

Inverse Perron values and connectivity of a uniform hypergraph

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Abstract

In this paper, we show that a uniform hypergraph \mathcal{G} is connected if and only if one of its inverse Perron values is larger than 0. We give some bounds on the bipartition width, isoperimetric number and eccentricities of \mathcal{G} in terms of inverse Perron values. By using the inverse Perron values, we give an estimation of the edge connectivity of a 2-design, and determine the explicit edge connectivity of a symmetric design. Moreover, relations between the inverse Perron values and resistance distance of a connected graph are presented.

Mathematics Subject Classifications: 05C50, 05C65, 05C40, 05C12, 15A69

1 Introduction

Let $V(\mathcal{G})$ and $E(\mathcal{G})$ denote the vertex set and edge set of a hypergraph \mathcal{G} , respectively. \mathcal{G} is k -uniform if $|e| = k$ for each $e \in E(\mathcal{G})$. In particular, 2-uniform hypergraphs are usual graphs. For $i \in V(\mathcal{G})$, $E_i(\mathcal{G})$ denotes the set of edges containing i , and $d_i = |E_i(\mathcal{G})|$

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denotes the degree of i . The adjacency tensor [8] of a k -uniform hypergraph \mathcal{G} , denoted by $\mathcal{A}_{\mathcal{G}}$, is an order k dimension $|V(\mathcal{G})|$ tensor with entries

$$a_{i_1 i_2 \dots i_k} = \begin{cases} \frac{1}{(k-1)!}, & \text{if } \{i_1, i_2, \dots, i_k\} \in E(\mathcal{G}), \\ 0, & \text{otherwise.} \end{cases}$$

The *Laplacian tensor* [27] of \mathcal{G} is $\mathcal{L}_{\mathcal{G}} = \mathcal{D}_{\mathcal{G}} - \mathcal{A}_{\mathcal{G}}$, where $\mathcal{D}_{\mathcal{G}}$ is the diagonal tensor of vertex degrees of \mathcal{G} . Recently, the research on spectral hypergraph theory via tensors has attracted much attention [7-10,14,19,24]. The spectral properties of the Laplacian tensor of hypergraphs are studied in [13,25,27,29,35].

For an order k dimension n tensor $\mathcal{T} = (t_{i_1 i_2 \dots i_k})$, let $\mathcal{T}\mathbf{x}^k = \sum_{i_1, \dots, i_k=1}^n t_{i_1 i_2 \dots i_k} x_{i_1} \cdots x_{i_k}$.

The algebraic connectivity of a graph plays important roles in spectral graph theory [11]. Analogue to the algebraic connectivity of a graph, Qi [27] defined the *analytic connectivity* of a k -uniform hypergraph \mathcal{G} as

$$\alpha(\mathcal{G}) = \min_{j=1, \dots, n} \min \left\{ \mathcal{L}_{\mathcal{G}}\mathbf{x}^k : \mathbf{x} \in \mathbb{R}_+^n, \sum_{i=1}^n x_i^k = 1, x_j = 0 \right\},$$

where $n = |V(\mathcal{G})|$, \mathbb{R}_+^n denotes the set of nonnegative vectors of dimension n . Qi proved that \mathcal{G} is connected if and only if $\alpha(\mathcal{G}) > 0$. In [20], some bounds on $\alpha(\mathcal{G})$ were presented in terms of degree, vertex connectivity, diameter and isoperimetric number. A feasible trust region algorithm of $\alpha(\mathcal{G})$ was given in [9].

For any vertex j of a k -uniform hypergraph \mathcal{G} , we define the *inverse Perron value* of j as

$$\alpha_j(\mathcal{G}) = \min \left\{ \mathcal{L}_{\mathcal{G}}\mathbf{x}^k : \mathbf{x} \in \mathbb{R}_+^n, \sum_{i=1}^n x_i^k = 1, x_j = 0 \right\}.$$

Clearly, the analytic connectivity $\alpha(\mathcal{G}) = \min_{j \in V(\mathcal{G})} \alpha_j(\mathcal{G})$ is the minimum inverse Perron value. For a connected graph G , $\alpha_j(G)$ is the minimum eigenvalue of $\mathcal{L}_G(j)$, where $\mathcal{L}_G(j)$ is the principal submatrix of \mathcal{L}_G obtained by deleting the row and column corresponding to j . $\mathcal{L}_G(j)$ is nonsingular and its inverse $\mathcal{L}_G(j)^{-1}$ is a nonnegative matrix [16]. It is easy to see that $\alpha_j^{-1}(G)$ is the spectral radius of $\mathcal{L}_G(j)^{-1}$, which is called the Perron value of G . All inverse Perron values of a tree T can determine the algebraic connectivity of T [1, 15].

The resistance distance [17, 34] is a distance function on graphs. For two vertices i, j in a connected graph G , the *resistance distance* between i and j , denoted by $r_{ij}(G)$, is defined to be the effective resistance between them when unit resistors are placed on every edge of G . The *Kirchhoff index* [17, 33] of G , denoted by $Kf(G)$, is defined as the sum of resistance distances between all pairs of vertices in G , i.e., $Kf(G) = \sum_{\{i,j\} \subseteq V(G)} r_{ij}(G)$. $Kf(G)$ is a global robustness index. The resistance distance and Kirchhoff index in graphs have been investigated extensively in mathematical and chemical literatures [3,4,6,12,23,31,36].

This paper is organized as follows. In Section 2, some auxiliary lemmas are introduced. In Section 3, we show that a uniform hypergraph \mathcal{G} is connected if and only if one of its inverse Perron values is larger than 0, and some inequalities among the inverse Perron values, bipartition width, isoperimetric number and eccentricities of \mathcal{G} are established. Partial results improve some bounds in [20, 27]. We also use the inverse Perron values to estimate the edge connectivity of 2-designs. In Section 4, some inequalities among the inverse Perron values, resistance distance and Kirchhoff index of a connected graph are presented.

2 Preliminaries

For a positive integer n , let $[n] = \{1, 2, \dots, n\}$. An order m dimension n tensor $\mathcal{T} = (t_{i_1 \dots i_m})$ consists of n^m entries, where $i_j \in [n]$, $j \in [m]$. When $m = 2$, T is an $n \times n$ matrix. Let $\mathbb{R}^{[m,n]}$ denote the set of order m dimension n real tensors, and let \mathbb{R}_+^n denote the cone of nonnegative vectors in \mathbb{R}^n . For $\mathcal{T} = (t_{i_1 i_2 \dots i_m}) \in \mathbb{R}^{[m,n]}$ and $\mathbf{x} = (x_1, \dots, x_n)^T \in \mathbb{R}^n$, let $\mathcal{T}\mathbf{x}^{m-1} \in \mathbb{R}^n$ denote the vector whose i -th component is

$$(\mathcal{T}\mathbf{x}^{m-1})_i = \sum_{i_2, i_3, \dots, i_m=1}^n t_{i i_2 \dots i_m} x_{i_2} x_{i_3} \cdots x_{i_m},$$

and let $\mathbf{x}^{[m-1]} = (x_1^{m-1}, \dots, x_n^{m-1})^T$. In 2005, Qi [26] and Lim [21] proposed the concept of eigenvalues of tensors, independently. For $\mathcal{T} = (t_{i_1 i_2 \dots i_m}) \in \mathbb{R}^{[m,n]}$, if there exist a number $\lambda \in \mathbb{R}$ and a nonzero vector $\mathbf{x} = (x_1, \dots, x_n)^T \in \mathbb{R}^n$ such that $\mathcal{T}\mathbf{x}^{m-1} = \lambda \mathbf{x}^{[m-1]}$, then λ is called an *H-eigenvalue* of \mathcal{T} , \mathbf{x} is called an *H-eigenvector* of \mathcal{T} corresponding to λ .

For a vertex j of a k -uniform hypergraph \mathcal{G} , let $\mathcal{L}_{\mathcal{G}}(j) \in \mathbb{R}^{[k,n-1]}$ denote the principal subtensor of $\mathcal{L}_{\mathcal{G}} \in \mathbb{R}^{[k,n]}$ with index set $V(\mathcal{G}) \setminus \{j\}$. By Lemma 2.3 in [32], we know that $\alpha_j(\mathcal{G})$ is the smallest H-eigenvalue of $\mathcal{L}_{\mathcal{G}}(j)$ for any $j \in V(\mathcal{G})$.

A path \mathcal{P} of a uniform hypergraph \mathcal{G} is an alternating sequence of vertices and edges $v_0 e_1 v_1 e_2 \cdots v_{l-1} e_l v_l$, where v_0, \dots, v_l are distinct vertices of \mathcal{G} , e_1, \dots, e_l are distinct edges of \mathcal{G} and $v_{i-1}, v_i \in e_i$, for $i = 1, \dots, l$. The number of edges in \mathcal{P} is the length of \mathcal{P} . For all $u, v \in V(\mathcal{G})$, if there exists a path starting at u and terminating at v , then \mathcal{G} is said to be *connected* [5].

Lemma 1. [27] *The uniform hypergraph \mathcal{G} is connected if and only if $\alpha(\mathcal{G}) > 0$.*

Let \mathcal{G} be a k -uniform hypergraph, $S \neq \emptyset$ be a proper subset of $V(\mathcal{G})$. Denote $\bar{S} = V(\mathcal{G}) \setminus S$. The edge-cut set $E(S, \bar{S})$ consists of edges whose vertices are in both S and \bar{S} . The minimum cardinality of such an edge-cut set is called *edge connectivity* of \mathcal{G} , denote by $e(\mathcal{G})$.

Lemma 2. [27] *Let \mathcal{G} be a k -uniform hypergraph with n vertices. Then*

$$e(\mathcal{G}) \geq \frac{n}{k} \alpha(\mathcal{G}).$$

The $\{1\}$ -inverse of a matrix M is a matrix X such that $MXM = M$. Let $M^{(1)}$ denote any $\{1\}$ -inverse of M , and let $(M)_{ij}$ denote the (i, j) -entry of M .

Lemma 3. [2, 34] Let G be a connected graph. Then

$$r_{ij}(G) = (\mathcal{L}_G^{(1)})_{ii} + (\mathcal{L}_G^{(1)})_{jj} - (\mathcal{L}_G^{(1)})_{ij} - (\mathcal{L}_G^{(1)})_{ji}.$$

Let $\text{tr}(A)$ denote the trace of the square matrix A , and let \mathbf{e} denote an all-ones column vector.

Lemma 4. [30] Let G be a connected graph of order n . Then

$$Kf(G) = n\text{tr}(\mathcal{L}_G^{(1)}) - \mathbf{e}^\top \mathcal{L}_G^{(1)} \mathbf{e}.$$

Lemma 5. [2] Let G be a connected graph with n vertices and $i \in [n]$. Let $\mathcal{L}_G =$

$$\begin{pmatrix} L_1 & \mathbf{x} & L_2 \\ \mathbf{x}^\top & d_i & \mathbf{y} \\ L_2^\top & \mathbf{y}^\top & L_3 \end{pmatrix}, \text{ where } L_1 \in \mathbb{R}^{(i-1) \times (i-1)}, L_3 \in \mathbb{R}^{(n-i) \times (n-i)}, \mathbf{x} \in \mathbb{R}^{i-1}, \mathbf{y}^\top \in \mathbb{R}^{n-i}.$$

Suppose $\mathcal{L}_G(i)^{-1} = \begin{pmatrix} \widetilde{L}_1 & \widetilde{L}_2 \\ \widetilde{L}_2^\top & \widetilde{L}_3 \end{pmatrix}$, where $\widetilde{L}_1 \in \mathbb{R}^{(i-1) \times (i-1)}$, $\widetilde{L}_3 \in \mathbb{R}^{(n-i) \times (n-i)}$. Then

$$\begin{pmatrix} \widetilde{L}_1 & \mathbf{0} & \widetilde{L}_2 \\ \mathbf{0} & 0 & \mathbf{0} \\ \widetilde{L}_2^\top & \mathbf{0} & \widetilde{L}_3 \end{pmatrix} \text{ is a symmetric } \{1\}\text{-inverse of } \mathcal{L}_G.$$

3 Inverse Perron values of uniform hypergraphs

In the following theorem, the relationship between inverse Perron values and connectivity of a hypergraph is presented.

Theorem 6. Let \mathcal{G} be a k -uniform hypergraph. Then the following statements are equivalent:

- (1) \mathcal{G} is connected.
- (2) $\alpha_j(\mathcal{G}) > 0$ for all $j \in V(\mathcal{G})$.
- (3) $\alpha_j(\mathcal{G}) > 0$ for some $j \in V(\mathcal{G})$.

Proof. (1) \implies (2). If \mathcal{G} is connected, then by Lemma 1, we know that $\alpha_j(\mathcal{G}) > 0$ for all $j \in V(\mathcal{G})$.

(2) \implies (3). Obviously.

(3) \implies (1). Suppose that \mathcal{G} is disconnected. For any $j \in V(\mathcal{G})$, let \mathcal{G}_1 be the component of \mathcal{G} such that $j \notin V(\mathcal{G}_1)$. Let $\mathbf{x} = (x_1, \dots, x_{|V(\mathcal{G})|})^\top$ be the vector satisfying

$$x_i = \begin{cases} |V(\mathcal{G}_1)|^{-\frac{1}{k}}, & \text{if } i \in V(\mathcal{G}_1), \\ 0, & \text{otherwise.} \end{cases}$$

Clearly, we have $\sum_{i=1}^n x_i^k = 1$. Then we have $0 \leq \alpha_j(\mathcal{G}) \leq \mathcal{L}_G \mathbf{x}^k = 0$ for any $j \in V(\mathcal{G})$, a contradiction to (3). Hence \mathcal{G} is connected if (3) holds. \square

The *bipartition width* of a hypergraph \mathcal{G} is defined as [18, 28]

$$\text{bw}(\mathcal{G}) = \min \left\{ |E(S, \overline{S})| : S \subseteq V(\mathcal{G}), |S| = \left\lfloor \frac{n}{2} \right\rfloor \right\},$$

where $\lfloor \frac{n}{2} \rfloor$ denotes the maximum integer not larger than $\frac{n}{2}$. The computation of $\text{bw}(\mathcal{G})$ is very difficult even for the graph case. In [22], Mohar and Poljak used the algebraic connectivity to obtain a lower bound on the bipartition width of a graph. In the following theorem, we use the inverse Perron values to obtain a lower bound on the bipartition width of a uniform hypergraph.

Theorem 7. *Let \mathcal{G} be a k -uniform hypergraph with n vertices. Then*

$$\text{bw}(\mathcal{G}) \geq \frac{n + (-1)^n}{k(n+1)} \sum_{j=1}^n \alpha_j(\mathcal{G}).$$

Proof. Suppose that $S_0 \subseteq V(\mathcal{G})$ satisfying $|S_0| = \lfloor \frac{n}{2} \rfloor$ and $|E(S_0, \overline{S_0})| = \text{bw}(\mathcal{G})$. Let $\mathbf{x} = (x_1, \dots, x_n)^T$ be the vector satisfying

$$x_i = \begin{cases} |S_0|^{-\frac{1}{k}}, & i \in S_0, \\ 0, & i \in \overline{S_0}. \end{cases}$$

Then $\sum_{i=1}^n x_i^k = 1$. For $j \in \overline{S_0}$, we get

$$\begin{aligned} \alpha_j(\mathcal{G}) &\leq \mathcal{L}_{\mathcal{G}} \mathbf{x}^k = \sum_{\{i_1, \dots, i_k\} \in E(\mathcal{G})} (x_{i_1}^k + \dots + x_{i_k}^k - kx_{i_1} \dots x_{i_k}) \\ &\leq \sum_{\{i_1, \dots, i_k\} \in E(S_0, \overline{S_0})} (x_{i_1}^k + \dots + x_{i_k}^k - kx_{i_1} \dots x_{i_k}) \\ &= \frac{1}{|S_0|} \sum_{e \in E(S_0, \overline{S_0})} |e \cap S_0| = \frac{t(S_0) \text{bw}(\mathcal{G})}{|S_0|}, \end{aligned} \tag{1}$$

where $t(S_0) = \frac{1}{|E(S_0, \overline{S_0})|} \sum_{e \in E(S_0, \overline{S_0})} |e \cap S_0|$.

Similarly, for $j \in S_0$, we obtain

$$\alpha_j(\mathcal{G}) \leq \frac{(k - t(S_0)) \text{bw}(\mathcal{G})}{|\overline{S_0}|}. \tag{2}$$

Combining (1) and (2), we get

$$\sum_{j=1}^n \alpha_j(\mathcal{G}) = \sum_{j \in S_0} \alpha_j(\mathcal{G}) + \sum_{j \in \overline{S_0}} \alpha_j(\mathcal{G}) \leq \frac{|S_0|(k - t(S_0)) \text{bw}(\mathcal{G})}{|S_0|} + \frac{|\overline{S_0}| t(S_0) \text{bw}(\mathcal{G})}{|S_0|}.$$

If n is even, then $|S_0| = |\overline{S_0}|$ and $\text{bw}(\mathcal{G}) \geq \frac{1}{k} \sum_{j=1}^n \alpha_j(\mathcal{G})$. If n is odd, then $|S_0| = |\overline{S_0}| - 1 = \frac{n-1}{2}$ and

$$\sum_{j=1}^n \alpha_j(\mathcal{G}) \leq k \frac{|\overline{S_0}|}{|S_0|} \text{bw}(\mathcal{G}) = \frac{k(n+1)\text{bw}(\mathcal{G})}{n-1}, \quad \text{bw}(\mathcal{G}) \geq \frac{n-1}{k(n+1)} \sum_{j=1}^n \alpha_j(\mathcal{G}). \quad \square$$

The *isoperimetric number* of a k -uniform hypergraph \mathcal{G} is defined as

$$i(\mathcal{G}) = \min \left\{ \frac{|E(S, \overline{S})|}{|S|} : S \subseteq V(\mathcal{G}), 0 < |S| \leq \frac{|V(\mathcal{G})|}{2} \right\}.$$

Let $\beta(\mathcal{G}) = \max_{j \in V(\mathcal{G})} \alpha_j(\mathcal{G})$ denote the maximum inverse Perron value of \mathcal{G} . In [20], it was shown that $i(\mathcal{G}) \geq \frac{2}{k} \alpha(\mathcal{G})$. We improve it as follows.

Theorem 8. *Let \mathcal{G} be a k -uniform hypergraph. Then*

$$i(\mathcal{G}) \geq \frac{\alpha(\mathcal{G}) + \beta(\mathcal{G})}{k}.$$

Proof. Suppose that $S_1 \subseteq V(\mathcal{G})$ satisfying $0 < |S_1| \leq \frac{|V(\mathcal{G})|}{2}$ and $\frac{|E(S_1, \overline{S_1})|}{|S_1|} = i(\mathcal{G})$. Let $\mathbf{x} = (x_1, \dots, x_n)^T$ be the vector satisfying

$$x_i = \begin{cases} |S_1|^{-\frac{1}{k}}, & i \in S_1, \\ 0, & i \in \overline{S_1}. \end{cases}$$

Then $\sum_{i=1}^n x_i^k = 1$. For $j \in \overline{S_1}$, we obtain

$$\alpha_j(\mathcal{G}) \leq \mathcal{L}_{\mathcal{G}} \mathbf{x}^k = \frac{t(S_1) |E(S_1, \overline{S_1})|}{|S_1|} = t(S_1) i(\mathcal{G}), \quad (3)$$

where $t(S_1) = \frac{1}{|E(S_1, \overline{S_1})|} \sum_{e \in E(S_1, \overline{S_1})} |e \cap S_1|$.

Similarly, for $j \in S_1$, we get

$$\alpha_j(\mathcal{G}) \leq \frac{(k - t(S_1)) |E(S_1, \overline{S_1})|}{|\overline{S_1}|} \leq (k - t(S_1)) i(\mathcal{G}). \quad (4)$$

Let $\alpha_s(\mathcal{G}) = \beta(\mathcal{G})$. If $s \in \overline{S_1}$, by (3), we get

$$\beta(\mathcal{G}) = \alpha_s(\mathcal{G}) \leq t(S_1) i(\mathcal{G}).$$

From (4), we have

$$\alpha(\mathcal{G}) = \min_{j \in V(\mathcal{G})} \alpha_j(\mathcal{G}) \leq \min_{j \in S_1} \alpha_j(\mathcal{G}) \leq (k - t(S_1)) i(\mathcal{G}).$$

Then

$$\alpha(\mathcal{G}) + \beta(\mathcal{G}) \leq t(S_1)i(\mathcal{G}) + (k - t(S_1))i(\mathcal{G}) = ki(\mathcal{G}).$$

Similarly, if $s \in S_1$, we can also obtain $\alpha(\mathcal{G}) + \beta(\mathcal{G}) \leq ki(\mathcal{G})$.

From the above discussion, we get $i(\mathcal{G}) \geq \frac{\alpha(\mathcal{G}) + \beta(\mathcal{G})}{k}$. □

The distance $d(u, v)$ between two distinct vertices u and v of \mathcal{G} is the length of the shortest path connecting them. The eccentricity of a vertex v is $\text{ecc}(v) = \max\{d(u, v) : u \in V(\mathcal{G})\}$. The diameter and radius of \mathcal{G} are defined as $\text{diam}(\mathcal{G}) = \max_{v \in V(\mathcal{G})} \text{ecc}(v)$ and $\text{rad}(\mathcal{G}) = \min_{v \in V(\mathcal{G})} \text{ecc}(v)$, respectively.

Theorem 9. *Let \mathcal{G} be a connected k -uniform hypergraph with n vertices. Then*

$$\text{ecc}(j) \geq \frac{k}{2(k-1)(n-1)\alpha_j(\mathcal{G})}, \quad j \in V(\mathcal{G}).$$

Proof. For $j \in V(\mathcal{G})$, let $\mathbf{x} = (x_1, \dots, x_n)^T \in \mathbb{R}_+^n$ satisfying $x_j = 0$, $\sum_{i=1}^n x_i^k = 1$ and $\alpha_j(\mathcal{G}) = \mathcal{L}_{\mathcal{G}}\mathbf{x}^k$. Then

$$\alpha_j(\mathcal{G}) = \mathcal{L}_{\mathcal{G}}\mathbf{x}^k = \sum_{\{i_1, \dots, i_k\} \in E(\mathcal{G})} (x_{i_1}^k + \dots + x_{i_k}^k - kx_{i_1} \dots x_{i_k}). \quad (5)$$

From AM-GM inequality, it yields that

$$\sum_{1 \leq s < t \leq k} x_{i_s}^{\frac{k}{2}} x_{i_t}^{\frac{k}{2}} \geq \frac{k(k-1)}{2} \left(\prod_{1 \leq s < t \leq k} x_{i_s}^{\frac{k}{2}} x_{i_t}^{\frac{k}{2}} \right)^{\frac{2}{k(k-1)}} = \frac{k(k-1)}{2} x_{i_1} \dots x_{i_k}. \quad (6)$$

By (5) and (6), we have

$$\begin{aligned} \alpha_j(\mathcal{G}) &\geq \sum_{\{i_1, \dots, i_k\} \in E(\mathcal{G})} \left(x_{i_1}^k + \dots + x_{i_k}^k - \frac{2}{k-1} \sum_{1 \leq s < t \leq k} x_{i_s}^{\frac{k}{2}} x_{i_t}^{\frac{k}{2}} \right) \\ &= \frac{1}{k-1} \sum_{\{i_1, \dots, i_k\} \in E(\mathcal{G})} \sum_{1 \leq s < t \leq k} \left(x_{i_s}^{\frac{k}{2}} - x_{i_t}^{\frac{k}{2}} \right)^2 \\ &= \frac{1}{k-1} \sum_{e \in E(\mathcal{G})} \sum_{s, t \in e} \left(x_s^{\frac{k}{2}} - x_t^{\frac{k}{2}} \right)^2. \end{aligned} \quad (7)$$

Let $v_0 \in \{i | x_i = \max_{p \in V(\mathcal{G})} x_p\}$. Let $\mathcal{P} = v_0 e_1 v_1 e_2 \dots v_{l-1} e_l v_l$ be the shortest path from vertex v_0 to vertex $v_l = j$. Then $x_{v_0}^k \geq \frac{1}{n-1}$, $x_{v_l} = 0$ and

$$\sum_{e \in E(\mathcal{G})} \sum_{s, t \in e} \left(x_s^{\frac{k}{2}} - x_t^{\frac{k}{2}} \right)^2 \geq \sum_{e \in E(\mathcal{P})} \sum_{s, t \in e} \left(x_s^{\frac{k}{2}} - x_t^{\frac{k}{2}} \right)^2$$

$$\begin{aligned}
&\geq \sum_{i=1}^l \left(\left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 + \sum_{u_j \in e_i \setminus \{v_{i-1}, v_i\}} \left(\left(x_{v_{i-1}}^{\frac{k}{2}} - x_{u_j}^{\frac{k}{2}} \right)^2 + \left(x_{u_j}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 \right) \right) \\
&\geq \sum_{i=1}^l \left(\left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 + \frac{1}{2} \sum_{u_j \in e_i \setminus \{v_{i-1}, v_i\}} \left(x_{v_{i-1}}^{\frac{k}{2}} - x_{u_j}^{\frac{k}{2}} + x_{u_j}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 \right) \\
&= \sum_{i=1}^l \left(\left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 + \frac{k-2}{2} \left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 \right) \\
&= \frac{k}{2} \sum_{i=1}^l \left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2.
\end{aligned}$$

By Cauchy-Schwarz inequality, we obtain

$$\begin{aligned}
\sum_{e \in E(\mathcal{G})} \sum_{s, t \in e} \left(x_s^{\frac{k}{2}} - x_t^{\frac{k}{2}} \right)^2 &\geq \frac{k}{2} \sum_{i=1}^l \left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right)^2 \geq \frac{k}{2l} \left(\sum_{i=1}^l \left(x_{v_{i-1}}^{\frac{k}{2}} - x_{v_i}^{\frac{k}{2}} \right) \right)^2 \\
&= \frac{k}{2l} \left(x_{v_0}^{\frac{k}{2}} - x_{v_l}^{\frac{k}{2}} \right)^2 \geq \frac{k}{2\text{ecc}(j)} \left(x_{v_0}^{\frac{k}{2}} - x_{v_l}^{\frac{k}{2}} \right)^2 \geq \frac{k}{2(n-1)\text{ecc}(j)}.
\end{aligned} \tag{8}$$

From (7) and (8), it yields that

$$\alpha_j(\mathcal{G}) \geq \frac{k}{2(k-1)(n-1)\text{ecc}(j)}, \quad \text{ecc}(j) \geq \frac{k}{2(k-1)(n-1)\alpha_j(\mathcal{G})}. \quad \square$$

For a connected k -uniform hypergraph \mathcal{G} with n vertices, [20] showed that

$$\text{diam}(\mathcal{G}) \geq \frac{4}{n^2(k-1)\alpha(\mathcal{G})}.$$

By Theorem 9, we obtain the following improved result.

Corollary 10. *Let \mathcal{G} be a connected k -uniform hypergraph with n vertices. Then*

$$\text{diam}(\mathcal{G}) \geq \frac{k}{2(k-1)(n-1)\alpha(\mathcal{G})}, \quad \text{rad}(\mathcal{G}) \geq \frac{k}{2(k-1)(n-1)\beta(\mathcal{G})}.$$

In [27], it was shown that $\alpha(\mathcal{G}) \leq \delta$, where δ is the minimum degree of \mathcal{G} . We improve it as follows.

Theorem 11. *Let \mathcal{G} be a k -uniform hypergraph with n vertices. Then*

$$\alpha_j(\mathcal{G}) \leq \frac{(k-1)d_j}{n-1}, \quad j \in V(\mathcal{G}).$$

Proof. For $j \in V(\mathcal{G})$, let $\mathbf{x} = (x_1, \dots, x_n)^T$ be the vector satisfying

$$x_i = \begin{cases} (n-1)^{-\frac{1}{k}}, & i \neq j, \\ 0, & i = j. \end{cases}$$

Then $\sum_{i=1}^n x_i^k = 1$, and we get

$$\begin{aligned} \alpha_j(\mathcal{G}) &\leq \mathcal{L}_{\mathcal{G}} \mathbf{x}^k = \sum_{\{i_1, \dots, i_k\} \in E(\mathcal{G})} (x_{i_1}^k + \dots + x_{i_k}^k - kx_{i_1} \dots x_{i_k}) \\ &= \sum_{\{i_1, \dots, i_k\} \in E_j(\mathcal{G})} (x_{i_1}^k + \dots + x_{i_k}^k) = \frac{(k-1)d_j}{n-1}, \end{aligned}$$

where $E_j(\mathcal{G})$ denotes the set of edges containing j . □

By Theorem 11, we obtain the following result.

Corollary 12. *Let \mathcal{G} be a k -uniform hypergraph with n vertices and m edges. Then*

$$\sum_{j=1}^n \alpha_j(\mathcal{G}) \leq \frac{(k-1)km}{n-1}, \quad j \in V(\mathcal{G}).$$

Let \mathcal{G} be a k -uniform hypergraph. For $x, y \in V(\mathcal{G})$, let $c(x, y) = |\{e \in E(\mathcal{G}) : x, y \in e\}|$. A 2 -(n, b, k, r, λ) design can be regarded as a k -uniform r -regular hypergraph \mathcal{G} on n vertices, b edges, and $c(x, y) = \lambda$ for any pair of distinct $x, y \in V(\mathcal{G})$. A 2-design satisfying $n = b$ is called a symmetric design.

Theorem 13. *Let \mathcal{G} be a connected k -uniform hypergraph with n vertices. Then \mathcal{G} is a 2-design if and only if $\alpha_1(\mathcal{G}) = \dots = \alpha_n(\mathcal{G}) = \frac{\Delta(k-1)}{n-1}$, where Δ is the maximum degree of \mathcal{G} .*

Proof. We first prove the necessity. If \mathcal{G} is a 2 -(n, b, k, r, λ) design, then $\lambda(n-1) = r(k-1)$ and $\Delta = r = d_1 = \dots = d_n$. For any $j \in V(\mathcal{G})$, by Theorem 11, we have

$$\alpha_j(\mathcal{G}) \leq \frac{r(k-1)}{n-1} = \lambda. \quad (9)$$

Let $\mathbf{x} = (x_1, \dots, x_n)^T \in \mathbb{R}_+^n$ satisfying $x_j = 0$, $\sum_{i=1}^n x_i^k = 1$ and $\alpha_j(\mathcal{G}) = \mathcal{L}_{\mathcal{G}} \mathbf{x}^k$. Then we get

$$\alpha_j(\mathcal{G}) = \mathcal{L}_{\mathcal{G}} \mathbf{x}^k \geq \sum_{\{i_1, \dots, i_k\} \in E_j(\mathcal{G})} (x_{i_1}^k + \dots + x_{i_k}^k - kx_{i_1} \dots x_{i_k}) = \lambda \sum_{i \neq j} x_i^k = \lambda. \quad (10)$$

Combining (9) and (10), we get

$$\alpha_1(\mathcal{G}) = \dots = \alpha_n(\mathcal{G}) = \lambda = \frac{r(k-1)}{n-1} = \frac{\Delta(k-1)}{n-1}.$$

Next we prove the sufficiency. Let $\alpha_1(\mathcal{G}) = \cdots = \alpha_n(\mathcal{G}) = \frac{\Delta(k-1)}{n-1}$. From Theorem 11, we obtain $d_1 = \cdots = d_n = \Delta$. Let $\mathbf{z} = \left((n-1)^{-\frac{1}{k}}, \dots, (n-1)^{-\frac{1}{k}} \right)^T \in \mathbb{R}_+^{n-1}$. For $j \in V(\mathcal{G})$, let $\mathbf{y} = (y_1, \dots, y_n)^T \in \mathbb{R}_+^n$ be a vector such that $y_i = 0$ if $i = j$ and $y_i = (n-1)^{-\frac{1}{k}}$ otherwise. Then

$$\begin{aligned} \mathcal{L}_{\mathcal{G}}\mathbf{y}^k &= \sum_{\{i_1, \dots, i_k\} \in E(\mathcal{G})} (y_{i_1}^k + \cdots + y_{i_k}^k - ky_{i_1} \cdots y_{i_k}) = \sum_{\{i_1, \dots, i_k\} \in E_j(\mathcal{G})} (y_{i_1}^k + \cdots + y_{i_k}^k) \\ &= \frac{\Delta(k-1)}{n-1} = \alpha_j(\mathcal{G}) = \alpha(\mathcal{G}). \end{aligned}$$

We know that $\alpha(\mathcal{G}) = \alpha_j(\mathcal{G})$ is the smallest H-eigenvalue of $\mathcal{L}_{\mathcal{G}}(j)$. Since $\mathcal{L}_{\mathcal{G}}(j)\mathbf{z}^k = \mathcal{L}_{\mathcal{G}}\mathbf{y}^k = \alpha(\mathcal{G})\mathbf{z}^k$, \mathbf{z} is an H-eigenvector corresponding to $\alpha(\mathcal{G})$, that is

$$\alpha(\mathcal{G})\mathbf{z}^{[k-1]} = \mathcal{L}_{\mathcal{G}}(j)\mathbf{z}^{k-1}.$$

For all $i \in V(\mathcal{G}) \setminus \{j\}$, we have

$$\begin{aligned} \alpha(\mathcal{G}) &= \frac{1}{z_i^{k-1}} (\mathcal{L}_{\mathcal{G}}(j)\mathbf{z}^{k-1})_i = \frac{1}{z_i^{k-1}} \sum_{i_2, \dots, i_k \neq j} (\mathcal{L}_{\mathcal{G}}(j))_{ii_2 \dots i_k} z_{i_2} \cdots z_{i_k} \\ &= \sum_{i_2, \dots, i_k \neq j} (\mathcal{L}_{\mathcal{G}})_{ii_2 \dots i_k} = c(i, j). \end{aligned}$$

So $c(i, j) = \alpha(\mathcal{G})$ for any pair of distinct $i, j \in V(\mathcal{G})$, which implies that \mathcal{G} is a 2-design. \square

We give an estimation of the edge connectivity of a 2-design as follows.

Theorem 14. *Let \mathcal{G} be a 2- (n, b, k, r, λ) design. Then*

$$\frac{n\lambda}{k} \leq e(\mathcal{G}) \leq \frac{(n-1)\lambda}{k-1}.$$

Moreover, if \mathcal{G} is a symmetric design, then $e(\mathcal{G}) = k = r$.

Proof. Since \mathcal{G} is a 2- (n, b, k, r, λ) design, we have $\lambda(n-1) = r(k-1)$. By Theorem 13, we have

$$\alpha(\mathcal{G}) = \frac{r(k-1)}{n-1} = \lambda.$$

It follows from Lemma 2 that

$$\frac{n\lambda}{k} = \frac{n}{k}\alpha(\mathcal{G}) \leq e(\mathcal{G}) \leq r = \frac{(n-1)\lambda}{k-1}. \quad (11)$$

Moreover, if \mathcal{G} is a symmetric design, then $n = b$. Since $nr = bk$, we have $r = k$. From $\lambda(n-1) = r(k-1)$ and (11), we have

$$\frac{n(k-1)}{n-1} \leq e(\mathcal{G}) \leq k.$$

Since $e(\mathcal{G})$ is a positive integer, we get $e(\mathcal{G}) = k = r$. \square

4 Inverse Perron values and resistance distance of graphs

For a vertex i of a connected graph G , we define its resistance eccentricity as $r_i(G) = \max_{j \in V(G)} r_{ij}$.

Theorem 15. *Let G be a connected graph. For any $i \in V(G)$, we have*

$$r_i(G) \leq \frac{1}{\alpha_i(G)}.$$

Proof. Without loss of generality, assume that i is the vertex corresponding to the last row of the Laplacian matrix \mathcal{L}_G . Since $\alpha_i(G)$ is the minimum eigenvalue of the principal submatrix $\mathcal{L}_G(i)$, $\alpha_i^{-1}(G)$ is the spectral radius of the symmetric nonnegative matrix $\mathcal{L}_G(i)^{-1}$. So $\alpha_i^{-1}(G) \geq \max_{j \neq i} (\mathcal{L}_G(i)^{-1})_{jj}$.

By Lemmas 5 and 3, we get $r_{ij}(G) = (\mathcal{L}_G(i)^{-1})_{jj}$ for any $j \neq i$. Hence

$$\alpha_i^{-1}(G) \geq \max_{j \neq i} (\mathcal{L}_G(i)^{-1})_{jj} = r_i(G),$$

$$r_i(G) \leq \frac{1}{\alpha_i(G)}. \quad \square$$

For a vertex i of a connected graph G , its resistance centrality is defined as $Kf_i(G) = \sum_{j \in V(G)} r_{ij}(G)$. It is used to measure the centrality of a network [4]. Note that $Kf(G) = \sum_{\{i,j\} \subseteq V(G)} r_{ij}(G) = \frac{1}{2} \sum_{i \in V(G)} Kf_i(G)$.

Theorem 16. *Let G be a connected graph with n vertices. For any $i \in V(G)$, we have*

$$nKf_i(G) - Kf(G) \leq \frac{n-1}{\alpha_i(G)}.$$

Proof. Note that $\alpha_i^{-1}(G)$ is the maximum eigenvalue of the symmetric matrix $\mathcal{L}_G(i)^{-1}$. Let \mathbf{e} be the all-ones column vector, then

$$\alpha_i^{-1}(G) \geq \frac{\mathbf{e}^\top \mathcal{L}_G(i)^{-1} \mathbf{e}}{\mathbf{e}^\top \mathbf{e}} = \frac{\mathbf{e}^\top \mathcal{L}_G(i)^{-1} \mathbf{e}}{n-1}.$$

By Lemmas 5 and 4, we have

$$Kf(G) = n \operatorname{tr}(\mathcal{L}_G(i)^{-1}) - \mathbf{e}^\top \mathcal{L}_G(i)^{-1} \mathbf{e}.$$

From Lemmas 5 and 3, we get $r_{ij}(G) = (\mathcal{L}_G(i)^{-1})_{jj}$ for any $j \neq i$. Hence $\operatorname{tr}(\mathcal{L}_G(i)^{-1}) = Kf_i(G)$ and

$$Kf(G) = nKf_i(G) - \mathbf{e}^\top \mathcal{L}_G(i)^{-1} \mathbf{e}.$$

By $\alpha_i^{-1}(G) \geq \frac{\mathbf{e}^\top \mathcal{L}_G(i)^{-1} \mathbf{e}}{n-1}$ we get

$$\alpha_i^{-1}(G) \geq \frac{\mathbf{e}^\top \mathcal{L}_G(i)^{-1} \mathbf{e}}{n-1} = \frac{nKf_i(G) - Kf(G)}{n-1},$$

$$nKf_i(G) - Kf(G) \leq \frac{n-1}{\alpha_i(G)}. \quad \square$$

Corollary 17. *Let G be a connected graph with n vertices. Then*

$$Kf(G) \leq \frac{n-1}{n} \sum_{i=1}^n \alpha_i^{-1}(G).$$

Proof. By Theorem 16, we have

$$\sum_{i=1}^n \frac{n-1}{\alpha_i(G)} \geq \sum_{i=1}^n (nKf_i(G) - Kf(G)) = nKf(G),$$

$$Kf(G) \leq \frac{n-1}{n} \sum_{i=1}^n \alpha_i^{-1}(G). \quad \square$$

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