3-uniform hypergraphs without a cycle of length five

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Abstract

In this paper we show that the maximum number of hyperedges in a 3-uniform hypergraph on n vertices without a (Berge) cycle of length five is less than $(0.254 + o(1))n^{3/2}$, improving an estimate of Bollobás and Győri.

We obtain this result by showing that not many 3-paths can start from certain subgraphs of the shadow.

Mathematics Subject Classifications: 05C65, 05D99

1 Introduction

A hypergraph H = (V, E) is a family E of distinct subsets of a finite set V. The members of E are called *hyperedges* and the elements of V are called *vertices*. A hypergraph is called 3-uniform if each member of E has size 3. A hypergraph H = (V, E) is called *linear* if every two hyperedges have at most one vertex in common.

A Berge cycle of length $k \ge 2$, denoted Berge- C_k , is an alternating sequence of distinct vertices and distinct edges of the form $v_1, h_1, v_2, h_2, \ldots, v_k, h_k$ where $v_i, v_{i+1} \in h_i$ for each $i \in \{1, 2, \ldots, k-1\}$ and $v_k, v_1 \in h_k$. (Note that if a hypergraph does not contain a Berge- C_2 , then it is linear.) This definition of a hypergraph cycle is the classical definition due to Berge. More generally, if F = (V(F), E(F)) is a graph and $\mathcal{Q} = (V(\mathcal{Q}), E(\mathcal{Q}))$ is a hypergraph, then we say \mathcal{Q} is *Berge-F* if there is a bijection $\phi : E(F) \to E(\mathcal{Q})$ such that $e \subseteq \phi(e)$ for all $e \in E(F)$. In other words, given a graph F we can obtain a Berge-F by replacing each edge of F with a hyperedge that contains it.

Given a family of graphs \mathcal{F} , we say that a hypergraph \mathcal{H} is *Berge-\mathcal{F}-free* if for every $F \in \mathcal{F}$, the hypergraph \mathcal{H} does not contain a Berge-F as a subhypergraph. The maximum possible number of hyperedges in a Berge- \mathcal{F} -free 3-uniform hypergraph on n vertices is the *Turán number* of Berge- \mathcal{F} , and is denoted by $ex_3(n, \mathcal{F})$. When $\mathcal{F} = \{F\}$ then we simply write $ex_3(n, F)$ instead of $ex_3(n, \{F\})$.

Determining $ex_3(n, \{C_2, C_3\})$ is basically equivalent to the famous (6, 3)-problem. This was settled by Ruzsa and Szemerédi in their classical paper [23], showing that $n^{2-\frac{c}{\sqrt{\log n}}} < ex_3(n, \{C_2, C_3\}) = o(n^2)$ for some constant c > 0. An important Turán-type extremal result for Berge cycles is due to Lazebnik and Verstraëte [21], who studied the maximum number of hyperedges in an *r*-uniform hypergraph containing no Berge cycle of length less than five (i.e., girth five). They showed the following.

Theorem 1 (Lazebnik, Verstraëte [21]).

$$ex_3(n, \{C_2, C_3, C_4\}) = \frac{1}{6}n^{3/2} + o(n^{3/2}).$$

The systematic study of the Turán number of Berge cycles started with the study of Berge triangles by Győri [15], and continued with the study of Berge five cycles by Bollobás and Győri [1] who showed the following.

Theorem 2 (Bollobás, Győri [1]).

$$(1+o(1))\frac{n^{3/2}}{3\sqrt{3}} \leq \exp_3(n, C_5) \leq \sqrt{2}n^{3/2} + 4.5n.$$

The following example of Bollobás and Győri proves the lower bound in Theorem 2.

Bollobás-Győri Example. Take a C_4 -free bipartite graph G_0 with n/3 vertices in each part and $(1 + o(1))(n/3)^{3/2}$ edges. In one part, replace each vertex u of G_0 by a pair of two new vertices u_1 and u_2 , and add the triple u_1u_2v for each edge uv of G_0 . It is easy to check that the resulting hypergraph H does not contain a Berge cycle of length 5. Moreover, the number of hyperedges in H is the same as the number of edges in G_0 .

In this paper, we improve Theorem 2 as follows.

Theorem 3.

$$ex_3(n, C_5) < (1 + o(1)) \ 0.254n^{3/2}.$$

Roughly speaking, our main idea in proving the above theorem is to analyze the structure of a Berge- C_5 -free hypergraph, and use this structure to efficiently bound the number of paths of length 3 that start from certain dense subgraphs (e.g., triangle, K_4) of the 2-shadow. This bound is then combined with the lower bound on the number of paths of length 3 provided by the Blakley-Roy inequality [2]. We prove Theorem 3 in Section 2.

Ergemlidze, Győri and Methuku [3] considered the analogous question for linear hypergraphs and proved that $ex_3(n, \{C_2, C_5\}) = n^{3/2}/3\sqrt{3} + o(n^{3/2})$. Surprisingly, even though their lower bound is the same as the lower bound in Theorem 2, the linear hypergraph that they constructed in [3] is very different from the hypergraph used in the Bollobás-Győri example discussed above – the latter is far from being linear. In [3], the authors also strengthened Theorem 1 by showing that $ex_3(n, \{C_2, C_3, C_4\}) \sim ex_3(n, \{C_2, C_4\})$. Recently, $ex_3(n, C_4)$ was studied in [5]. See [6] for results on the maximum number of hyperedges in an *r*-uniform hypergraph of girth six.

Győri and Lemons [16, 17] generalized Theorem 2 to Berge cycles of any given length and proved bounds on $ex_r(n, C_{2k+1})$ and $ex_r(n, C_{2k})$. These bounds were improved by Füredi and Özkahya [9], Jiang and Ma [19], Gerbner, Methuku and Vizer [11]. Recently Füredi, Kostochka and Luo [7] started the study of the maximum size of an *n*-vertex *r*-uniform hypergraph without any Berge cycle of length at least *k*. This study has been continued in [8, 18, 20, 4].

General results for Berge-F-free hypergraphs have been obtained in [12, 13, 10] and the Turán numbers of Berge- $K_{2,t}$ and Berge cliques, among others, were studied in [24, 22, 11, 14, 10].

Notation

We introduce some important notations and definitions used throughout the paper.

- Length of a path is the number of edges in the path. We usually denote a path v_0, v_1, \ldots, v_k , simply as $v_0v_1 \ldots v_k$.
- For convenience, an edge $\{a, b\}$ of a graph or a pair of vertices a, b is referred to as ab. A hyperedge $\{a, b, c\}$ is written simply as abc.
- For a hypergraph H (or a graph G), for convenience, we sometimes use H (or G) to denote the edge set of the hypergraph H (or G respectively). Thus the number of edges in H is |H|.
- Given a graph G and a subset of its vertices S, let the subgraph of G induced by S be denoted by G[S].
- For a hypergraph H, let $\partial H = \{ab \mid ab \subset e \in E(H)\}$ denote its 2-shadow graph.
- For a hypergraph H, the *neighborhood* of v in H is defined as

$$N(v) = \{x \in V(H) \setminus \{v\} \mid v, x \in h \text{ for some } h \in E(H)\}.$$

• For a hypergraph H and a pair of vertices $u, v \in V(H)$, let codeg(v, u) denote the number of hyperedges of H containing the pair $\{u, v\}$.

2 Proof of Theorem 3

Let H be a hypergraph on n vertices without a Berge 5-cycle and let $G = \partial H$ be the 2-shadow of H. First we introduce some definitions.

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Definition 4. A pair $xy \in \partial H$ is called *thin* if $\operatorname{codeg}(xy) = 1$, otherwise it is called *fat*. We say a hyperedge $abc \in H$ is *thin* if at least two of the pairs ab, bc, ac are thin.

Definition 5. We say a set of hyperedges (or a hypergraph) is tightly-connected if it can be obtained by starting with a hyperedge and adding hyperedges one by one, such that every added hyperedge intersects with one of the previous hyperedges in 2 vertices.

Definition 6. A *block* in *H* is a maximal set of tightly-connected hyperedges.

Definition 7. For a block B, a maximal subhypergraph of B without containing thin hyperedges is called the *core* of the block.

Let K_4^3 denote the complete 3-uniform hypergraph on 4 vertices. A crown of size k is a set of $k \ge 1$ hyperedges of the form $abc_1, abc_2, \ldots, abc_k$. Below we define 2 specific hypergraphs:

- Let F_1 be a hypergraph consisting of exactly 3 hyperedges on 4 vertices (i.e., K_4^3 minus an edge).
- For distinct vertices a, b, c, d and o, let F_2 be the hypergraph consisting of hyperedges oab, obc, ocd and oda.

Lemma 8. Let B be a block of H, and let \mathcal{B} be a core of B. Then \mathcal{B} is either \emptyset , K_4^3 , F_1 , F_2 or a crown of size k for some $k \ge 1$.

Proof. If $\mathcal{B} = \emptyset$, we are done, so let us assume $\mathcal{B} \neq \emptyset$. Since B is tightly-connected and it can be obtained by adding thin hyperedges to \mathcal{B} , it is easy to see that \mathcal{B} is also tightly-connected. Thus if \mathcal{B} has at most two hyperedges, then it is a crown of size 1 or 2 and we are done. Therefore, in the rest of the proof we will assume that \mathcal{B} contains at least 3 hyperedges.

If \mathcal{B} contains at most 4 vertices then it is easy to see that \mathcal{B} is either K_4^3 or F_1 . So assume that \mathcal{B} has at least 5 vertices (and at least 3 hyperedges). Since \mathcal{B} is not a crown, there exists a tight path of length 3, say *abc*, *bcd*, *cde*. Since *abc* is in the core, one of the pairs *ab* or *ac* is fat, so there exists a hyperedge $h \neq abc$ containing either *ab* or *ac*. Similarly there exists a hyperedge $f \neq cde$ and f contains *ed* or *ec*. If h = f then $\mathcal{B} \supseteq F_2$. However, it is easy to see that F_2 cannot be extended to a larger tightly-connected set of hyperedges without creating a Berge 5-cycle, so in this case $\mathcal{B} = F_2$. If $h \neq f$ then the hyperedges h, *abc*, *bcd*, *cde*, f create a Berge 5-cycle in H, a contradiction. This completes the proof of the lemma.

Observation 9. Let B be a block of H and let \mathcal{B} be the core of B. If $\mathcal{B} = \emptyset$ then the block B is a crown, and if $\mathcal{B} \neq \emptyset$ then every fat pair of B is contained in $\partial \mathcal{B}$.

Edge Decomposition of $G = \partial H$. We define a decomposition \mathcal{D} of the edges of G into paths of length 2, triangles and K_4 's such as follows: Let B be a block of H and \mathcal{B} be its core. If $\mathcal{B} = \emptyset$, then B is a crown-block $\{abc_1, abc_2, \dots, abc_k\}$ (for some $k \ge 1$); we partition ∂B into the triangle abc_1 and paths ac_ib where $2 \le i \le k$.

If $\mathcal{B} \neq \emptyset$, then our plan is to first partition $\partial B \setminus \partial \mathcal{B}$. If $abc \in B \setminus \mathcal{B}$, then abc is a thin hyperedge, so it contains at least 2 thin pairs, say ab and bc. We claim that the pair ac is in $\partial \mathcal{B}$. Indeed, ac has to be a fat pair, otherwise the block B consists of only one hyperedge abc, so $\mathcal{B} = \emptyset$ contradicting the assumption. So by Observation 9, achas to be a pair in $\partial \mathcal{B}$. For every $abc \in B \setminus \mathcal{B}$ such that ab and bc are thin pairs, add the 2-path abc to the edge decomposition \mathcal{D} . This partitions all the edges in $\partial B \setminus \partial \mathcal{B}$ into paths of length 2. So all we have left is to partition the edges of $\partial \mathcal{B}$.

- If \mathcal{B} is a crown $\{abc_1, abc_2, \ldots, abc_k\}$ for some $k \ge 1$, then we partition ∂B into the triangle abc_1 and paths ac_ib where $2 \le i \le k$.
- If $\mathcal{B} = F_1 = \{abc, bcd, acd\}$ then we partition $\partial \mathcal{B}$ into 2-paths abc, bdc and cad.
- If $\mathcal{B} = F_2 = \{oab, obc, ocd, oda\}$ then we partition $\partial \mathcal{B}$ into 2-paths abo, bco, cdo and dao.
- Finally, if $\mathcal{B} = K_4^3 = \{abc, abd, acd, bcd\}$ then we partition $\partial \mathcal{B}$ as K_4 , i.e., we add $\partial \mathcal{B} = K_4$ as an element of \mathcal{D} .

Clearly, by Lemma 8 we have no other cases left. Thus all of the edges of the graph G are partitioned into paths of length 2, triangles and K_4 's.

Observation 10.

- (a) If D is a triangle that belongs to \mathcal{D} , then there is a hyperedge $h \in H$ such that $D = \partial h$.
- (b) If abc is a 2-path that belongs to \mathcal{D} , then $abc \in H$. Moreover ac is a fat pair.
- (c) If D is a K_4 that belongs to \mathcal{D} , then there exists $F = K_4^3 \subseteq H$ such that $D = \partial F$.

Let $\alpha_1 |G|$ and $\alpha_2 |G|$ be the number of edges of G that are contained in triangles and 2-paths of the edge-decomposition \mathcal{D} of G, respectively. So $(1 - \alpha_1 - \alpha_2) |G|$ edges of G belong to the K_4 's in \mathcal{D} .

Claim 11. We have,

$$|H| = \left(\frac{\alpha_1}{3} + \frac{\alpha_2}{2} + \frac{2(1 - \alpha_1 - \alpha_2)}{3}\right) |G|.$$

Proof. Let B be a block with the core \mathcal{B} . Recall that for each hyperedge $h \in B \setminus \mathcal{B}$, we have added exactly one 2-path or a triangle to \mathcal{D} .

Moreover, because of the way we partitioned $\partial \mathcal{B}$, it is easy to check that in all of the cases except when $\mathcal{B} = K_4^3$, the number of hyperedges of \mathcal{B} is the same as the number of elements of \mathcal{D} that $\partial \mathcal{B}$ is partitioned into; these elements being 2-paths and triangles. On the other hand, if $\mathcal{B} = K_4^3$, then the number of hyperedges of \mathcal{B} is 4 but we added only one element to \mathcal{D} (namely K_4).

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This shows that the number of hyperedges of H is equal to the number of elements of \mathcal{D} that are 2-paths or triangles plus the number of hyperedges which are in copies of K_4^3 in H, i.e., 4 times the number of K_4 's in \mathcal{D} . Since $\alpha_1 |G|$ edges of G are in 2-paths, the number of elements of \mathcal{D} that are 2-paths is $\alpha_1 |G|/2$. Similarly, the number of elements of \mathcal{D} that are triangles is $\alpha_2 |G|/3$, and the number of K_4 's in \mathcal{D} is $(1 - \alpha_1 - \alpha_2) |G|/6$. Combining this with the discussion above finishes the proof of the claim.

The link of a vertex v is the graph consisting of the edges $\{uw \mid uvw \in H\}$ and is denoted by L_v .

Claim 12. $|L_v| \leq 2 |N(v)|$.

Proof. First let us notice that there is no path of length 5 in L_v . Indeed, otherwise, there exist vertices v_0, v_1, \ldots, v_5 such that $vv_{i-1}v_i \in H$ for each $1 \leq i \leq 5$ which means there is a Berge 5-cycle in H formed by the hyperedges containing the pairs $vv_1, v_1v_2, v_2v_3, v_3v_4, v_4v$, a contradiction. So by the Erdős-Gallai theorem $|L_v| \leq \frac{5-1}{2} |N(v)|$, proving the claim. \Box

Lemma 13. Let $v \in V(H)$ be an arbitrary vertex, then the number of edges in G[N(v)] is less than 8 |N(v)|.

Proof. Let G_v be a subgraph of G on a vertex set N(v), such that $xy \in G_v$ if and only if there exists a vertex $z \neq v$ such that $xyz \in H$. Then each edge of G[N(v)] belongs to either L_v or G_v , so $|G[N(v)]| \leq |L_v| + |G_v|$. Combining this with Claim 12, we get $|G[N(v)]| \leq |G_v| + 2 |N(v)|$. So it suffices to prove that $|G_v| < 6 |N(v)|$.

First we will prove that there is no path of length 12 in G_v . Let us assume by contradiction that $P = v_0, v_1, \ldots, v_{12}$ is a path in G_v . Since for each pair of vertices v_i, v_{i+1} , there is a hyperedge $v_i v_{i+1} x$ in H where $x \neq v$, we can conclude that there is a subsequence u_0, u_1, \ldots, u_6 of v_0, v_1, \ldots, v_{12} and a sequence of distinct hyperedges h_1, h_2, \ldots, h_6 , such that $u_{i-1}u_i \subset h_i$ and $v \notin h_i$ for each $1 \leq i \leq 6$. Since $u_0, u_3, u_6 \in N(v)$ there exist hyperedges $f_1, f_2, f_3 \in H$ such that $vu_0 \subset f_1, vu_3 \subset f_2$ and $vu_6 \subset f_3$. Clearly, either $f_1 \neq f_2$ or $f_2 \neq f_3$. In the first case the hyperedges f_1, h_1, h_2, h_3, f_2 , and in the second case the hyperedges f_2, h_4, h_5, h_6, f_3 form a Berge 5-cycle in H, a contradiction.

Therefore, there is no path of length 12 in G_v , so by the Erdős-Gallai theorem, the number of edges in G_v is at most $\frac{12-1}{2}|N(v)| < 6|N(v)|$, as required.

2.1 Relating the hypergraph degree to the degree in the shadow

For a vertex $v \in V(H) = V(G)$, let d(v) denote the degree of v in H and let $d_G(v)$ denote the degree of v in G (i.e., $d_G(v)$ is the degree in the shadow).

Clearly $d_G(v) \leq 2d(v)$. Moreover, $d(v) = |L_v|$ and $d_G(v) = |N(v)|$. So by Claim 12, we have

$$\frac{d_G(v)}{2} \leqslant d(v) \leqslant 2d_G(v). \tag{1}$$

Let \overline{d} and \overline{d}_G be the average degrees of H and G respectively.

Suppose there is a vertex v of H, such that $d(v) < \overline{d}/3$. Then we may delete v and all the edges incident to v from H to obtain a graph H' whose average degree is more than

 $3(n\overline{d}/3 - \overline{d}/3)/(n-1) = \overline{d}$. Then it is easy to see that if the theorem holds for H', then it holds for H as well. Repeating this procedure, we may assume that for every vertex v of H, $d(v) \ge \overline{d}/3$. Therefore, by (1), we may assume that the degree of every vertex of G is at least $\overline{d}/6$.

2.2 Counting paths of length 3

Definition 14. A 2-path in ∂H is called *bad* if both of its edges are contained in a triangle of ∂H , otherwise it is called *good*.

Lemma 15. For any vertex $v \in V(G)$ and a set $M \subseteq N(v)$, let \mathcal{P} be the set of the good 2-paths vxy such that $x \in M$. Let $M' = \{y \mid vxy \in \mathcal{P}\}$ then $|\mathcal{P}| < 2|M'| + 48d_G(v)$.

Proof. Let $B_{\mathcal{P}} = \{xy \mid x \in M, y \in M', xy \in G\}$ be a bipartite graph, clearly $|B_{\mathcal{P}}| = |\mathcal{P}|$. Let $E = \{xyz \in H \mid x, y \in N(v), \operatorname{codeg}(x, y) \leq 2\}$. By Lemma 13, $|E| \leq 2 \cdot 8 |N(v)|$ so the number of edges of 2-shadow of E is $|\partial E| \leq 48 |N(v)|$. Let $B = \{xy \in B_{\mathcal{P}} \mid \exists z \in V(H), xyz \in H \setminus E\}$. Then clearly,

$$|B| \ge |B_{\mathcal{P}}| - |\partial E| \ge |\mathcal{P}| - 48 |N(v)| = |\mathcal{P}| - 48 d_G(v).$$
⁽²⁾

Let $d_B(x)$ denote the degree of a vertex x in the graph B.

Claim 16. For every $y \in M'$ such that $d_B(y) = k \ge 3$, there exists a set of k - 2 vertices $S_y \subseteq M'$ such that $\forall w \in S_y$ we have $d_B(w) = 1$. Moreover, $S_y \cap S_z = \emptyset$ for any $y \ne z \in M'$ (with $d_B(y), d_B(z) \ge 3$).

Proof. Let $yx_1, yx_2, \ldots, yx_k \in B$ be the edges of B incident to y. For each $1 \leq j \leq k$ let $f_j \in H$ be a hyperedge such that $vx_j \subset f_j$. For each $yx_i \in B$ clearly there is a hyperedge $yx_iw_i \in H \setminus E$.

We claim that for each $1 \leq i \leq k$, $w_i \in M'$. It is easy to see that $w_i \in N(v)$ or $w_i \in M'$ (because vx_iw_i is a 2-path in G). Assume for a contradiction that $w_i \in N(v)$, then since $yx_iw_i \notin E$ we have, $\operatorname{codeg}(x_i, w_i) \geq 3$. Let $f \in H$ be a hyperedge such that $vw_i \subset f$. Now take $j \neq i$ such that $x_j \neq w_i$. If $f_j \neq f$ then since $\operatorname{codeg}(x_i, w_i) \geq 3$ there exists a hyperedge $h \supset x_iw_i$ such that $h \neq f$ and $h \neq x_iw_iy$, then the hyperedges $f, h, x_iw_iy, yx_jw_j, f_j$ form a Berge 5-cycle. So $f_j = f$, therefore $f_j \neq f_i$. Similarly in this case, there exists a hyperedge $h \supset x_iw_i$ such that $h \neq f_i$ and $h \neq x_iw_iy$, therefore the hyperedges $f_i, h, x_iw_iy, yx_jw_j, f_j$ form a Berge 5-cycle, a contradiction. So we proved that $w_i \in M'$ for each $1 \leq i \leq k$.

Claim 17. For all but at most 2 of the w_i 's (where $1 \leq i \leq k$), we have $d_B(w_i) = 1$.

Proof. If $d_B(w_i) = 1$ for all $1 \le i \le k$ then we are done, so we may assume that there is $1 \le i \le k$ such that $d_B(w_i) \ne 1$.

For each $1 \leq i \leq k$, $w_i \in M'$ and $x_i w_i \in \partial(H \setminus E)$ (because $x_i w_i y \in H \setminus E$), so it is clear that $d_B(w_i) \geq 1$. So $d_B(w_i) > 1$. Then there is a vertex $x \in M \setminus \{x_i\}$ such that $w_i x \in B$. Let $f, h \in H$ be hyperedges with $w_i x \in h$ and $xv \in f$. If there are $j, l \in \{1, 2, \ldots, k\} \setminus \{i\}$ such that x, x_j and x_l are all different from each other, then clearly, either $f \neq f_j$ or $f \neq f_l$, so without loss of generality we may assume $f \neq f_j$. Then the hyperedges $f, h, w_i x_i y, y w_j x_j, f_j$ create a Berge cycle of length 5, a contradiction. So there are no $j, l \in \{1, 2, ..., k\} \setminus \{i\}$ such that x, x_j and x_l are all different from each other. Clearly this is only possible when k < 4 and there is a $j \in \{1, 2, 3\} \setminus \{i\}$ such that $x = x_j$. Let $l \in \{1, 2, 3\} \setminus \{i, j\}$. If $f_j \neq f_l$ then the hyperedges $f_j, h, w_i x_i y, y w_l x_l, f_l$ form a Berge 5-cycle. Therefore $f_j = f_l$. So we proved that $d_B(w_i) \neq 1$ implies that k = 3 and for $\{j, l\} = \{1, 2, 3\} \setminus \{i\}$, we have $f_j = f_l$. So if $d_B(w_i) \neq 1$ and $d_B(w_j) \neq 1$ we have $f_j = f_l$ and $f_i = f_l$, which is impossible. So $d_B(w_j) = 1$. So we proved that if for any $1 \leq i \leq k, d_B(w_i) \neq 1$ then k = 3 and all but at most 2 of the vertices in $\{w_1, w_2, w_3\}$ have degree 1 in the graph B, as desired.

We claim that for any $i \neq j$ where $d_B(w_i) = d_B(w_j) = 1$ we have $w_i \neq w_j$. Indeed, if there exists $i \neq j$ such that $w_i = w_j$ then $w_i x_j$ and $w_i x_i$ are both adjacent to w_i in the graph B which contradicts to $d_B(w_i) = 1$. So using the above claim, we conclude that the set $\{w_1, w_2, \ldots, w_k\}$ contains at least k - 2 distinct elements with each having degree one in the graph B, so we can set S_y to be the set of these k - 2 elements. (Then of course $\forall w_i \in S_y$ we have $d_B(w_i) = 1$.)

Now we have to prove that for each $z \neq y$ we have $S_y \cap S_z = \emptyset$. Assume by contradiction that $w_i \in S_z \cap S_y$ for some $z \neq y$. That is, there is some hyperedge $uw_i z \in H \setminus E$ where $u \in M$, moreover $u = x_i$ otherwise $d_B(w_i) > 1$. So we have a hyperedge $x_i w_i z \in H \setminus E$ for some $z \in M' \setminus \{y\}$. Let $j, l \in \{1, 2, ..., k\} \setminus \{i\}$ such that $j \neq l$. Recall that $x_j v \subset f_j$ and $x_l v \subset f_l$. Clearly either $f_j \neq f_i$ or $f_l \neq f_i$ so without loss of generality we can assume $f_j \neq f_i$. Then it is easy to see that the hyperedges $f_j, x_j w_j y, y x_i w_i, w_i z x_i, f_i$ are all different and they create a Berge 5-cycle $(x_j w_j y \neq y x_i w_i$ because $x_j \neq w_i$).

For each $x \in M'$ with $d_B(x) = k \ge 3$, let S_x be defined as in Claim 16. Then the average of the degrees of the vertices in $S_x \cup \{x\}$ in B is $(k+|S_x|)/(k-1) = (2k-2)(k-1) = 2$. Since the sets $S_x \cup x$ (with $x \in M'$, $d_B(x) \ge 3$) are disjoint, we can conclude that average degree of the set M' is at most 2. Therefore $2|M'| \ge |B|$. So by (2) we have $2|M'| \ge |B| > |\mathcal{P}| - 48d_G(V)$, which completes the proof of the lemma. \Box

Claim 18. We may assume that the maximum degree in the graph G is less than $160\sqrt{n}$ when n is large enough.

Proof. Let v be an arbitrary vertex with $d_G(v) = C\overline{d}$ for some constant C > 0. Let \mathcal{P} be the set of the good 2-paths starting from the vertex v. Then applying Lemma 15 with M = N(v) and $M' = \{y \mid vxy \in \mathcal{P}\}$, we have $|\mathcal{P}| < 2 |M'| + 48d_G(v) < 2n + 48 \cdot C\overline{d}$. Since the minimum degree in G is at least $\overline{d}/6$, the number of (ordered) 2-paths starting from v is at least $d(v) \cdot (\overline{d}/6 - 1) = C\overline{d} \cdot (\overline{d}/6 - 1)$. Notice that the number of (ordered) bad 2-paths starting at v is the number of 2-paths vxy such that $x, y \in N(v)$. So by Lemma 13, this is at most $2 \cdot 8 |N(v)| = 16C\overline{d}$, so the number of good 2-paths is at least $C\overline{d} \cdot (\overline{d}/6 - 17)$. So $|\mathcal{P}| \ge C\overline{d} \cdot (\overline{d}/6 - 17)$. Thus we have

$$C\overline{d} \cdot (\overline{d}/6 - 17) \leq |\mathcal{P}| < 2n + 48C\overline{d}.$$

So $C\overline{d}(\overline{d}/6 - 65) < 2n$. Therefore, $6C(\overline{d}/6 - 65)^2 < 2n$, i.e., $\overline{d} < 6\sqrt{n/3C} + 390$, so $|H| = n\overline{d}/3 \leq 2n\sqrt{n/3C} + 130n$. If $C \geq 36$ we get that $|H| \leq \frac{n^{3/2}}{3\sqrt{3}} + 130n = \frac{n^{3/2}}{3\sqrt{3}} + O(n)$, proving Theorem 3. So we may assume C < 36.

Theorem 2 implies that

$$|H| = n\bar{d}/3 \leqslant \sqrt{2}n^{3/2} + 4.5n,\tag{3}$$

so $\overline{d} \leq 3\sqrt{2}\sqrt{n} + 13.5$. So combining this with the fact that C < 36, we have $d_G(v) = C\overline{d} < 108\sqrt{2}\sqrt{n} + 486 < 160\sqrt{n}$ for large enough n.

Combining Lemma 15 and Claim 18, we obtain the following.

Lemma 19. For any vertex $v \in V(G)$ and a set $M \subseteq N(v)$, let \mathcal{P} be the set of good 2-paths vxy such that $x \in M$. Let $M' = \{y \mid vxy \in \mathcal{P}\}$ then $|\mathcal{P}| < 2|M'| + 7680\sqrt{n}$ when n is large enough.

Definition 20. A 3-path x_0, x_1, x_2, x_3 is called *good* if both 2-paths x_0, x_1, x_2 and x_1, x_2, x_3 are good 2-paths.

Claim 21. The number of (ordered) good 3-paths in G is at least $n\overline{d}_G^3 - C_0 n^{3/2}\overline{d}_G$ for some constant $C_0 > 0$ (for large enough n).

Proof. First we will prove that the number of (ordered) 3-walks that are not good 3-paths is at most $5440n^{3/2}\overline{d}_G$.

For any vertex $x \in V(H)$ if a path yxz is a bad 2-path then zy is an edge of G, so the number of (ordered) bad 2-paths whose middle vertex is x, is at most 2 times the number of edges in G[N(x)], which is less than $2 \cdot 8 |N(x)| = 16d_G(x)$ by Lemma 13. The number of 2-walks which are not 2-paths and whose middle vertex is x is exactly $d_G(x)$. So the total number of (ordered) 2-walks that are not good 2-paths is at most $\sum_{x \in V(H)} 17d_G(x) = 17n\overline{d}_G$.

Notice that, by definition, any (ordered) 3-walk that is not a good 3-path must contain a 2-walk that is not a good 2-path. Moreover, if xyz is a 2-walk that is not a good 2-path, then the number of 3-walks in G containing it is at most $d_G(x) + d_G(z) < 320\sqrt{n}$ (for large enough n) by Claim 18. Therefore, the total number of (ordered) 3-walks that are not good 3-paths is at most $17n\overline{d}_G \cdot 320\sqrt{n} = 5440n^{3/2}\overline{d}_G$.

By the Blakley-Roy inequality, the total number of (ordered) 3-walks in G is at least $n\overline{d}_G^3$. By the above discussion, all but at most $5440n^{3/2}\overline{d}_G$ of them are good 3-paths, so letting $C_0 = 5440$ completes the proof of the claim.

Claim 22. Let $\{a, b, c\}$ be the vertex set of a triangle that belongs to \mathcal{D} . (By Observation 10 (a) $abc \in H$.) Then the number of good 3-paths whose first edge is ab, bc or ca is at most $8n + C_1\sqrt{n}$ for some constant C_1 and for large enough n.

Proof. For each $\{x, y\} \subset \{a, b, c\}$, let $S_{xy} = N(x) \cap N(y) \setminus \{a, b, c\}$. For each $x \in \{a, b, c\}$, let $S_x = N(x) \setminus (N(y) \cup N(z) \cup \{a, b, c\})$ where $\{y, z\} = \{a, b, c\} \setminus \{x\}$.

For each $x \in \{a, b, c\}$, let \mathcal{P}_x be the set of good 2-paths xuv where $u \in S_x$. Let $S'_x = \{v \mid xuv \in \mathcal{P}_x\}$. For each $\{x, y\} \subset \{a, b, c\}$, let \mathcal{P}_{xy} be the set of good 2-paths xuv and yuv where $u \in S_{xy}$. Let $S'_{xy} = \{v \mid xuv \in \mathcal{P}_{xy}\}$.

Let $\{x, y\} \subset \{a, b, c\}$ and $z = \{a, b, c\} \setminus \{x, y\}$. Notice that each 2-path $yuv \in \mathcal{P}_{xy}$ ($xuv \in \mathcal{P}_{xy}$), is contained in at most one good 3-path zyuv (respectively zxuv) whose first edge is in the triangle abc. Indeed, since $u \in S_{xy}$, xyuv (respectively yxuv) is not a good 3-path. Therefore, the number of good 3-paths whose first edge is in the triangle abc, and whose third vertex is in S_{xy} is at most $|\mathcal{P}_{xy}|$. The number of paths in \mathcal{P}_{xy} that start with the vertex x is less than $2|S'_{xy}| + 7680\sqrt{n}$, by Lemma 19. Similarly, the number of paths in \mathcal{P}_{xy} that start with the vertex y is less than $2|S'_{xy}| + 7680\sqrt{n}$. Since every path in \mathcal{P}_{xy} starts with either x or y, we have $|\mathcal{P}_{xy}| < 4|S'_{xy}| + 15360\sqrt{n}$. Therefore, for any $\{x, y\} \subset \{a, b, c\}$, the number of good 3-paths whose first edge is in the triangle abc, and whose third vertex is in S_{xy} is less than $4|S'_{xy}| + 15360\sqrt{n}$.

In total, the number of good 3-paths whose first edge is in the triangle abc and whose third vertex is in $S_{ab} \cup S_{bc} \cup S_{ac}$ is at most

$$4(|S'_{ab}| + |S'_{bc}| + |S'_{ac}|) + 46080\sqrt{n}.$$
(4)

Let $x \in \{a, b, c\}$ and $\{y, z\} = \{a, b, c\} \setminus \{x\}$. For any 2-path $xuv \in \mathcal{P}_x$ there are 2 good 3-paths with the first edge in the triangle *abc*, namely yxuv and zxuv. So the total number of 3-paths whose first edge is in the triangle *abc* and whose third vertex is in $S_a \cup S_b \cup S_c$ is $2(|\mathcal{P}_a| + |\mathcal{P}_b| + |\mathcal{P}_c|)$, which is at most

$$4(|S'_a| + |S'_b| + |S'_c|) + 46080\sqrt{n},\tag{5}$$

by Lemma 19.

Now we will prove that every vertex is in at most 2 of the sets $S'_a, S'_b, S'_c, S'_{ab}, S'_{bc}, S'_{ac}$. Let us assume by contradiction that a vertex $v \in V(G) \setminus \{a, b, c\}$ is in at least 3 of them. We claim that there do not exist 3 vertices $u_a \in N(a) \setminus \{b, c\}, u_b \in N(b) \setminus \{a, c\}$ and $u_c \in N(c) \setminus \{a, b\}$ such that $xu_x v$ is a good 3-path for each $x \in \{a, b, c\}$. Indeed, otherwise, consider hyperedges h_a, h'_a containing the pairs au_a and $u_a v$ respectively (since au_av is a good 2-path, note that $h_a \neq h'_a$, and hyperedges h_b, h'_b, h_c, h'_c containing the pairs bu_b, u_bv, cu_c, u_cv respectively. Then either $h'_a \neq h'_b$ or $h'_a \neq h'_c$, say $h'_a \neq h'_b$ without loss of generality. Then the hyperedges $h_a, h'_a, h'_b, h_b, abc$ create a Berge 5-cycle in H, a contradiction, proving that it is impossible to have 3 vertices $u_a \in N(a) \setminus \{b, c\}, u_b \in$ $N(b) \setminus \{a, c\}$ and $u_c \in N(c) \setminus \{a, b\}$ with the above mentioned property. Without loss of generality let us assume that there is no vertex $u_a \in N(a) \setminus \{b, c\}$ such that $au_a v$ is a good 2-path – in other words, $v \notin S'_a \cup S'_{ab} \cup S'_{ac}$. However, since we assumed that v is contained in at least 3 of the sets $S'_a, S'_b, S'_c, S'_{ab}, S'_{bc}, S'_{ac}$, we can conclude that v is contained in all 3 of the sets S'_b , S'_c , S'_{bc} , i.e., there are vertices $u_b \in S_b$, $u_c \in S_c$, $u \in S_{bc}$ such that vu_bb , vu_cc , vub, vuc are good 2-paths. Using a similar argument as before, if $vu \in h, vu_b \in h_b$ and $vu_c \in h_c$, without loss of generality we can assume that $h \neq h_b$, so the hyperedges abc, h, h_b together with hyperedges containing uc and $u_b b$ form a Berge 5-cycle in H, a contradiction.

So we proved that

$$2 \left| S'_a \cup S'_b \cup S'_c \cup S'_{ab} \cup S'_{bc} \cup S'_{ac} \right| \ge \left| S'_a \right| + \left| S'_b \right| + \left| S'_c \right| + \left| S'_{ab} \right| + \left| S'_{bc} \right| + \left| S'_{ac} \right|$$

This together with (4) and (5), we get that the number of good 3-paths whose first edge is in the triangle *abc* is at most

$$8 |S'_a \cup S'_b \cup S'_c \cup S'_{ab} \cup S'_{bc} \cup S'_{ac}| + 92160\sqrt{n} < 8n + C_1\sqrt{n}$$

for $C_1 = 92160$ and large enough n, finishing the proof of the claim.

Claim 23. Let P = abc be a 2-path and $P \in \mathcal{D}$. (By Observation 10 (b) $abc \in H$.) Then the number of good 3-paths whose first edge is ab or bc is at most $4n + C_2\sqrt{n}$ for some constant $C_2 > 0$ and large enough n.

Proof. First we bound the number of 3-paths whose first edge is ab. Let $S_{ab} = N(a) \cap N(b)$. Let $S_a = N(a) \setminus (N(b) \cup \{b\})$ and $S_b = N(b) \setminus (N(a) \cup \{a\})$. For each $x \in \{a, b\}$, let \mathcal{P}_x be the set of good 2-paths xuv where $u \in S_x$, and let $S'_x = \{v \mid xuv \in \mathcal{P}_x\}$. The set of good 3-paths whose first edge is ab is $\mathcal{P}_a \cup \mathcal{P}_b$, because the third vertex of a good 3-path starting with an edge ab can not belong to $N(a) \cap N(b)$ by the definition of a good 3-path.

We claim that $|S'_a \cap S'_b| \leq 160\sqrt{n}$. Let us assume by contradiction that $v_0, v_1, \ldots, v_k \in S'_a \cap S'_b$ for $k > 160\sqrt{n}$. For each vertex v_i where $0 \leq i \leq k$, there are vertices $a_i \in S_a$ and $b_i \in S_b$ such that aa_iv_i, bb_iv_i are good 2-paths. For each $0 \leq i \leq k$, the hyperedge $a_iv_ib_i$ is in H, otherwise we can find distinct hyperedges containing the pairs $aa_i, a_iv_i, v_ib_i, b_ib$ and these hyperedges together with abc, would form a Berge 5-cycle in H, a contradiction. We claim that there are $j, l \in \{0, 1, \ldots, k\}$ such that $a_j \neq a_l$, otherwise there is a vertex x such that $x = a_i$ for each $0 \leq i \leq k$. Then $xv_i \in G$ for each $0 \leq i \leq k$, so we get that $d_G(x) > k > 160\sqrt{n}$ which contradicts Claim 18.

So there are $j, l \in \{0, 1, ..., k\}$ such that $a_j \neq a_l$ and $a_j v_j b_j, a_l v_l b_l \in H$. By observation 10 (b), there is a hyperedge $h \neq abc$ such that $ac \subset h$. Clearly either $a_j \notin h$ or $a_l \notin h$. Without loss of generality let $a_j \notin h$, so there is a hyperedge h_a with $aa_j \subset h_a \neq h$. Let $h_b \supset b_j b$, then the hyperedges $abc, h, h_a, a_j v_j b_j, h_b$ form a Berge 5-cycle, a contradiction, proving that $|S'_a \cap S'_b| \leq 160\sqrt{n}$.

Notice that $|S'_a| + |S'_b| = |S'_a \cup S'_b| + |S'_a \cap S'_b| \leq n + 160\sqrt{n}$. So by Lemma 19, we have

$$|\mathcal{P}_a| + |\mathcal{P}_b| \leq 2(|S'_a| + |S'_b|) + 2 \cdot 7680\sqrt{n} \leq 2(n + 160\sqrt{n}) + 2 \cdot 7680\sqrt{n} = 2n + 15680\sqrt{n}$$

for large enough n. So the number of good 3-paths whose first edge is ab is at most $2n + 15680\sqrt{n}$. By the same argument, the number of good 3-paths whose first edge is bc is at most $2n + 15680\sqrt{n}$. Their sum is at most $4n + C_2\sqrt{n}$ for $C_2 = 31360$ and large enough n, as desired.

Claim 24. Let $\{a, b, c, d\}$ be the vertex set of a K_4 that belongs to \mathcal{D} . Let $F = K_4^3$ be a hypergraph on the vertex set $\{a, b, c, d\}$. (By Observation 10 (c) $F \subseteq H$.) Then the number of good 3-paths whose first edge belongs to ∂F is at most $6n + C_3\sqrt{n}$ for some constant $C_3 > 0$ and large enough n.

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Proof. First, let us observe that there is no Berge path of length 2, 3 or 4 between distinct vertices $x, y \in \{a, b, c, d\}$ in the hypergraph $H \setminus F$, because otherwise this Berge path together with some edges of F will form a Berge 5-cycle in H. This implies, that there is no path of length 3 or 4 between x and y in $G \setminus \partial F$, because otherwise we would find a Berge path of length 2, 3 or 4 between x and y in $H \setminus F$.

Let $S = \{u \in V(H) \setminus \{a, b, c, d\} \mid \exists \{x, y\} \subset \{a, b, c, d\}, u \in N(x) \cap N(y)\}$. For each $x \in \{a, b, c, d\}$, let $S_x = N(x) \setminus (S \cup \{a, b, c, d\})$. Let \mathcal{P}_S be the set of good 2-paths xuv where $x \in \{a, b, c, d\}$ and $u \in S$. Let $S' = \{v \mid xuv \in \mathcal{P}_S\}$. For each $x \in \{a, b, c, d\}$, let \mathcal{P}_x be the set of good 2-paths xuv where $u \in S_x$, and let $S'_x = \{v \mid xuv \in \mathcal{P}_x\}$.

Let $v \in S'$. By definition, there exists a pair of vertices $\{x, y\} \subset \{a, b, c, d\}$ and a vertex u, such that xuv and yuv are good 2-paths.

Suppose that zu'v is a 2-path different from xuv and yuv where $z \in \{a, b, c, d\}$. If u' = u then $z \notin \{x, y\}$ so there is a Berge 2-path between x and y or between x and z in $H \setminus F$, which is impossible. So $u \neq u'$. Either $z \neq x$ or $z \neq y$, without loss of generality let us assume that $z \neq x$. Then zu'vux is a path of length 4 in $G \setminus \partial F$, a contradiction. So for any $v \in S'$ there are only 2 paths of $\mathcal{P}_a \cup \mathcal{P}_b \cup \mathcal{P}_c \cup \mathcal{P}_d \cup \mathcal{P}_S$ that contain v as an end vertex – both of which are in \mathcal{P}_S – which means that $v \notin S'_a \cup S'_b \cup S'_c \cup S'_d$, so $S' \cap (S'_a \cup S'_b \cup S'_c \cup S'_d) = \emptyset$. Moreover,

$$|\mathcal{P}_S| \leqslant 2 |S'| \,. \tag{6}$$

We claim that S'_a and S'_b are disjoint. Indeed, otherwise, if $v \in S'_a \cap S'_b$ there exists $x \in S_a$ and $y \in S_b$ such that vxa and vyb are paths in G, so there is a 4-path axvyb between vertices of F in $G \setminus \partial F$, a contradiction. Similarly we can prove that S'_a, S'_b, S'_c and S'_d are pairwise disjoint. This shows that the sets S', S'_a, S'_b, S'_c and S'_d are pairwise disjoint. So we have

$$|S' \cup S'_a \cup S'_b \cup S'_c \cup S'_d| = |S'| + |S'_a| + |S'_b| + |S'_c| + |S'_d|.$$
(7)

By Lemma 19, we have $|\mathcal{P}_a| + |\mathcal{P}_b| + |\mathcal{P}_c| + |\mathcal{P}_d| \leq 2(|S'_a| + |S'_b| + |S'_c| + |S'_d|) + 4 \cdot 7680\sqrt{n}$. Combining this inequality with (6), we get

$$|\mathcal{P}_S| + |\mathcal{P}_a| + |\mathcal{P}_b| + |\mathcal{P}_c| + |\mathcal{P}_d| \leqslant 2 |S'| + 2(|S'_a| + |S'_b| + |S'_c| + |S'_d|) + 4 \cdot 7680\sqrt{n}.$$
 (8)

Combining (7) with (8) we get

$$|\mathcal{P}_{S}| + |\mathcal{P}_{a}| + |\mathcal{P}_{b}| + |\mathcal{P}_{c}| + |\mathcal{P}_{d}| \leq 2 |S' \cup S'_{a} \cup S'_{b} \cup S'_{c} \cup S'_{d}| + 30720\sqrt{n} < 2n + 30720\sqrt{n}, \quad (9)$$

for large enough n.

Each 2-path in $\mathcal{P}_S \cup \mathcal{P}_a \cup \mathcal{P}_b \cup \mathcal{P}_c \cup \mathcal{P}_d$ can be extended to at most three good 3-paths whose first edge is in ∂F . (For example, $auv \in \mathcal{P}_a$ can be extended to bauv, cauv and dauv.) On the other hand, every good 3-path whose first edge is in ∂F must contain a 2-path of $\mathcal{P}_a \cup \mathcal{P}_b \cup \mathcal{P}_c \cup \mathcal{P}_d \cup \mathcal{P}_S$ as a subpath. So the number of good 3-paths whose first edge is in ∂F is at most $3 |\mathcal{P}_a \cup \mathcal{P}_b \cup \mathcal{P}_c \cup \mathcal{P}_d \cup \mathcal{P}_S| = 3(|\mathcal{P}_S| + |\mathcal{P}_a| + |\mathcal{P}_b| + |\mathcal{P}_c| + |\mathcal{P}_d|)$ which is at most $6n + C_3\sqrt{n}$ by (9), for $C_3 = 92160$ and large enough n, proving the desired claim.

2.3 Combining bounds on the number of 3-paths

Recall that $\alpha_1 |G|$, $\alpha_2 |G|$, $(1 - \alpha_1 - \alpha_2) |G|$ are the number of edges of G that are contained in triangles, 2-paths and K_4 's of the edge-decomposition \mathcal{D} of G, respectively. Then the number of triangles, 2-paths and K_4 's in \mathcal{D} is $\alpha_1 |G|/3$, $\alpha_2 |G|/2$ and $(1 - \alpha_1 - \alpha_2) |G|/6$ respectively. Therefore, using Claim 22, Claim 23 and Claim 24, the total number of (ordered) good 3-paths in G is at most

$$\frac{\alpha_1}{3} |G| (8n + C_1 \sqrt{n}) + \frac{\alpha_2}{2} |G| (4n + C_2 \sqrt{n}) + \frac{(1 - \alpha_1 - \alpha_2)}{6} |G| (6n + C_3 \sqrt{n}) \leqslant$$
$$\leqslant |G| n \left(\frac{8\alpha_1}{3} + 2\alpha_2 + (1 - \alpha_1 - \alpha_2)\right) + (C_1 + C_2 + C_3)\sqrt{n} |G| =$$
$$= \frac{n^2 \overline{d}_G}{2} \left(\frac{5\alpha_1 + 3\alpha_2 + 3}{3}\right) + (C_1 + C_2 + C_3) \frac{n^{3/2} \overline{d}_G}{2}.$$

Combining this with the fact that the number of good 3-paths is at least $n\overline{d}_G^3 - C_0 n^{3/2}\overline{d}_G$ (see Claim 21), we get

$$n\overline{d}_{G}^{3} - C_{0}n^{3/2}\overline{d}_{G} \leqslant \frac{n^{2}\overline{d}_{G}}{2} \left(\frac{5\alpha_{1} + 3\alpha_{2} + 3}{3}\right) + (C_{1} + C_{2} + C_{3})\frac{n^{3/2}\overline{d}_{G}}{2}.$$

Rearranging and dividing by $n\overline{d}_G$ on both sides, we get

$$\overline{d}_G^2 \leqslant \left(\frac{5\alpha_1 + 3\alpha_2 + 3}{6}\right)n + \frac{1}{2}\sqrt{n}((C_1 + C_2 + C_3) + 2C_0).$$

Since $(5\alpha_1 + 3\alpha_2 + 3)/6 \ge 1/2$, we may replace 1/2 with $(5\alpha_1 + 3\alpha_2 + 3)/6$ in the above inequality to obtain

$$\overline{d}_{G}^{2} \leqslant \left(\frac{5\alpha_{1} + 3\alpha_{2} + 3}{6}\right) n \left(1 + \frac{(C_{1} + C_{2} + C_{3}) + 2C_{0}}{\sqrt{n}}\right).$$

So letting $C_4 = (C_1 + C_2 + C_3) + 2C_0$ we have,

$$\overline{d}_G \leqslant \sqrt{1 + \frac{C_4}{\sqrt{n}}} \sqrt{\frac{5\alpha_1 + 3\alpha_2 + 3}{6}} \sqrt{n} < \left(1 + \frac{C_4}{2\sqrt{n}}\right) \sqrt{\frac{5\alpha_1 + 3\alpha_2 + 3}{6}} \sqrt{n}, \quad (10)$$

for large enough n. By Claim 11, we have

$$|H| \leqslant \frac{\alpha_1}{3} |G| + \frac{\alpha_2}{2} |G| + \frac{2(1 - \alpha_1 - \alpha_2)}{3} |G| = \frac{4 - 2\alpha_1 - \alpha_2}{6} \frac{n\overline{d}_G}{2}.$$

Combining this with (10) we get

$$|H| \leqslant \left(1 + \frac{C_4}{2\sqrt{n}}\right) \frac{(4 - 2\alpha_1 - \alpha_2)}{12} \sqrt{\frac{5\alpha_1 + 3\alpha_2 + 3}{6}} n^{3/2},$$

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for sufficiently large n. So we have

$$\exp_3(n, C_5) \leqslant (1+o(1)) \frac{(4-2\alpha_1-\alpha_2)}{12} \sqrt{\frac{5\alpha_1+3\alpha_2+3}{6}} n^{3/2}.$$

The right hand side is maximized when $\alpha_1 = 0$ and $\alpha_2 = 2/3$, so we have

$$ex_3(n, C_5) \leq (1+o(1))\frac{4-2/3}{12}\sqrt{\frac{5}{6}}n^{1.5} < (1+o(1))0.2536n^{3/2}.$$

This finishes the proof.

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