6.1, we have

$$\tilde{\mathfrak{a}}_{k}(x) = \frac{1}{(k+1)!} \cdot \sum_{\sigma \in \mathfrak{S}_{k+1}} \widetilde{\operatorname{CT}_{k}^{XY}}(\sigma)(x) + \frac{1}{k!} \cdot \sum_{\sigma \in \mathfrak{S}_{k}} \widetilde{\operatorname{CT}_{k}^{Y}}(\sigma)(x) - \frac{1}{k!} \cdot \sum_{\sigma \in \mathfrak{S}_{k}} \widetilde{\operatorname{CT}_{k}^{XY}}(\sigma)(x).$$
(15)

Then, as a consequence of Lemma 5.9 and Corollary 6.2, we can instead write

Corollary 6.5 (Generating function for unlabeled k-trees). For \mathfrak{a}_k the species of k-trees and $\widetilde{\operatorname{CT}}_k^Y$ and $\widetilde{\operatorname{CT}}_k^{XY}$ as in Corollary 6.2, we have

$$\tilde{\mathfrak{a}}_{k}(x) = \sum_{\lambda \vdash k+1} \frac{1}{z_{\lambda}} \widetilde{\mathfrak{CT}_{k}^{XY}}(\lambda)(x) + \sum_{\lambda \vdash k} \frac{1}{z_{\lambda}} \widetilde{\mathfrak{CT}_{k}^{Y}}(\lambda)(x) - \sum_{\lambda \vdash k} \frac{1}{z_{\lambda}} \widetilde{\mathfrak{CT}_{k}^{XY}}(\lambda \cup \{1\})(x).$$
(16)

This direct characterization of the ordinary generating function of unlabeled k-trees, while still recursive, is much simpler computationally than the characterization of the full cycle index in Eq. (12). For computation of the number of unlabeled k-trees, it is therefore much preferred. Classical methods for working with recursively-defined power series suffice to extract the coefficients quickly and efficiently. The results of some such explicit calculations are presented in Appendix A.

7 Special-case behavior for small k

Many of the complexities of the preceding analysis apply only for $k \ge 3$. In the cases k = 1 and k = 2, our analysis simplifies dramatically, and effectively reduces to previous work.

7.1 Ordinary trees (k = 1)

When k = 1, an \mathfrak{a}_k -structure is merely an ordinary tree with X-labels on its edges and Y-labels on its vertices. There is no internal symmetry of the form that the actions of \mathfrak{S}_k are intended to break. The actions of \mathfrak{S}_2 act on ordinary trees rooted at a *directed* edge, with the nontrivial element $\tau \in \mathfrak{S}_2$ acting by reversing this orientation. The resulting decomposition from the dissymmetry theorem in Theorem 2.1 and the recursive functional equations of Theorem 3.5 then clearly reduce to the classical dissymmetry analysis of ordinary trees.

7.2 2-trees

When k = 2, there is a nontrivial symmetry at fronts (edges); two triangles may be joined at an edge in two distinct ways. The imposition of a coherent orientation on a 2-tree by directing one of its edges breaks this symmetry; the action of \mathfrak{S}_2 by reversal of these orientations gives unoriented 2-trees as its orbits. The defined action of \mathfrak{S}_3 on edge-rooted oriented triangles is simply the classical action of the dihedral group D_6 on a triangle, and its orbits are unoriented, unrooted triangles. We further note that ρ_i is the trivial map on \mathfrak{S}_2 and that $\rho_i(\sigma) = (12)$ for $\sigma \in \mathfrak{S}_3$ if and only if σ is an odd permutation, both regardless of *i*. We then have that:

$$Z_{\mathcal{CT}_{2}^{Y}}^{\mathfrak{S}_{2}} = p_{1}[y] \cdot Z_{\mathcal{E}} \circ \left(p_{1}[x] \cdot \prod_{c \in C(\sigma)} Z_{\mathcal{CT}_{2}^{Y}}^{\mathfrak{S}_{2}}(e) \left(p_{|c|}[x], p_{2|c|}[x], \dots; p_{|c|}[y], p_{2|c|}[y], \dots \right) \right)$$
(17a)

and

$$Z_{\mathcal{CT}_{2}^{XY}}^{\mathfrak{S}_{3}} = p_{1}[x] \cdot \prod_{c \in C(\sigma)} Z_{\mathcal{CT}_{2}^{Y}}^{\mathfrak{S}_{2}} \left(\rho(\sigma)^{|c|} \right) \left(p_{|c|}[x], p_{2|c|}[x], \dots; p_{|c|}[y], p_{2|c|}[y], \dots \right).$$
(17b)

where, by abuse of notation, we let ρ represent any ρ_i . By the previous, the argument $\rho(\sigma)^{|c|}$ in Eq. (17b) is τ if and only if σ is an odd permutation and c is of odd length. This analysis and the resulting formulas for the cycle index $Z_{\mathfrak{a}_2}$ are essentially equivalent to those derived in [4].

A Enumerative tables

With the recursive functional equations for cycle indices of Section 5, we can calculate the explicit cycle index for the species \mathfrak{a}_k to any finite degree we choose using computational methods; this cycle index can then be used to enumerate both unlabeled and labeled (at fronts, hedra, or both) k-trees up to a specified number n of hedra (or, equivalently, kn + 1 fronts). We have done so here for $k \leq 7$ and $n \leq 30$ using Sage 5.0 [13] using code available in Appendix B. The resulting values appear in Table 1.

We note that both unlabeled and hedron-labeled enumerations of k-trees stabilize in k:

Theorem A.1. For $k \ge n-2$, the numbers of unlabeled and hedron-labeled k-trees with n hedra are independent of k.

Proof. We show that the species \mathfrak{a}_k and \mathfrak{a}_{k+1} have contact up to order k+2 by explicitly constructing a natural bijection. We note that in a (k+1)-tree with no more than k+2 hedra, there will exist at least one vertex which is common to *all* hedra. For any k-tree with no more than k+2 hedra, we can construct a (k+1)-tree with the same number of hedra by adding a single vertex and connecting it by edges to every existing vertex; we can then pass labels up from the (k+1)-cliques which are the hedra of the k-tree to the (k+2)-cliques which now sit over them. The resulting graph will be a (k+1)-tree whose (k+1)-tree hedra are adjacent exactly when the k-tree hedra they came from were adjacent.

Therefore, any two distinct k-trees will pass to distinct (k + 1)-trees. Similarly, for any (k + 1)-tree with no more than k + 2 hedra, choose one of the vertices common to all the hedra and remove it, passing the labels of (k + 1)-tree hedra down to the k-tree hedra constructed from them; again, adjacency of hedra is preserved. This of course creates a k-tree, and for distinct (k + 1)-trees the resulting k-trees will be distinct. Moreover, by symmetry the result is independent of the choice of common vertex, in the case there is more than one.

However, thus far we have neither determined a direct method for computing these stabilization numbers nor identified a straightforward combinatorial characterization of the structures they represent.

	(a) $k = 1$		(b) $k = 2$
n	Unlabeled 1-trees	n	Unlabeled 2-trees
0	1	0	1
1	1	1	1
2	1	2	1
3	2	3	2
4	3	4	5
5	6	5	12
6	11	6	39
7	23	7	136
8	47	8	529
9	106	9	2171
10	235	10	9368
11	551	11	41534
12	1301	12	188942
13	3159	13	874906
14	7741	14	4115060
15	19320	15	19602156
16	48629	16	94419351
17	123867	17	459183768
18	317955	18	2252217207
19	823065	19	11130545494
20	2144505	20	55382155396
21	5623756	21	277255622646
22	14828074	22	1395731021610
23	39299897	23	7061871805974
24	104636890	24	35896206800034
25	279793450	25	183241761631584
26	751065460	26	939081790240231
27	2023443032	27	4830116366008952
28	5469566585	28	24927175920361855
29	14830871802	29	129047003236769110
30	40330829030	30	670024248072778235

Table 1: Enumerative data for k-trees with n hedra

	(c) $k = 3$		(d) $k = 4$
n	Unlabeled 3-trees	n	Unlabeled 4-trees
0	1	0	1
1	1	1	1
2	1	2	1
3	2	3	2
4	5	4	5
5	15	5	15
6	58	6	64
7	275	7	331
8	1505	8	2150
9	9003	9	15817
10	56931	10	127194
11	372973	11	1077639
12	2506312	12	9466983
13	17165954	13	85252938
14	119398333	14	782238933
15	841244274	15	7283470324
16	5993093551	16	68639621442
17	43109340222	17	653492361220
18	312747109787	18	6276834750665
19	2286190318744	19	60759388837299
20	16826338257708	20	592227182125701
21	124605344758149	21	5808446697002391
22	927910207739261	22	57289008242377068
23	6945172081954449	23	567939935463185078
24	52225283886702922	24	5656700148512008902
25	394398440097305861	25	56583199285317631541
26	2990207055800156659	26	568236762643725657852
27	22753619938517594709	27	5727423267612393252616
28	173727411594289881739	28	57924486783495226147615
29	1330614569159767263501	29	587672090447840337304025
30	10221394007530945428347	30	5979782184127687211698807

Enumerative data for k-trees with n hedra, continued

Enumerative d	lata for	<i>k</i> -trees	with η	n hedra,	continued
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	(e) $k = 5$		(f) $k = 6$
n	Unlabeled 5-trees	$n \mid$	Unlabeled 6-trees
0	1	0	1
1	1	1	1
2	1	2	1
3	2	3	2
4	5	4	5
5	15	5	15
6	64	6	64
7	342	7	342
8	2321	8	2344
9	18578	9	19090
10	168287	10	179562
11	1656209	11	1878277
12	17288336	12	21365403
13	188006362	13	258965451
14	2105867058	14	3294561195
15	24108331027	15	43472906719
16	280638347609	16	589744428065
17	3310098377912	17	8171396893523
18	39462525169310	18	115094557122380
19	474697793413215	19	1642269376265063
20	5754095507495584	20	23679803216530017
21	70216415130786725	21	344396036645439675
22	861924378411516159	22	5045351124912000756
23	10636562125193377459	23	74375422235109338507
24	131890971196221692874	24	1102368908826371717478
25	1642577274341274449247	25	16417712341047912048640
26	20538830517384955820622	26	245566461812077209025580
27	257767439475728146293796	27	3687384661929075391318298
28	3246108646710813383678978	28	55566472746158319169779382
29	41008581189552637540038747	29	840092106663809502446963972
30	519599497193547405843864376	30	12739517442131428048314937036

Enumerative data for k -trees with n hedra, contin	nued
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(g) $k = 7$			
n	Unlabeled 7-trees		
0	1		
1	1		
2	1		
3	2		
4	5		
5	15		
6	64		
7	342		
8	2344		
9	19137		
10	181098		
11	1922215		
12	22472875		
13	284556458		
14	3849828695		
15	54974808527		
16	819865209740		
17	12655913153775		
18	200748351368185		
19	3253193955012557		
20	53619437319817482		
21	895778170144927928		
22	15129118461773051724		
23	257812223121779545108		
24	4426056869082751747930		
25	76463433541541506345648		
26	1328088941166844504424628		
27	23175796698013212039339479		
28	406103563562864890670029228		
29	7142350290468621849814034057		
30	126034923903699365819345698783		

Enumerative data for k-trees with n hedra, continu
--

	(ii) $n = 0$
n	Unlabeled 8-trees
0	1
1	1
2	1
3	2
4	5
5	15
6	64
7	342
8	2344
9	19137
10	181204
11	1926782
12	22638677
13	289742922
14	3996857019
15	58854922207
16	916955507587
17	14988769972628
18	255067524402905
19	4487202163529135
20	81112295567987808
21	1498874117898285574
22	28195965395340358096
23	538126404726276758908
24	10391826059632904271057
25	202624626664206041379718
26	3982593421723767068438772
27	78804180647706388187446055
28	1568191570016583843925943321
29	31359266621157738864915907470
30	629755261439815181073415721542

(h) k = 8

Enumerative data for	k-trees	with n	hedra,	continued
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	(1) $\kappa = 9$
n	Unlabeled 9-trees
0	1
1	1
2	1
3	2
4	5
5	15
6	64
7	342
8	2344
9	19137
10	181204
11	1927017
12	22652254
13	290351000
14	4019973352
15	59642496465
16	941751344429
17	15724551551655
18	275926445572426
19	5057692869843759
20	96275031338911591
21	1892687812366295682
22	38234411627616084843
23	790120238796588845615
24	16638524087850961727575
25	355878246778832856290372
26	7710423952280397990026132
27	168843592748278228259801752
28	3730285520855433827693340329
29	83027821492843727307516904184
30	1859625249087075723295908757282

(i) k = 9

Enumerative data for k-trees with n hedra, continued

(j) $k = 10$		
n	Unlabeled 10-trees	
0	1	
1	1	
2	1	
3	2	
4	5	
5	15	
6	64	
7	342	
8	2344	
9	19137	
10	181204	
11	1927017	
12	22652805	
13	290391147	
14	4022154893	
15	59741455314	
16	945737514583	
17	15871943695637	
18	281035862707569	
19	5226147900656616	
20	101612006684523937	
21	2056425123910104429	
22	43127730369661586804	
23	933229734601789336024	
24	20749443766669472108394	
25	472211306357077710523863	
26	10961384502758318928846970	
27	258737420965101611169934566	
28	6193917223279376307682721853	
29	150039339181032274342778699887	
30	3670778410024403632885217999313	

B Code listing

The recursive functional equations in Eqs. (14a), (14b) and (16) characterize the ordinary generating function $\tilde{\mathfrak{a}}_k(x)$ for unlabeled general k-trees. Code to compute the coefficients of this generating function using the computer algebra system Sage 5.0 [13] explicitly follows in listing 1. Specifically, the generating function for unlabeled k-trees may be computed to degree n by copying the included code into a Sage notebook, modifying the final line with the desired values of k and n, and executing.

This code takes full advantage of Lemmas 5.8 and 5.9 to minimize the number of distinct calculations which must be performed; as a result, it is able to compute the number of k-trees on up to n hedra quickly even for relatively large k and n. For example, the first thirty terms of the generating function for 8-trees in Appendix A were computed on a modern desktop-class computer in approximately two minutes.

Listing 1: Sage code to compute numbers of k-trees)

```
1 # Set up a ring of formal power series
   psr = PowerSeriesRing(QQ, 'x')
3 x = psr.gen()
  # Compute the generating function for unlabeled Y-rooted k-trees fixed
5
     by permutations of a given cycle type mu.
   # Note that mu should partition k
  @cached_function
7
   def unlY(mu, n):
9
       if n \le 0:
           return psr(1)
11
       else:
           ystretcher = lambda c, part:
              unlY(Partition(partition_power(part, c)),
              floor((n-1)/c)).subs({x:x**c})
13
           descendant_pseries = lambda part: prod(ystretcher(c, part) for
              c in part)
           return sum(x**i/i * descendant_pseries(partition_power(mu,
              i)).subs({x:x**i}) for i in xrange(1, n+1)).exp(n+1)
15
   # Compute the generating function for unlabeled XY-rooted k-trees
     fixed by permutations of a given cycle type mu.
17 # Note that mu should partition k+1
   @cached_function
19 def unlXY(mu, n):
       if n \le 0:
21
           return psr(0)
       else:
23
           ystretcher = lambda c: unlY(Partition(partition_power(mu,
              c)[:-1]), floor((n-1)/c)).subs({x:x**c})
           return (x * prod(ystretcher(c) for c in mu)).add_bigoh(n+1)
25
   # Compute the generating functions for unlabeled X-, Y-, and XY-rooted
     k-trees using quotients
```

```
27 ax = lambda k, n: sum(1/mu.aut() * unlXY(mu, n) for mu in
Partitions(k+1))
ay = lambda k, n: sum(1/mu.aut() * unlY(mu, n) for mu in Partitions(k))
29 axy = lambda k, n: sum(1/mu.aut() * unlXY(Partition(mu + [1]), n) for
mu in Partitions(k))
31 # Compute the generating function for unlabeled un-rooted k-trees
using the dissymmetry theorem
a = lambda k, n: ax(k, n) + ay(k, n) - axy(k, n)
33
# Print the result
35 # ALERT: User must substitue values for k and n (number of hedra)
print a(kval, nval)
```

Acknowledgments

The author wishes to express his profound gratitude to Ira Gessel, who served as advisor and committee chair for the dissertation in which this research originated. His guidance, advice, and insightful mentorship were invaluable throughout that process.

The author also offers his thanks to an anonymous referee, whose close reading and thorough commentary led to the correction of numerous small errors and oversights and whose broader recommendations led to significant improvement of the structure and narrative flow of the paper as a whole.

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