Biclique Covers and Partitions

Trevor Pinto^{*}

School of Mathematical Sciences Queen Mary University of London London E1 4NS, UK.

tpinto@maths.qmul.ac.uk

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Abstract

The biclique cover number (resp. biclique partition number) of a graph G, bc(G) (resp. bp(G)), is the least number of bicliques—complete bipartite subgraphs—that are needed to cover (resp. partition) the edges of G.

The local biclique cover number (resp. local biclique partition number) of a graph G, lbc(G) (resp. lbp(G)), is the least r such that there is a cover (resp. partition) of the edges of G by bicliques with no vertex in more than r of these bicliques.

We show that bp(G) may be bounded in terms of bc(G), in particular, $bp(G) \leq \frac{1}{2}(3^{bc(G)} - 1)$. However, the analogous result does not hold for the local measures. Indeed, in our main result, we show that lbp(G) can be arbitrarily large, even for graphs with lbc(G) = 2. For such graphs, G, we try to bound lbp(G) in terms of additional information about biclique covers of G. We both answer and leave open questions related to this.

There is a well known link between biclique covers and subcube intersection graphs. We consider the problem of finding the least r(n) for which every graph on n vertices can be represented as a subcube intersection graph in which every subcube has dimension r. We reduce this problem to the much studied question of finding the least d(n) such that every graph on n vertices is the intersection graph of subcubes of a d-dimensional cube.

Keywords: biclique cover; biclique partition; local biclique cover number; subcube intersection graphs;

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1 Introduction

1.1 Biclique covers and partitions

A biclique cover of a (simple) graph, G, is a collection of bicliques (complete bipartite subgraphs) of G, the union of whose edges is E(G). A related notion is that of a biclique partition: a collection of bicliques of G, whose edges partition the edges of G. We say that a biclique cover (resp. partition) is a k-cover (resp. k-partition) if it contains at most k bicliques. The biclique cover number of G, bc(G), is the least k for which there exists a k-cover of G. Analogously, the biclique partition number of G, bp(G), is the least k for which there exists a k-partition of G. The biclique cover number and biclique partition number have been studied extensively, see for instance [16] and [11], motivated by applications to many other areas such as combinatorial geometry [1], communication complexity [13], network addressing [8] and even immunology [17].

One of the early results on the biclique partition number is the Graham-Pollak Theorem [8]—see [20] for an elegant proof—which concerns K_n , the complete graph on n vertices. More specifically, it states that $bp(K_n) = n-1$. In contrast to this, it is easy to show that $bc(K_n) = \lceil \log n \rceil$. (Here and elsewhere, log refers to \log_2). These results demonstrate that the trivial inequality $bc(G) \leq bp(G)$ may be quite loose: $bp(K_n) \geq 2^{bc(K_n)-1} - 1$.

In Section 3, we investigate how large bp can be in terms of bc. Indeed, we prove that if bc(G) = t, $bp(G) \leq (3^t - 1)/2$ and exhibit a graph showing that this is tight.

We call a biclique cover (resp. partition) r-local if every vertex is in at most r of the bicliques involved in the cover (resp. partition) of G. The local biclique cover number, lbc(G), is the least r such that there exists an r-local cover of the graph. Similarly, the local biclique partition number, lbp(G), is the least r such that there exists an r-local partition of the graph. Variants of lbp and lbc have long been studied, starting in 1967 with Katona and Szemerédi [12] in connection with diameter 2 graphs with few edges. An easy corollary of their result is:

$$\operatorname{lbc}(K_n) = \operatorname{lbp}(K_n) = \lceil \log n \rceil,$$

which was also shown in full by Dong and Liu [4]. This result is perhaps surprising given that $bc(K_n)$ and $bp(K_n)$ differ so vastly.

One of the key results in the theory of the local biclique cover/ partition numbers, is that of Lupanov [14] in 1956, stating that for all n, there exists a graph G on n vertices with $lbc(G) \ge \frac{c_1 n}{\log n}$, where c_1 is a positive constant. Indeed his result is slightly stronger than stated—he actually showed the same bound holds if lbc(G) is replaced by w(G), the minimum over all covers, κ , of G of the average number of bicliques in κ to which each vertex belongs. This result was later reproved by Chung, Erdős and Spencer [3] in 1983 and Tuza [19] in 1984, among others.

Each of the above papers also included corresponding upper bounds for w(G), but it was not until Erds and Pyber [6] in 1997 that a similar upper bound for lbp, or lbc, was proved. More formally, they showed that for all n, $lbp(G) \leq \frac{c_2n}{\log n}$, for some constant c_2 . Combined with Lupanov's result above, this means that if lbc(G) is large compared to the number of vertices of G, then lbp(G) is not too much larger.

This is one motivation for seeking to bound lbp in terms of lbc, hoping for a similar result to our previously mentioned bound for bc in terms of bp. Perhaps surprisingly, however, we find that no such bound is possible. More specifically, in Section 4 we exhibit, for all $k \ge 2$, a graph G with lbc(G) = 2, but $lbp(G) \ge k$.

Given this, it is natural to ask 'how fast' can lbp(G) tend to infinity if lbc(G) = 2. One way of formalizing this question to ask 'if G has a 2-local *m*-cover, how large can lbp(G) be?' While we are able to answer this, up to a constant factor, in Section 5, we leave open the equally natural question of 'if lbc(G) = 2 and bc(G) = m, how large can lbp be?' A second question that we leave open is whether, loosely speaking, all graphs have a single cover that is close to optimal for both bc and lbc.

1.2 Subcube Intersection Graph Representations

The biclique cover number discussed above has a natural interpretation in the setting of hypercube intersection graphs.

The *d*-dimensional hypercube, Q_d , is the graph with vertex set $\{0, 1\}^d$ and with two vertices connected by an edge if and only if the vectors representing them differ in exactly one coordinate. A subset, S, of the vertices of the hypercube, Q_d , is called a subcube if it induces a graph isomorphic to a hypercube of some dimension. In other words, there is some set $J \subseteq [d] = \{1, 2, 3, \ldots, d\}$, and constants $a_j \in \{0, 1\}$ for each $j \in J$ such that $(x_1, \ldots, x_n) \in S$ if and only if for all $j \in J$, $x_j = a_j$. Fixed coordinates are those coordinates in J, whereas free coordinates are coordinates that are not fixed.

A subcube is said to have dimension r if it has r free coordinates. We shall write subcubes as vectors in $\{0, 1, *\}^d$, where the asterisks denote free coordinates and the 0's and 1's denote fixed coordinates, j, with $a_j = 0$ or 1 respectively.

An *intersection graph* of a family of subsets of some groundset has one vertex for each of the sets, and an edge between two vertices if and only if the corresponding sets intersect. See [10] or [15] for more on intersection graphs for various set families. In this paper, we are interested in intersection graphs where all sets in the family are subcubes of a hypercube. Such graphs are called *subcube intersection graphs*. Note that two subcubes intersect if and only if they agree on all coordinates where both are fixed. See Johnson and Markström [9] for some background.

Let I(n, d) be the set of all graphs on n vertices that are the intersection graph of some family of subcubes of a d-dimensional hypercube. It is easy to see that $G \in I(n, d)$ if and only if its complement has a d-cover, as was apparently first pointed out by Fishburn and Hammer [7], and also noted in [9]. Indeed, a representation of $G \in I(n, d)$ assigns a vector in $\{0, 1, *\}^d$ to each vertex. Generate a biclique for each of the d coordinates, having as one class all the vertices with a zero in this coordinate, and as the other class all vertices with a one in this coordinate. Then, the union of these bicliques has an edge exactly where G does not have an edge. This produces a d-cover of G^c and the converse is similar.

If G is a graph on n vertices, we write $\rho(G)$ for the smallest d such that $G \in I(n, d)$ —in

other words, $\rho(G)$ is the smallest d such that G may be represented as an intersection graph of a family of subcubes of Q_d . By the previous paragraph, $\rho(G) = \operatorname{bc}(G^c)$. We also write $\rho(n) = \max\{\rho(G) : |G| = n\}$, where |G| denotes the number of vertices of G.

Similarly, we define $\tau(G)$ to be the smallest r such that G may be represented as an intersection graph of a family of r dimensional subcubes of some hypercube. We also let $\tau(n) = \max\{\tau(G) : |G| = n\}$. Using the relationship between subcube intersection graphs and biclique coverings, we may see that $\tau(G)$ is the least r such that G^c has a biclique cover where every vertex lies in all but r of the bicliques in the cover. (Note in this interpretation, bicliques with one empty class are allowed). This is bears a resemblance to, but is distinct from $\operatorname{lbc}(G^c)$ —the least r such that G^c has a cover with every biclique in no more than r bicliques of the cover.

Many authors have placed various bounds on ρ —see for instance [2], [5], [18] and [19]. The best known bounds are:

$$n - c \log n \leqslant \rho(n) \leqslant n - \lfloor \log n \rfloor + 1,$$

where c is some positive constant. The upper bound is due to Tuza [19] and the probabilistically proved lower bound to Rödl and Ruciński [18]. The question of obtaining similar bounds for $\tau(n)$ was posed by Johnson and Markström [9]. In Section 2, we shall prove a very close relationship between $\tau(n)$ and $\rho(n)$ and thus we obtain bounds on τ similar to those on ρ .

We remark briefly that the other biclique covering/partitioning measures may be interpreted in this way. For instance, $lbc(G^c)$ is the least r such that G can be represented as the intersection graph of a family of codimension r subcubes of some hypercube. Equally, $bp(G^c)$ is the least d such that G can be represented in I(n, d) such that the subcubes of non-adjacent vertices differ in only one coordinate.

2 Relationship between τ and ρ

We shall prove the upper bounds and lower bounds on τ separately.

Lemma 1. For any graph G, $\tau(G) \leq \rho(G)$, and hence $\tau(n) \leq \rho(n)$.

Proof. Let G be a graph on n vertices with $\rho(G) = d$. This means that there are subcubes A_1, \ldots, A_n of Q_d whose intersection graph is G. Let M (resp. m) be the maximum (resp. minimum) dimension of these subcubes. For all *i*, replace A_i by A'_i , a subcube of Q_{d+M-m} by appending 0's and *'s to the vector of A_i . We ensure that we add as many *'s as needed so the resultant vector has precisely M *'s and as many 0's as needed to give the vector length (d+M-m). The intersection graph of the A'_i is G. This shows that G may be represented as an intersection graph of a family of subcubes of dimension M. Since $M \leq d$, the result follows.

Lemma 2. $\tau(n+1) \ge \rho(n)$.

Proof. Let G be a graph with vertices v_1, \ldots, v_n and with $\rho(G) = \rho(n)$. Now form G' from G by adding a single vertex, v_{n+1} , adjacent to all vertices of G. Let $A_1, \ldots, A_n, A_{n+1}$ be r-dimensional subcubes of Q_d such that G' is the intersection graph of the A_i (the vertex v_i being represented by A_i) and such that r is equal to $\tau(G')$. Without loss of generality, let A_{n+1} be free in the first r coordinates, and hence fixed in the other d - rcoordinates. Where A_{n+1} is fixed, all the other A_i must be either free (have an asterisk in that coordinate) or have the same fixed value as A_{n+1} , as v_{n+1} is adjacent to all the other vertices. Thus restricting the first n subcubes to the first r coordinates does not change which pairs of subcubes intersect. So the intersection graph of these restricted subcubes is G, implying that $G \in I(n, r)$.

Therefore, using the definition of ρ , $\tau(n+1) \ge \tau(G') = r \ge \rho(G) = \rho(n)$.

Combining these two lemmas with the bounds on ρ gives:

Theorem 3. There is some absolute constant, c, such that for all n,

$$n - c \log n \leqslant \tau(n) \leqslant n - \lfloor \log n \rfloor + 1.$$

3 Bounding bp **in terms of** bc

As the earlier example of K_n shows, bp can be as large as exponential in bc. Here, we show that bp grows no faster than exponentially in bc, and calculate the best upper bound exactly. We do this by reducing the problem to proving the upper bound for a single graph. In Theorem 6, we use ideas from Tverberg's proof [20] of the Graham-Pollak Theorem to calculate bp of this graph and show that the previous bound is tight.

Theorem 4. If bc(G) = m, then $bp(G) \leq \frac{1}{2}(3^m - 1)$.

Proof. Fix an *m*-cover of G, $\kappa = \{B_1, \ldots, B_m\}$, and for each biclique B_i we label the vertex classes as class 0 and class 1. We now represent each vertex of G as an element of $\{0, 1, *\}^m$, based on its membership of elements of κ . If $v \in V(G)$, define $\tilde{v} \in \{0, 1, *\}^m$ by:

$$\tilde{v}_i = \begin{cases} 0 & \text{if } v \text{ is in class } 0 \text{ of } B_i, \\ 1 & \text{if } v \text{ is in class } 1 \text{ of } B_i, \\ * & \text{otherwise.} \end{cases}$$

If $\tilde{u} = \tilde{v}$ then u and v have the same neighbours. Therefore, if $u \neq v$, this would mean that bp(G) = bp(G-u), so we could replace G with G-u and thus we may assume vertices all have different representations. Using the identification $v \mapsto \tilde{v}$, $V(G) \subseteq \{0, 1, *\}^m$.

Since the biclique partition number of a graph is at least that of each of its induced subgraphs, we may assume that $V(G) = \{0, 1, *\}^m$. Note that there is an edge between two vertices u and v if and only if there is some i for which $\{u_i, v_i\} = \{0, 1\}$. In other words, this is the complement of the intersection graph of all subcubes of an m-dimensional cube. We write G_m for this graph.

Before proceeding by induction on m, we introduce some notation. We write (α, β) for the biclique with vertex classes α and β . If $\alpha \subseteq \{0, 1\}^m$ and $x, y \in \{0, 1, *\}$, we define the following subsets of $\{0, 1, *\}^{m+1}$:

$$x\alpha := \{(x, v_1, v_2, \dots, v_m) : (v_1, \dots, v_m) \in \alpha\} \quad \stackrel{x}{y}\alpha := x\alpha \cup y\alpha$$

Let π_m be an optimal partition of G_m —i.e. π_m is a bc(G)-cover of G. Add B_1 to π_{m+1} . For every $(\alpha, \beta) \in \pi_m$, place the following three bicliques in π_{m+1} : $({}^*_0\alpha, {}^*_0\beta)$, $({}^*_1\alpha, 1\beta)$ and $(1\alpha, *\beta)$.

Claim 5. π_{m+1} is a partition of G_{m+1} and contains $3bp(G_m) + 1$ bicliques.

Proof of claim. Let $u = (u_1, \ldots, u_{m+1})$ and $v = (v_1, \ldots, v_{m+1})$ be two vertices that are not joined by an edge of G_{m+1} . This means there is no *i* such that $\{u_i, v_i\} = \{0, 1\}$ so uvis not an edge of B_1 . Moreover, if we define $u' = \{u_2, \ldots, u_{m+1}\}$ and $v' = \{v_2, \ldots, v_{m+1}\}$, i.e. *u* and *v* with the first coordinate removed, then u'v' is not an edge of G_m . Hence u'v'is not an edge in any biclique of π_m ; this implies uv is not an edge of any biclique of π_{m+1} .

Conversely, let e be an edge of G_{m+1} . If e is also an edge of B_1 , it is contained in no other bicliques of π_{m+1} . If e is not an edge of B_1 , then removing the first coordinate from each of its endpoints forms an edge e' of G_m . As π_m is a partition of G_m , e' lies in exactly one $(\alpha, \beta) \in \pi_m$. Then, by inspection, e must lie in exactly one of the corresponding bicliques in π_{m+1} .

As $bp(G_1) = 1$, our proof is concluded by applying the above claim inductively.

Theorem 6. There is a graph G with bc(G) = m, but $bp(G) = \frac{1}{2}(3^m - 1)$.

Proof. For a graph G, we write A(G) for the adjacency matrix of the graph, and rank(G) for the rank of A(G). Suppose P_1, \ldots, P_k form a biclique partition of G. Then $A(G) = \sum_{i=1}^k A(P_i)$. Since rank is subadditive and bicliques have rank 2, rank $(G) \leq 2k$. This implies $bp(G) \geq rank(G)/2$.

We conclude the proof by calculating rank (G_m) , where G_m is the same graph as defined above. We define the following order: 0 < 1 < * and use lexicographical order for the product spaces $\{0, 1, *\}^m$, for all m.

Note that $A_m = A(G_m)$ has 3^m rows (and columns) and the final row (and column) in the above order consists solely of 0's. We shall show by induction on m that the first $3^m - 1$ rows of $A_m = A(G_m)$, together with the all 1 vector, form a linearly independent set. This is trivial for m = 1.

In block matrix notation, and using the convention that $A_0 = (0)$,

$$A_{m} = \begin{pmatrix} A_{m-1} & 1 & A_{m-1} \\ 1 & A_{m-1} & A_{m-1} \\ A_{m-1} & A_{m-1} & A_{m-1} \end{pmatrix} \rightarrow \begin{pmatrix} 0 & 1 - A_{m-1} & 0 \\ 1 - A_{m-1} & 0 & 0 \\ A_{m-1} & A_{m-1} & A_{m-1} \end{pmatrix}$$

The second matrix is obtained from the first by simple row operations. We define A'_m as the matrix formed by replacing the final row of this second matrix by a row of only 1's.

$$A'_{m} = \begin{pmatrix} 0 & 1 - A_{m-1} & 0\\ 1 - A_{m-1} & 0 & 0\\ A'_{m-1} & A'_{m-1} & A'_{m-1} \end{pmatrix}$$

Then our induction hypothesis is equivalent to the claim that A'_m is of full rank in other words the linear span of the rows of A'_m , $\operatorname{span}(A'_m)$, is \mathbb{R}^{3^m} . Since $\operatorname{span}(1 - A_{m-1}) = \operatorname{span}(A'_{m-1}) = \mathbb{R}^{3^{m-1}}$, $\operatorname{span}(A'_m) = \mathbb{R}^m$. Therefore, $\operatorname{rank}(A_m) \ge 3^m - 1$ and $\operatorname{bp}(G_m) \ge \frac{1}{2}(3^m - 1)$. Combined with Theorem 4, this concludes the proof.

4 Relationship between lbc and lbp

As before, we write (α, β) for a biclique with vertex classes α and β .

In a contrast to the previous section's results, we show that lbp cannot be bounded by lbc, indeed, even for lbp(G) = 2, it is possible for lbc(G) to be arbitrarily large. More concretely:

Theorem 7. For all $m \ge 2$, there is a graph G with a 2-local m-cover for which $lbp(G) \ge \frac{1}{2} log(\frac{m-1}{3})$.

We explicitly construct graphs satisfying this inequality. For $m \ge 2$, let G_m have vertex set

$$V(G_m) = \{ v \in \{0, 1, *\}^m | v \text{ has exactly } m - 2 *'s \}.$$

We let $uv \in E(G_m)$ if and only if there is some *i* for which $\{u_i, v_i\} = \{0, 1\}$, where $u = \{u_1, \ldots, u_m\}$ and $v = \{v_1, \ldots, v_m\}$. Although we do not use this fact, G_m is the complement of the intersection graph of all codimension 2 subcubes of a cube of dimension *m*. Our definition of G_m may appear strange but it arises from simple reverse engineering of a graph with a 2-local *m*-cover. See the proof of Theorem 10 for more details.

The graph G_m may also be viewed as the union of m bicliques, B_1, \ldots, B_m , where B_i is the subgraph induced by the vertex set $\{v \in V(G) | v_i \neq *\}$. These B_i are known henceforth as *covering bicliques*, as they form a 2-local cover, although not partition, of G_m . Thus $\operatorname{lbc}(G_m) = 2$.

A crown graph on 2t vertices, H_t , is $K_{t,t}$ with a perfect matching removed. We may label the vertices of H_t as u^1, \ldots, u^t and v^1, \ldots, v^t ; $u^i v^j$ is an edge if and only if $i \neq j$, and there are no other edges.

We shall show that an r-local partition of G_m can be altered to form an r-local partition of H_t , for t linear in m. We then show that the local biclique cover number of crown graphs tends to infinity as the size of the graph does, indeed, $lbc(H_t) = \Omega(\log t)$. Since $lbp \ge lbc$, this suffices to finish the proof. Proof of Theorem 7. Let G_m and its subgraphs B_1, \ldots, B_m be as above. We say that an edge uv is *shared* if it is an edge in two of the covering bicliques; equivalently there are two *i* such that $\{u_i, v_i\} = \{0, 1\}$.

Let π be a partition of G—we call bicliques in π partitioning bicliques. Let e be a shared edge contained in B_i and let $A = (\alpha, \beta)$ be the partitioning biclique containing it. Consider the graph $A \cap B_i$ —i.e. the graph with vertex set $V(A) \cap V(B_i)$, and edge set $E(A) \cap E(B_i)$. It consists of up to two disjoint bicliques—let $A_i^e = (\alpha_i^e, \beta_i^e)$ denote the one containing the edge e. We term this the restriction of A onto B_i (with respect to the edge e). Note that all vertices in α_i^e agree in the i^{th} coordinate, as do all vertices in β_i^e .

Lemma 8. Let e be a shared edge of B_i and B_j , and $A = (\alpha, \beta)$ a biclique containing it. Then one of A_i^e or A_j^e is a star, or both A_i^e and A_j^e are $K_{2,2}$.

This lemma, though slightly technical, is key to our approach. A more intuitive formulation is that all bicliques containing an edge shared by two covering bicliques are either 'close' to being contained in one of these covering bicliques or 'very small'.

Proof. By permuting the coordinates and swapping 0 with 1 in the first two coordinates, we may assume that, the endpoints of e are (0, 0, *, ..., *) and (1, 1, *, ..., *), with the former being in α and the latter in β .

Notice that if any of α_1^e , α_2^e , β_1^e or β_2^e are singletons, we are done, so we assume the contrary. Suppose at least one of $(0, 1, *, \ldots, *)$ and $(1, 0, *, \ldots, *)$ is in A—without loss of generality $(0, 1, *, \ldots, *)$ is in α . Then every vertex in β must have a 1 in the first coordinate. But any vertex in β_2^e has 1 in the second coordinate. Thus there is only one vertex in β_2^e , contradicting our earlier assumption.

Therefore neither (0, 1, *, ..., *) nor (1, 0, *, ..., *) is in A. A second vertex in α_1^e must have a 0 in the first coordinate, a * in the second coordinate and without loss of generality, a 0 in the third coordinate. Equally, a second vertex in β_2^e must be (*, 1, 1, *, ..., *) in order to be adjacent to the vertices we have assumed are in α .

A further vertex in α_2^e must start with *, followed by 0 and have a further non-asterisk digit, in the i^{th} place, say, (i > 2). Our previous assumptions determine that there is only one further vertex in A—the vector with 1's in the first and i^{th} coordinates. It can be seen that A_1 and A_2 are both $K_{2,2}$.

We now define for each i, a colouring c_i of the shared edges of B_i . If e is a shared edge of B_i and B_j and A is the partitioning biclique containing it, $c_i(e)$ is red if A_i^e is a star or a $K_{2,2}$ and blue otherwise. Note that the preceding lemma shows that if $c_i(e)$ is blue then A_i^e is a star.

A shared edge is blue in at most one colouring. Hence, each of the $2\binom{m}{2} = m(m-1)$ shared edges is red in at least one of the *m* colourings, implying that at least one of the *m* covering bicliques, B_1 say, must contain at least m-1 shared edges coloured red.

All shared edges are vertex-disjoint—as seen by the form of the edges as vectors—so we may label the edges as $u^1v^1, \ldots, u^{m-1}v^{m-1}$, where the u^i and the v^j are from different classes of B_1 . We seek a large $B \subseteq B_1$ such that all shared edges in B are in partitioning

bicliques that restrict to stars on B_1 . Pick $I \subseteq [m-1]$ as follows. Let $1 \in I$. If the shared edge $e = u^1 v^1 \in A \in \pi$ and $A_1^e \cong K_{2,2}$ then discard from I the (at most two) indices other than 1 for which one of u^i and v^i is a vertex of this $K_{2,2}$. Then proceed to the next index not already discarded, and continue in the same way. Since at each stage, we throw away at most two indices, the set of surviving indices, I, satisfies $|I| \ge (m-1)/3$.

Let *B* be the induced subgraph of B_1 with the u^i and v^j as vertices, for $i \in I$, $j \in I$. Let *H* be the subgraph of *B* formed by removing all shared edges, $u^i v^i$ for $i \in I$. Clearly, *H* is isomorphic to $H_{|I|}$, the crown graph on 2|I| vertices. Let *A* be a biclique containing one of these removed edges, *e*, and note that $A \cap B$ is the union of at most two disjoint bicliques. By the construction of *B*, the component of $A \cap B$ containing *e*, $A_1^e \cap B$, is a star. We label the other (possibly empty) biclique component *A'*. From π , we induce a partition π' of *H* as follows. If $A \in \pi$ contains an edge of *H*, then place $A \cap H$ inside π' . Now, $A \cap H = A \cap B$ and is hence a union of bicliques, unless *A* contains a removed edge. Using G - e to denote the graph *G* with the edge *e* deleted, in this case, $A \cap H = A' \cup (A_1^e \cap B - e)$, the latter component remaining a star. Note that if π is a *k*-local partition of *G*_m, then π' is a *k*-local partition of *H*. This argument is valid for all partitions π of *G*, so $lbp(H) \leq lbp(G_m)$.

We now complete the proof in the following lemma:

Lemma 9. $\operatorname{lbc}(H_t) \ge \frac{1}{2} \lceil \log t \rceil$.

Proof. A biclique cover of H_t induces a biclique cover of K_t when two corresponding vertices (i.e. u^i and v^i for each i) are identified as one. As mentioned earlier, $lbc(K_t) = \lceil \log t \rceil$ so, there must be some vertex of the K_t in at least $\lceil \log t \rceil$ bicliques belonging to the induced cover. This implies one of the vertices of H_t that are identified to it is in at least $\lfloor \frac{1}{2} \lceil \log t \rceil$ bicliques, in the original covering.

5 Bounding lbp(G), for lbc(G) = 2

The previous section showed that even for lbc(G) = 2, we cannot place an absolute bound on lbp(G). In this section, we instead aim for bounds involving additional covering information about G. One way of doing this is by looking at the number of bicliques in a 2-local cover of G:

Theorem 10. If G has a 2-local m-cover then $lbp(G) \leq 2\lceil log(m-1) \rceil + 2$.

This upper bound is best possible up to a constant factor by Theorem 7. We proceed by using an argument similar to the start of the proof of Theorem 4 to reduce proving the upper bound for all graphs with a 2-local m-cover to proving it for one particular graph. This 'worst' graph will turn out to be similar to the graph used in the proof of Theorem 7. *Proof.* Fix a 2-local *m*-cover of *G*. We may identify vertices $v \in V(G)$ with vectors $\tilde{v} \in \{0, 1, *\}^m$ with exactly m - 2 s's, similar to the proof of Theorem 4. As shown there, we may assume no two vertices have the same representation, and that all such vectors are the representation of some vertex, i.e. we assume that

$$V(G) \subseteq \{v \in \{0, 1, *\}^m | v \text{ has at least } m - 2 *'s\}.$$

Note that there is an edge between two vertices u and v if and only if there is some i for which $\{u_i, v_i\} = \{0, 1\}$.

For i = 1, ..., m, let B_i be the induced subgraph on the vertices whose i^{th} coordinate is not *. We define shared edges in an identical manner to in Section 4 and we shall define C_i , a subgraph of B_i , such that every edge of G is an edge in exactly one C_i .

Let C_i contain all non-shared edges of B_i . Additionally, for all i and j, let C_i contain exactly one of the edges shared between B_i and B_j , and let C_j contain the other. Thus any vertex is in at most two C_i , and the C_i are isomorphic, so $lbp(G) \leq 2lbp(C_1)$. But C_i is a complete bipartite graph minus a matching of size m - 1. It is easy to show that $lbp(C_1) \leq lbp(H_{m-1}) + 1$, where H_{m-1} is the crown graph on 2(m-1) vertices.

We now show that $lbp(H_t) \leq \lceil \log t \rceil$, for any t. We name the classes of $H_t U$ and V. We assign to each vertex a unique label consisting of its class name and a binary vector of length $\lceil \log t \rceil$. For each position, $i = 1, \ldots, \lceil \log t \rceil$, we generate two bicliques corresponding to the difference in the i^{th} position of the binary vector. More specifically, we let B_i^1 be the induced biclique formed by all the vertices of class U with a 0 in the i^{th} coordinate and all vertices of class V with a 1 in the i^{th} coordinate. Similarly, we let B_i^2 be the induced biclique formed by all vertices of class U with a 1 in the i^{th} coordinate and all vertices of class V with a 1 in the i^{th} coordinate and all vertices of class V with a 0 in the i^{th} coordinate. The collection of all these bicliques is a $\lceil \log t \rceil$ -local partition of H_t . This concludes the proof.

Since bc(G) is a measure of independent interest to local measures, one may try to use the separate existence of a *m*-cover in a graph with lbc(G) = 2, rather than the existence of a 2-local *m*-cover, to bound lbp(G). More formally, we ask:

Question 11. What is the smallest k(m) such that if the graph G has bc(G) = m and lbc(G) = 2, we have $lbp(G) \leq k$?

Theorem 7 tells us that $k(m) \ge \frac{1+o(1)}{2} \log m$. However, an upper bound does not follow directly from Theorem 10 as the cover that shows lbc(G) = 2 may be genuinely different from the cover that shows bc(G) = m. The following result shows this:

Theorem 12. For all $m \ge 4$, there is a bipartite graph, G, with lbp(G) = lbc(G) = 2and bp(G) = bc(G) = m, but there are no (m-1)-local m-covers.

Proof. First, we make a general observation about biclique covers. Two edges, $\{v_1, v_2\}$ and $\{u_1, u_2\} \in E(G)$ are called *strongly independent* if they are independent and the minimum degree of $G[u_1, u_2, v_1, v_2]$ is 1—in other words, they are vertex disjoint and do

not lie in the same $K_{2,2}$ subgraph. No biclique can contain both edges so if E(G) contains a set of k pairwise strongly independent edges, $bc(G) \ge k$.

Let $X = \{x_1, \ldots, x_m\}$, $Y = \{y_{i,j} : 1 \le i < j \le m\} \cup \{y_{[m]}\}$ and let $V(G) = X \cup Y$. Let there be an edge from x_i to $y_{i,j}$ and to $y_{j,i}$ for all appropriate j. Also, let $y_{[m]}$ have edges to all of the x_i . This defines a bipartite graph, G. Let the partition π consist of a $K_{1,m}$ with its centre at $y_{[m]}$, and stars at each of the x_i each with leaves at all adjacent vertices other than $y_{[m]}$. Thus lbp(G) = lbc(G) = 2. Note that $I = \{\{x_1, y_{1,2}\}, \{x_2, y_{2,3}\}, \ldots, \{x_m, y_{m,1}\}\}$ is a set of strongly independent edges. Therefore $bc(G) \ge m$, and so bp(G) = bc(G) = m as we can cover (indeed partition) with stars that have their centre in X.

Suppose κ is a biclique cover of G using just m bicliques. Each of these contains exactly one of the strongly independent edges listed above.

We can see that $I \cup \{x_1, y_{1,3}\} \setminus \{x_1, y_{1,2}\}$ is also a strongly independent set of m edges. Therefore, x_1 must be at the centre of a star in κ . Similarly, every x_i must be the centre of a star in κ , and these m stars account for all the bicliques in κ . So the star with centre x_i , must contain all vertices adjacent to x_i . Thus $y_{[m]}$ is in m bicliques of κ .

So we can see that insisting that a cover of a graph G attains bc(G) may ensure the cover is not close to being optimal for lbc. This leads to the following question—if we insist that the cover of a graph attains lbc(G), could this ensure that it is far from being optimal for bc? More formally:

Question 13. Suppose bc(G) = m, lbc(G) = k, what is the smallest r = r(m, k) such that we can guarantee G has a k-local r-cover?

An answer to this for k = 2, combined with Theorem 10, would lead to a partial solution of Question 11.

The following example of the hypercube is instructive, in giving bounds on the problem. In particular, it shows $r(2^{d-1}, d/2) \ge d2^{d-3}$, for d even.

Example 14. Consider Q_d , for d even. It is easy to show that $bp(Q_d) = bc(Q_d) = 2^{d-1}$. Indeed, each biclique contains at most d edges, and $e(Q_d) = d2^{d-1}$, so $bc(Q_d) \ge 2^{d-1}$. On the other hand, we may partition using $|Q_d|/2 = 2^{d-1}$ stars since Q_d is bipartite. Dong and Liu [4] showed that $lbp(Q_d) = lbc(Q_d) = d/2$. In fact, their proof showed that a cover achieves this only if it consists solely of $K_{2,2}$'s so every d/2-local cover contains at least $d2^{d-3}$ bicliques.

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