# Lower Bounds for Cover-Free Families

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#### Abstract

Let  $\mathcal{F}$  be a set of blocks of a *t*-set X. A pair  $(X, \mathcal{F})$  is called an (w, r)-cover-free family ((w, r)-CFF) provided that, the intersection of any w blocks in  $\mathcal{F}$  is not contained in the union of any other r blocks in  $\mathcal{F}$ .

We give new asymptotic lower bounds for the number of minimum points t in a (w, r)-CFF when  $w \leq r = |\mathcal{F}|^{\epsilon}$  for some constant  $\epsilon \geq 1/2$ .

Keywords: Cover-Free Family, Lower Bound.

## 1 Introduction

Let  $\mathcal{F}$  be a set of blocks (subsets) of a *t*-set X. A pair  $(X, \mathcal{F})$  is called a (w, r)-cover-free family ((w, r) - CFF) provided that, for any w blocks  $A_1, A_2, \ldots, A_w \in \mathcal{F}$  and any other r blocks  $B_1, B_2, \ldots, B_r \in \mathcal{F}$  we have

$$\bigcap_{i=1}^{w} A_i \not\subseteq \bigcup_{j=1}^{r} B_j.$$

Since using De Morgan, a (w, r)-CFF can be turned into (r, w)-CFF, throughout the paper we assume that  $w \leq r$ . Cover-free families were first introduced in 1964 by Kautz and Singleton [5].

Let N(n, (w, r)) denote the minimum number of points |X| in any (w, r)-CFF having  $|\mathcal{F}| = n$  blocks. The best known lower bound for N(n, (1, r)) is [2, 4, 7]

$$N(n,(1,r)) = \Omega\left(\frac{r^2}{\log r}\log n\right) \tag{1}$$

when  $r \leq \sqrt{n}$ , and,  $\Omega(n)$  when  $r > \sqrt{n}$ . The constant of the  $\Omega()$  is asymptotically 1/2, 1/4 and 1/8, respectively. Stinson et. al, [8], proved that

$$N(n, (w, r)) \ge N(n - 1, (w - 1, r)) + N(n - 1, (w, r - 1)).$$
(2)

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They then use it with (1) to prove two bounds. The first bound is

$$N(n, (w, r)) \ge \Omega\left(\frac{\binom{w+r}{w}(w+r)}{\log\binom{w+r}{w}}\log n\right)$$
(3)

when  $r \leqslant \sqrt{n}$ , [8, 6], and

$$N(n, (w, r)) \ge \Omega\left(\frac{\binom{w+r}{w}}{\log(w+r)}\log n\right)$$
(4)

for any  $r \leq n$ , [8]. To the best of our knowledge (4) is the best bound known when  $\sqrt{n} \leq r \leq n$ . D'yachkov et. al. breakthrough result, [3], implies that for  $r \leq \sqrt{n}$  and  $r, n \to \infty$ 

$$N(n, (w, r)) = \Theta\left(\frac{\binom{w+r}{w}(w+r)}{\log\binom{w+r}{w}}\log n\right)$$
(5)

and for  $r \ge \sqrt{n}$  and  $r, n \to \infty$ 

$$N(n, (w, r)) \leqslant O\left(\frac{r}{w} \cdot \frac{\binom{w+r}{w}}{\log(w+r)}\log n\right).$$
(6)

In this paper we give a new lower bound for (w, r)-CFF when  $r > \sqrt{n}$ . We combine the two techniques used in [8, 6] and [1] to give the following asymptotic lower bound.

**Theorem 1.** For any  $2 \leq k \leq w < r \leq n/2$  we have

$$N(n, (w, r)) \ge \frac{k^k k!}{2(k+1)^{2k}} \frac{r^{w+1}}{(w+1)! \ln^k r} = \Omega\left(\frac{\sqrt{k}}{e^k} \cdot \frac{r^{w+1}}{(w+1)! \ln^{k+1} r} \log n\right)$$

for

$$(n+k-1-w)^{\frac{k-1}{k}} \leqslant r \leqslant (n+k-w)^{\frac{k}{k+1}}$$

and

$$N(n, (w, r)) = \Theta\left(\binom{n}{w}\right)$$

for

$$r = \Omega\left((n\log n)^{\frac{w}{w+1}}\right).$$

Our bound is

$$\Theta\left(\frac{\sqrt{k}\cdot r}{w(e\ln r)^k}\right)$$

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times greater than the previous bound in (4). In particular, when k is constant, our lower bound improves the bound in (4) to

$$N(n, (w, r)) \ge \Omega\left(\frac{r}{w\log^k r} \cdot \frac{\binom{w+r}{w}}{\log(w+r)}\log n\right).$$
(7)

A slightly better bound can be achieved when

$$(n+k-w)^{\frac{k}{k+1}} \leqslant r \leqslant (n+k-w)^{\frac{k}{k+1}} \ln^{1/(k+1)} n.$$

For example, let w = 4. Table 1 compares our results with the previous results (asymptotic values).

	Previous Lower	Upper	Our Lower
<i>r</i>	Bounds $(3), (4)$	Bound [3]	Bound
$r\leqslant n^{1/2}$	$r^{5} \frac{\log n}{\log r}$	$r^5 \frac{\log n}{\log r}$	
$n^{1/2}\leqslant r\leqslant n^{2/3}$	$r^4 \frac{\log n}{\log r}$	$r^{5} \frac{\log n}{\log r}$	$r^5 \frac{\log n}{\log^3 r}$
$n^{2/3} \leqslant r \leqslant n^{3/4}$	$r^4 \frac{\log n}{\log r}$	$r^{5} rac{\log n}{\log r}$	$r^5 \frac{\log n}{\log^4 r}$
$n^{3/4}\leqslant r\leqslant n^{4/5}$	$r^4 \frac{\log n}{\log r}$	$r^{5} \frac{\log n}{\log r}$	$r^5 \frac{\log n}{\log^5 r}$
$n > r \ge (n \log n)^{4/5}$	$r^4$	$n^4$	$n^4$

Table 1: Results for w = 4.

## 2 First Lower Bound

In this section, we prove

Lemma 2. Let  $w \leq r \leq n/2$ . If

$$r = \Omega\left((n\log n)^{\frac{w}{w+1}}\right)$$

then

$$N(n, (w, r)) = \Theta\left(\binom{n}{w}\right).$$
(8)

Otherwise,

$$N(n, (w, r)) \ge \Omega\left(\left(\frac{r}{(w+1)\ln r}\right)^{w+1}\log n\right).$$
(9)

Lemma 2 follows from the following.

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**Lemma 3.** Let  $\epsilon < 1$  be any constant. For  $w \leq r \leq n/2$  we have

$$N(n,(w,r)) \ge \min\left((1-\epsilon)\frac{w^w}{(w+1)^{2w+1}} \cdot \frac{r^{w+1}}{\ln^w r}, \ \epsilon\binom{n}{w}\right).$$
(10)

*Proof.* Let  $(X, \mathcal{F})$  be an optimal (w, r)-CFF. Let  $\mathcal{F} = \{F_1, \ldots, F_n\},\$ 

$$|X| = N = N(n, (w, r))$$

and assume without loss of generality that  $X = [N] := \{1, \ldots, N\}$ . Define  $v^{(i)} \in \{0, 1\}^n$ ,  $i = 1, \ldots, N$  where  $v_j^{(i)} = 1$  if and only if  $i \in F_j$ . Let  $V = \{v^{(i)} | i = 1, \ldots, N\}$ . Let  $V_0$  be the set of  $v^{(i)}$  of weight  $wt(v^{(i)})$  (i.e.,  $\sum_j v_j^{(i)}$ ) equal to w. Let

$$m = \frac{(w+1)^2 n \ln r}{wr}$$

and consider the two sets  $V_1 = \{v^{(i)} \mid w < wt(v^{(i)}) < m\}$  and  $V_2 = \{v^{(i)} \mid wt(v^{(i)}) \ge m\}$ . Obviously,  $V = V_0 \cup V_1 \cup V_2$  is a partition of V. Suppose by contradiction that

$$|V_0| \leqslant \epsilon \binom{n}{w}$$

and

$$\max(|V_1|, |V_2|) \le (1 - \epsilon) \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{r^{w+1}}{\ln^w r}$$

Consider  $W = \{(j_1, \ldots, j_w) \mid 1 \leq j_1 < \cdots < j_w \leq n\}$  and  $W' \subset W$  the set of all  $(j_1, \ldots, j_w)$  where no  $v^{(i)} \in V_0, i = 1, \ldots, N$ , satisfies  $v_{j_1}^{(i)} = \cdots = v_{j_w}^{(i)} = 1$ . Obviously,

$$|W'| = \binom{n}{w} - |V_0| \ge (1-\epsilon)\binom{n}{w}.$$

Fix an element  $v \in V_1$  and randomly and uniformly choose  $j = (j_1, \ldots, j_w) \in W'$ . We have

$$\Pr_{j \in W'}[v_{j_1} = \dots = v_{j_w} = 1] \leqslant \frac{\binom{wt(v)}{w}}{|W'|} \leqslant \frac{\binom{m}{w}}{(1-\epsilon)\binom{n}{w}}$$

Therefore, the expectation of the number of  $v \in V_1$  for which  $v_{j_1} = \cdots = v_{j_w} = 1$  is at most

$$\frac{\binom{m}{w}|V_1|}{(1-\epsilon)\binom{n}{w}} \leq \frac{1}{1-\epsilon} \left(\frac{m}{n}\right)^w |V_1|$$
$$\leq \frac{1}{1-\epsilon} \frac{(w+1)^{2w} \ln^w r}{w^w r^w} \cdot (1-\epsilon) \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{r^{w+1}}{\ln^w r}$$
$$= \frac{r}{w+1}.$$

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Therefore, there is  $j' = (j'_1, \ldots, j'_w) \in W'$  such that the number of  $v \in V_1$  that satisfies  $v_{j'_1} = \cdots = v_{j'_w} = 1$  is  $r_1 \leq r/(w+1)$ . Since the weight of every  $v \in V_1$  is greater than w, we can choose  $r_1$  new entries  $j''_1, \ldots, j''_{r_1} \notin \{j'_1, \ldots, j'_w\}$  such that for every  $v \in V_1$  where  $v_{j'_1} = \cdots = v_{j'_w} = 1$  there is  $j''_\ell$  such that  $v_{j''_\ell} = 1$ .

Now randomly and uniformly choose

$$r_2 := \left\lceil \frac{wr}{w+1} \right\rceil$$

distinct  $k_1, \ldots, k_{r_2} \in [n]$ . Let A be the event that  $\{k_1, \ldots, k_{r_2}\} \cap \{j'_1, \ldots, j'_w\} \neq \emptyset$ . The probability that A does not happen is

$$\frac{\binom{n-w}{r_2}}{\binom{n}{r_2}} \ge \frac{\binom{n-w}{r_2}}{2^w\binom{n-w}{r_2}} = \frac{1}{2^w}.$$

Then

$$\Pr[A \lor (\exists v \in V_2) \ v_{k_1} = \dots = v_{k_{r_2}} = 0] \leqslant 1 - \frac{1}{2^w} + |V_2| \frac{\binom{n-m}{r_2}}{\binom{n}{r_2}}$$
$$\leqslant 1 - \frac{1}{2^w} + |V_2| \left(\frac{n-m}{n}\right)^{r_2}$$
$$\leqslant 1 - \frac{1}{2^w} + |V_2| e^{-\frac{mr_2}{n}}$$

and

$$\begin{aligned} V_2 | e^{-\frac{mr_2}{n}} &\leqslant (1-\epsilon) \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{r^{w+1}}{\ln^w r} \cdot e^{-\frac{(w+1)^2 \ln r}{wr} r_2} \\ &\leqslant (1-\epsilon) \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{r^{w+1}}{\ln^w r} \cdot e^{-(w+1)\ln r} \\ &= (1-\epsilon) \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{1}{\ln^w r} \\ &< \frac{1}{2^w}. \end{aligned}$$

Therefore,

$$\Pr[A \lor (\exists v \in V_2) \ v_{k_1} = \dots = v_{k_{r_2}} = 0] < 1.$$

Therefore, there is  $\{k_1, \ldots, k_{r_2}\}$  such that  $\{k_1, \ldots, k_{r_2}\} \cap \{j'_1, \ldots, j'_w\} = \emptyset$  and for every  $v \in V_2$  there is  $k_\ell \in \{k_1, \ldots, k_{r_2}\}$  where  $v_{k_\ell} = 1$ .

Now it is easy to see that there is no  $v \in V$  where  $v_{j'_1} = \cdots = v_{j'_w} = 1$ ,  $v_{j''_1} = \cdots = v_{j''_{r_1}} = 0$  and  $v_{k_1} = \cdots = v_{k_{r_2}} = 0$ . This implies that

$$\bigcap_{i=1}^{w} F_{j_i'} \subseteq \bigcup_{i=1}^{r_1} F_{j_i''} \cup \bigcup_{i=1}^{r_2} F_{k_i},$$

which is a contradiction.

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# 3 The Second Bound

In this section we prove Theorem 1.

**Lemma 4.** For any  $2 \leq k \leq w \leq r \leq n/2$  and

$$2 \leqslant r \leqslant (n+k-w)^{\frac{\kappa}{k+1}}$$
$$N(n,(w,r)) \geqslant \frac{k^k k!}{2(k+1)^{2k}} \frac{r^{w+1}}{(w+1)! \ln^k r} = \Omega\left(\frac{r^{w+1}}{(w+1)! \ln^k r}\right).$$

*Proof.* We prove the lemma by induction on w.

From Lemma 3 the lemma holds for w = k. Now assume the bound holds for some w and every r that satisfies  $r \leq (n + k - w)^{\frac{k}{k+1}}$ . We now prove the bound for w + 1 and  $r \leq (n + k - w - 1)^{\frac{k}{k+1}}$ . We have

$$N(n, (w+1, r)) \ge N(n-1, (w, r))) + N(n-1, (w+1, r-1))$$
(11)

$$\geq \sum_{j=1}^{N} N(n-r+j-1,(w,j))$$

$$\geq N(n-r,(w,1)) +$$

$$(12)$$

$$N(n-r,(w,1)) + \sum_{j=2}^{r} \frac{k^{k}k!}{2(k+1)^{2k}} \frac{j^{w+1}}{(w+1)! \ln^{k} j}$$
(13)

$$\geq \frac{k^{k}k!}{2(k+1)^{2k}(w+1)!\ln^{k}r} \sum_{j=1}^{r} j^{w+1} \\ \geq \frac{k^{k}k!}{2(k+1)^{2k}(w+1)!\ln^{k}r} \int_{0}^{r} x^{w+1} dx \\ \geq \frac{k^{k}k!}{2(k+1)^{2k}} \frac{r^{w+2}}{(w+2)!\ln^{k}r}.$$

Here, inequality (11) comes from [8]. Inequality (12) follows from the fact that  $N(n - r + 1, (w + 1, 1)) \ge N(n - r, (w, 1))$ . Inequality (13) follows from the induction hypothesis since

$$j = r - (r - j)$$

$$\leqslant (n + k - w - 1)^{\frac{k}{k+1}} - (r - j)$$

$$\leqslant (n + k - w - 1 - (r - j))^{\frac{k}{k+1}}$$

$$= ((n - r + j - 1) + k - w)^{\frac{k}{k+1}}.$$

Lemma 5. Let  $w \leq r \leq n/2$ . If

$$r = \Omega\left((n\log n)^{\frac{w}{w+1}}\right)$$

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then

$$N(n, (w, r)) = \Theta\left(\binom{n}{w}\right).$$

*Proof.* Let  $r > c(n \log n)^{\frac{w}{w+1}}$  for large enough constant c > 2e and  $c' = c^{(w+1)/w}$ . Since

$$\begin{aligned} \frac{1}{2} \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{r^{w+1}}{\ln^w r} & \geqslant \quad \frac{1}{2} \frac{w^w}{(w+1)^{2w+1}} \cdot \frac{c^{w+1} n^w \log^w n}{\left(\frac{w}{w+1}\right)^w \ln^w (c'n \log n)} \\ & \geqslant \quad \frac{1}{2} \frac{1}{(w+1)\left(1+\frac{1}{w}\right)^w} \cdot \frac{c^{w+1} n^w \ln^w n}{w^w \ln^w (n^2)} \\ & \geqslant \quad \frac{c}{2e} \frac{1}{(w+1)} \left(\frac{c}{2e}\right)^w \cdot \left(\frac{en}{w}\right)^w \\ & \geqslant \quad \frac{c}{2e} \frac{1}{(w+1)} \left(\frac{c}{2e}\right)^w \cdot \binom{n}{w} \geqslant \binom{n}{w}, \end{aligned}$$

by Lemma 3, the result follows.

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