

# On vertex-disjoint paths in regular graphs

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## Abstract

Let  $c \in (0, 1]$  be a real number and let  $n$  be a sufficiently large integer. We prove that every  $n$ -vertex  $cn$ -regular graph  $G$  contains a collection of  $\lfloor 1/c \rfloor$  paths whose union covers all but at most  $o(n)$  vertices of  $G$ . The constant  $\lfloor 1/c \rfloor$  is best possible when  $1/c \notin \mathbb{N}$  and off by 1 otherwise. Moreover, if in addition  $G$  is bipartite, then the number of paths can be reduced to  $\lfloor 1/(2c) \rfloor$ , which is best possible.

**Mathematics Subject Classifications:** 05C38

## 1 Introduction

Paths and cycles are fundamental objects in graph theory. The *path cover number* is the minimum number of vertex-disjoint paths whose union covers all vertices of  $G$ . Note that we allow paths of length 0 (single vertices) in the definition above. Trivially the path cover number of a graph  $G$  is upper bounded by the independence number of  $G$ , because the set of the (arbitrary one out of the two) end vertices of the paths in a minimal path cover form an independent set in  $G$ . It is evident that for general graphs, determining the path cover number is NP-hard, because deciding if the path cover number equals 1 is equivalent to the decision problem for a Hamiltonian path, which is NP-complete. For bounds on the path cover number for general graphs, see e.g. [3, 6]. For regular graphs, Magnant and Martin [5] made the following conjecture and confirmed it for  $0 \leq k \leq 5$ .

**Conjecture 1.** If  $G$  be a  $k$ -regular graph of order  $n$ , then the path cover number of  $G$  is at most  $n/(k+1)$ .

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If true, the bound in Conjecture 1 would be tight as seen by disjoint copies of complete graph  $K_{k+1}$  (if  $n \equiv j$  modulo  $k + 1$  and  $j \neq 0$ , then we change one copy of  $K_{k+1}$  to a  $k$ -regular subgraph of  $K_{k+1+j}$ ). By the celebrated Dirac theorem on Hamiltonian paths [2], Conjecture 1 is true for  $k \geq (n - 1)/2$ . As far as we know, Conjecture 1 is open for all other cases. To provide more evidence on the validity of the conjecture, in this note we prove the following result for dense regular graphs.

**Theorem 2.** *For any  $c, \alpha > 0$ , there exists  $n_0 \in \mathbb{N}$  such that the following holds for every integer  $n \geq n_0$ .*

1. *Every  $\lceil cn \rceil$ -regular graph of order  $n$  contains a collection of at most  $\lfloor 1/c \rfloor$  vertex-disjoint paths whose union covers all but  $\alpha n$  vertices.*
2. *Every bipartite  $\lceil cn \rceil$ -regular graph of order  $n$  contains a collection of at most  $\lfloor 1/(2c) \rfloor$  vertex-disjoint paths whose union covers all but  $\alpha n$  vertices.*

Note that Part (2) of the theorem corresponds to the bipartite version of Conjecture 1: if  $G$  is a bipartite  $k$ -regular graph of order  $n$ , then the path cover number of  $G$  can be as large as  $n/(2k)$ , as seen by the vertex-disjoint copies of  $K_{k,k}$ . Note that

$$\frac{-2}{c(\lceil cn \rceil + 1)} \leq \frac{n}{\lceil cn \rceil + 1} - \frac{1}{c} = \frac{cn - \lceil cn \rceil - 1}{c(\lceil cn \rceil + 1)} < 0.$$

So when  $n$  is large, if  $1/c \notin \mathbb{N}$ , then  $\lfloor \frac{n}{\lceil cn \rceil + 1} \rfloor = \lfloor 1/c \rfloor$ , i.e., the number of paths in Theorem 2 matches the quantity in Conjecture 1; however if  $1/c \in \mathbb{N}$ , then the quantity  $\lfloor 1/c \rfloor$  is off by 1. On the other hand, the quantity  $\lfloor 1/(2c) \rfloor$  in Part (2) is optimal.

At last, we remark that the bound in Conjecture 1 is not tight if we restrict the problem on connected regular graphs, see [7] for connected cubic graphs.

## 2 A weaker result

We first prove the following weaker result. For reals  $a, b, c$ , we write  $a = (1 \pm b)c$  if there exists a real  $x \in (1 - b, 1 + b)$  such that  $a = xc$ .

**Theorem 3.** *Given any reals  $c, \alpha > 0$ , there exists  $\epsilon > 0$  and integer  $C > 0$  such that the following holds for sufficiently large integer  $n$ . Let  $G$  be a graph of order  $n$  such that  $\deg(v) = (1 \pm \epsilon)cn$  for every  $v \in V(G)$ . Then there exists a collection of  $C$  vertex-disjoint cycles in  $G$  whose union covers all but  $\alpha n$  vertices of  $G$ .*

Our main tools for embedding the cycles are the Regularity Lemma of Szemerédi [11] and the Blow-up Lemma of Komlós et al. [4]. For any two disjoint vertex-sets  $A$  and  $B$  of a graph  $G$ , the density of  $A$  and  $B$  is defined as  $d(A, B) := e(A, B)/(|A||B|)$ , where  $e(A, B)$  is the number of edges with one end vertex in  $A$  and the other in  $B$ . Let  $\epsilon$  and  $\delta$  be two positive real numbers. The pair  $(A, B)$  is called  $\epsilon$ -regular if for every  $X \subseteq A$  and  $Y \subseteq B$  satisfying  $|X| > \epsilon|A|$ ,  $|Y| > \epsilon|B|$ , we have  $|d(X, Y) - d(A, B)| < \epsilon$ . Moreover, the pair  $(A, B)$  is called  $(\epsilon, \delta)$ -super-regular if  $(A, B)$  is  $\epsilon$ -regular and  $\deg_B(a) > \delta|B|$  for all  $a \in A$  and  $\deg_A(b) > \delta|A|$  for all  $b \in B$ .

**Lemma 4** (Regularity Lemma – Degree Form). *For every  $\epsilon > 0$  there is an  $M = M(\epsilon)$  such that for any graph  $G = (V, E)$  and any real number  $d \in [0, 1]$ , there is a partition of the vertex set  $V$  into  $t + 1$  clusters  $V_0, V_1, \dots, V_t$ , and there is a subgraph  $G'$  of  $G$  with the following properties:*

- $t \leq M$ ,
- $|V_i| \leq \epsilon|V|$  for  $0 \leq i \leq t$  and  $|V_1| = |V_2| = \dots = |V_t|$ ,
- $\deg_{G'}(v) > \deg_G(v) - (d + \epsilon)|V|$  for all  $v \in V$ ,
- $G'[V_i] = \emptyset$  for all  $i$ ,
- each  $(V_i, V_j)$ ,  $1 \leq i < j \leq t$ , is  $\epsilon$ -regular with  $d(V_i, V_j) = 0$  or  $d(V_i, V_j) \geq d$  in  $G'$ .

The Blow-up Lemma allows us to regard a super regular pair as a complete bipartite graph when embedding a graph with bounded degree. Since we will always use it to embed a cycle, we state it in the following special form.

**Lemma 5.** *For every  $\delta > 0$ , there exists an  $\epsilon > 0$  such that the following holds for sufficiently large integer  $N$ . Let  $(X, Y)$  be an  $(\epsilon, \delta)$ -super-regular pair with  $|X| = |Y| = N$ . Then  $(X, Y)$  contains a spanning cycle (a cycle of length  $2N$ ).*

A *fractional matching* is a function  $f$  that assigns to each edge of a graph a real number in  $[0, 1]$  so that, for each vertex  $v$ , we have  $\sum f(e) \leq 1$  where the sum is taken over all edges incident to  $v$ . The *fractional matching number*  $\mu_f(G)$  of a graph  $G$  is the supremum of  $\sum_{e \in E(G)} f(e)$  over all fractional matchings  $f$ . We use the following so-called ‘fractional Berge-Tutte formula’ of Scheinerman and Ullman [10, Theorem 2.2.6]. Note that it is also proved in [10] that (see Theorem 2.1.5) in a graph  $G$ , the maximum fractional matching, i.e., with weight  $\mu_f(G)$ , can be achieved with weights only chosen from  $\{0, 1/2, 1\}$ .

**Theorem 6** ([10]). *For any graph  $G$ ,*

$$\mu_f(G) = \frac{1}{2} (|V(G)| - \max\{i(V(G) \setminus S) - |S|\}),$$

where  $i(X)$  denotes the number of isolated vertices in  $G[X]$ , and the maximum is taken over all  $S \subseteq V(G)$ .

*Proof of Theorem 3.* Given  $c, \alpha > 0$ , let  $d = \alpha c/9$ . We apply Lemma 5 with  $\delta = d/2$  and obtain  $\epsilon_1 > 0$ . Let  $\epsilon = \min\{\epsilon_1, d/6, 3d/(2c)\}$ . We then apply Lemma 4 with  $\epsilon$  and obtain  $M = M(\epsilon)$ . Let  $n \in \mathbb{N}$  be sufficiently large. Let  $G = (V, E)$  be a graph of order  $n$  such that  $\deg(v) = (1 \pm \epsilon)cn$  for every  $v \in V$ . We apply the Regularity Lemma (Lemma 4) on  $G$  with the constants  $\epsilon, d$  chosen as above and obtain a partition of  $V$  into  $V_0, V_1, \dots, V_t$  for some  $t \leq M$ , and a subgraph  $G'$  of  $G$  with the properties as described in Lemma 4. By moving at most one vertex from each  $V_i$ ,  $i \in [t]$  to  $V_0$ , we can assume that  $m := |V_i|$  is even. Thus we have  $|V_0| \leq \epsilon n + t \leq 2\epsilon n$ . Now for any  $v \in V$ ,

$$\deg_{G'}(v) - |V_0| > \deg_G(v) - (d + \epsilon)n - 2\epsilon n \geq (c - 2d)n.$$

Let  $\beta = 3d/c$ . Let  $H$  be the graph on  $[t]$  such that  $ij \in E(H)$  if and only if  $d(V_i, V_j) \geq d$ . We first assume that there exists a set  $S \subseteq [t]$ , such that  $i(V(H) \setminus S) - |S| \geq \beta t$ . In particular, let  $T$  be the collection of  $|S| + \beta t$  isolated vertices in  $H - S$ . Thus we have

$$e_G(T, S) \geq |T|m(\deg_{G'}(v) - |V_0|) \geq |T|m(c - 2d)n.$$

However, by averaging, this implies that there exists a vertex  $v \in V$  such that

$$\deg_G(v) \geq \deg_{G'}(v) \geq \frac{|T|m(c - 2d)n}{|S|m} \geq \frac{t}{t - \beta t}(c - 2d)n > (1 + \epsilon)cn,$$

by the definition of  $\beta$  and  $\epsilon$ , a contradiction. Thus we have  $i(V(H) \setminus S) - |S| \leq \beta t$  for any  $S \subseteq [t]$ . So by Theorem 6, we get  $\mu_f(H) \geq (1 - \beta)t/2$ . Moreover, there exists a fractional matching  $f$  such that  $\sum_{e \in E(H)} f(e) = \mu_f(H) \geq (1 - \beta)t/2$  and  $f(e) \in \{0, 1/2, 1\}$  for every edge  $e \in E(H)$ .

For each  $i \in [t]$  we arbitrarily split  $V_i$  into  $V_i^1$  and  $V_i^2$  each of size  $m/2$ . Thus the existence of  $f$  implies that we can partition  $V \setminus V_0$  into at least  $(1 - \beta)t/2$  pairs of sets each of form  $(V_i^a, V_j^b)$  with density at least  $d$ , where  $i, j \in [t]$ ,  $i \neq j$  and  $a, b \in [2]$ , and a set of at most  $2\beta t \cdot m$  vertices. Note that here (to simplify the argument) even if an edge  $ij \in E(H)$  receives weight 1, we still split it, e.g., as  $(V_i^1, V_j^1)$  and  $(V_i^2, V_j^2)$ . Thus every vertex of  $H$  is in at most two pairs so there are at most  $t$  pairs.

We will show that each such pair contains a cycle that covers all but at most  $2\epsilon m$  vertices. Indeed, fix any pair  $(V_i^a, V_j^b)$ , let  $A$  be the set of vertices in  $V_i^a$  whose degree to  $V_j^b$  is less than  $(d - \epsilon)|V_j^b|$ . Since  $d(A, V_j^b) < d - \epsilon$  and  $|V_j^b| > \epsilon m$ , the regularity of  $(V_i, V_j)$  implies that  $|A| \leq \epsilon m$ . Similarly let  $B$  be the set of vertices in  $V_j^b$  whose degree to  $V_i^a$  is less than  $(d - \epsilon)|V_i^a|$  and we have  $|B| \leq \epsilon m$ . Let  $A' \supseteq A$  and  $B' \supseteq B$  be arbitrary subsets of  $V_i^a$  and  $V_j^b$ , respectively, of size exactly  $\epsilon m$ . Now let  $X = V_i^a \setminus A'$  and  $Y = V_j^b \setminus B'$ , we get that  $(X, Y)$  is  $(\epsilon, d - 3\epsilon)$ -super-regular with density at least  $d - 3\epsilon$ , and  $|X| = |Y| = m - \epsilon m$ . Since  $d - 3\epsilon \geq d/2$ , by Lemma 5,  $(X, Y)$  contains a spanning cycle and we are done.

Let  $C = M$ . Thus we obtain a set of at most  $t \leq M = C$  vertex-disjoint cycles in  $G$  that covers all but at most

$$t \cdot 2\epsilon m + |V_0| + 2\beta t \cdot m \leq 3\beta n = \alpha n$$

vertices, completing the proof. □

### 3 Proof of Theorem 2

In the proof of Theorem 2 we use the trick of a ‘reservoir lemma’ from [8, 9]. Roughly speaking, we will reserve a random set  $R$  of vertices at the beginning of the proof, and use them to connect the paths returned by applying Theorem 3 on  $G - R$ . We first recall the following Chernoff’s bounds (see, e.g., [1]) for binomial random variables and for  $x > 0$ :

$$\begin{aligned} \mathbb{P}[\text{Bin}(n', \zeta) \geq n'\zeta + x] &< e^{-x^2/(2n'\zeta + x/3)} \\ \mathbb{P}[\text{Bin}(n', \zeta) \leq n'\zeta - x] &< e^{-x^2/(2n'\zeta)}. \end{aligned}$$

**Lemma 7.** *Given any  $c, \gamma, \epsilon > 0$ , the following holds for sufficiently large integer  $n$ . Let  $G$  be a  $\lceil cn \rceil$ -regular graph of order  $n$ . Then there exists a set  $R \subseteq V(G)$  such that  $|R| = (1 \pm \epsilon)\gamma n$  and every vertex of  $G$  has degree  $(1 \pm \epsilon)c\gamma n$  in  $R$ .*

*Proof.* We select the set  $R$  by including each vertex of  $G$  independently and randomly with probability  $\gamma$ . Note that  $|R|$  and  $\deg(v, R)$  for each  $v \in V(G)$  are both binomial random variables with expectation  $\gamma n$  and  $\gamma \lceil cn \rceil$ , respectively. By Chernoff's bounds, we get

$$\begin{aligned} \mathbb{P}[|R| > (1 + \epsilon)\gamma n] &< e^{-\epsilon^2 \gamma n / 3}, \quad \mathbb{P}[|R| < (1 - \epsilon)\gamma n] < e^{-\epsilon^2 \gamma n / 3}, \\ \mathbb{P}[x_v > (1 + \epsilon)c\gamma n] &< e^{-\epsilon^2 c\gamma n / 3}, \quad \mathbb{P}[x_v < (1 - \epsilon)c\gamma n] < e^{-\epsilon^2 c\gamma n / 3}, \end{aligned}$$

where  $x_v := \deg(v, R)$  for all  $v \in V(G)$ . Since  $(2n + 2)e^{-\epsilon^2 c\gamma n / 3} < 1$  because  $n$  is large enough, there is a choice of  $R$  with the desired properties.  $\square$

*Proof of Theorem 2.* Given  $c, \alpha \in (0, 1)$ , let  $\gamma = \alpha/4$ . We apply Theorem 3 with  $c$  and  $\alpha/2$  in place of  $\alpha$  and obtain  $\epsilon_1$  and  $C \in \mathbb{N}$ . Let  $\epsilon = \min\{\epsilon_1, ((\lfloor 1/c \rfloor + 1)c - 1)/3\}$ . Let  $G$  be a  $\lceil cn \rceil$ -regular graph of order  $n$ . We first pick the set  $R$  by Lemma 7. Let  $G_1 = G - R$  and  $n_1 = n - |R|$ . Thus for every vertex  $v \in V(G_1)$ , we know that  $\deg_{G_1}(v) = \lceil cn \rceil - (1 \pm \epsilon)c\gamma n$ . Since  $|R| = (1 \pm \epsilon)\gamma n$  we know that  $\deg_{G_1}(v) = (1 \pm \epsilon)cn_1$ . Indeed, for the upper bound we have

$$\deg_{G_1}(v) \leq c(n_1 + |R|) + 1 - (1 - \epsilon)c\gamma n \leq cn_1 + 1 + 2\epsilon c\gamma n \leq (1 + \epsilon)cn_1,$$

where in the last inequality we use  $n < 2n_1$  and  $\gamma \leq 1/4$ ; and the lower bound can be shown similarly. By applying Theorem 3 with  $\alpha/2$  in place of  $\alpha$ , we obtain a collection of at most  $C$  vertex-disjoint paths whose union covers all but at most  $\alpha|V(G_1)|/2 \leq \alpha n/2$  vertices of  $G_1$ .

Next we iteratively use the property of  $R$  to connect some pair of paths. We first explain the general case. Suppose there are at least  $\lfloor 1/c \rfloor + 1$  paths left. Indeed, let  $v_1, v_2, \dots, v_{\lfloor 1/c \rfloor + 1}$  be the (arbitrary one out of the two) ends of the  $\lfloor 1/c \rfloor + 1$  paths. Note that throughout the iteration there are at most  $C$  vertices in  $R$  that have been used for connecting and thus removed from  $R$ . So for each  $i$ , by  $\deg(v_i, R) - C \geq (1 - \epsilon)c\gamma n - C \geq (1 - 2\epsilon)c|R|$ , we have

$$(\lfloor 1/c \rfloor + 1)(\deg(v_i, R) - C) \geq (1 + 3\epsilon)(1 - 2\epsilon)|R| > |R|,$$

by the definition of  $\epsilon$ . Thus there exist two vertices  $v_i, v_j$  which have a common neighbor  $w$  in  $R$  so that we can connect the corresponding two paths by  $w$ . At the end, we obtain a collection of at most  $\lfloor 1/c \rfloor$  vertex-disjoint paths whose union covers all but at most  $\alpha n/2 + (1 + \epsilon)\gamma n \leq \alpha n$  vertices in  $G$ .

Second, assume that there are at least  $\lfloor 1/(2c) \rfloor + 1$  paths left and in addition that  $G$  is bipartite with bipartition  $X$  and  $Y$ . Note that since  $G$  is regular we have  $|X| = |Y| = n/2$  (so in particular  $n$  must be even). Fix  $\lfloor 1/(2c) \rfloor + 1$  paths. By throwing away at most one vertex from each path we can assume that each path has exactly one end vertex in  $X$  and

one in  $Y$ . Let  $v_1, v_2, \dots, v_{\lfloor 1/(2c) \rfloor + 1}$  be the end vertices in  $X$ . By the similar calculation, we can find a vertex  $w \in R \cap Y$  which connects some pair of paths. At the end, we obtain a collection of at most  $\lfloor 1/(2c) \rfloor$  vertex-disjoint paths whose union covers all but at most  $\alpha n/2 + (1 + \epsilon)\gamma n + C^2 \leq \alpha n$  vertices in  $G$ , because the iteration has at most  $C$  steps and in each step we threw away at most  $C$  vertices from the current paths.  $\square$

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