

Some Properties of the k -bonacci Words on the Infinite Alphabet

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Abstract

The Fibonacci word W on an infinite alphabet was introduced in [Zhang et al., *Electronic J. Combinatorics* 2017 24(2), 2-52] as a fixed point of the morphism $2i \rightarrow (2i)(2i + 1)$, $(2i + 1) \rightarrow (2i + 2)$, $i \geq 0$. Here, for any integer $k > 2$, we define the infinite k -bonacci word $W^{(k)}$ on the infinite alphabet as $\varphi_k^\omega(0)$, where the morphism φ_k on the alphabet \mathbb{N} is defined for any $i \geq 0$ and any $0 \leq j \leq k - 1$, by

$$\varphi_k(ki + j) = \begin{cases} (ki)(ki + j + 1) & \text{if } j = 0, \dots, k - 2, \\ (ki + j + 1) & \text{otherwise.} \end{cases}$$

We consider the sequence of finite words $(W_n^{(k)})_{n \geq 0}$, where $W_n^{(k)}$ is the prefix of $W^{(k)}$ whose length is the $(n + k)$ -th k -bonacci number. We then provide a recursive formula for the number of palindromes that occur in different positions of $W_n^{(k)}$. Finally, we obtain the structure of all palindromes occurring in $W^{(k)}$ and based on this, we compute the palindrome complexity of $W^{(k)}$, for any $k > 2$.

Mathematics Subject Classifications: 68R15, 11B50

1 Introduction

Finite and infinite Fibonacci words are among the most studied ones in combinatorics of words and have important roles in computer science, based on their optimal properties and various applications, see for example [13, 3, 14, 5]. The sequence of finite Fibonacci words $(F_n)_{n \geq -1}$ is given by $F_{-1} = 1, F_0 = 0$ and the recurrence relation $F_n = F_{n-1}F_{n-2}$ which holds for $n \geq 1$. An equivalent way to give these words for $n \geq 0$, is using $F_n = \psi^n(0)$, where ψ is the binary morphism $0 \rightarrow 01, 1 \rightarrow 0$. The infinite Fibonacci word is then given by $F_\infty = \lim_{n \rightarrow \infty} F_n$ or equivalently by $F_\infty = \psi^\omega(0)$.

The infinite Fibonacci word belongs to the class of infinite aperiodic binary words having the minimal complexity (i.e. the minimal number of factors of each given length); any such word is called a Sturmian word. Sturmian words are well-studied in the literature; they admit some equivalent definitions and have several interesting properties, see [4, 12, 10] for instance.

A natural extension of finite Fibonacci words to k -letter alphabet, $k > 2$, is defining finite k -bonacci words $(F_n^{(k)})_{n \geq 0}$ by

$$F_n^{(k)} = \begin{cases} 0 & \text{if } n = 0, \\ F_{n-1}^{(k)} \dots F_0^{(k)} n & \text{if } 1 \leq n < k, \\ F_{n-1}^{(k)} \dots F_{n-k}^{(k)} & \text{if } n \geq k. \end{cases}$$

Alternatively, these words may be given by $F_n^{(k)} = \psi_k^n(0)$, for $n \geq 0$, where $\psi_k : \{0, \dots, k-1\}^* \rightarrow \{0, \dots, k-1\}^*$ is the morphism

$$\psi_k(i) = \begin{cases} 0(i+1) & \text{if } i = 0, \dots, k-2, \\ 0 & \text{if } i = k-1. \end{cases} \quad (1)$$

The infinite k -bonacci word is then given by $F_\infty^{(k)} = \lim_{n \rightarrow \infty} F_n^{(k)}$ or equivalently by $F_\infty^{(k)} = \psi_k^\omega(0)$.

While the infinite Fibonacci word is the simplest example of Sturmian words, the infinite k -bonacci word is similarly related to the most natural extension of Sturmian words, namely episturmian words. More precisely, the k -bonacci word is the simplest example of non-ultimately periodic episturmian words on the k -letter alphabet; to see the definition and properties of episturmian words see [6, 4, 11, 8, 7].

The infinite Fibonacci word over the infinite alphabet of nonnegative integers, \mathbb{N} , denoted here as $W^{(2)}$, is presented in [15] as the fixed point of the morphism φ_2 starting with 0, where φ_2 is given by $\varphi_2(2i) = (2i)(2i+1)$ and $\varphi_2(2i+1) = 2i+2$ for $i \geq 0$. More precisely, we have $W^{(2)} = \varphi_2^\omega(0)$. The authors of [15] have also studied the finite Fibonacci words over \mathbb{N} , namely the words $W_n^{(2)} = \varphi_2^n(0)$. It is obvious that when the digits of $W_n^{(2)}$ and $W^{(2)}$ are calculated mod 2, these words are reduced to F_n and F_∞ , respectively. Among several properties of words $W_n^{(2)}$ and $W^{(2)}$ studied in [15], the authors characterized palindromic factors of $W_n^{(2)}$ and $W^{(2)}$. Particularly, the authors showed that in contrast to the ordinary infinite Fibonacci word which contains palindromic factors of

arbitrary length, the word $W^{(2)}$ has no palindrome of length greater than 3. Some more properties of these words were consequently studied by Glen et al. in [9]. Among other results, they computed the number of palindromes in $W_n^{(2)}$.

In this paper, we introduce finite and infinite k -bonacci words on the infinite alphabet \mathbb{N} , denoted respectively as $W_n^{(k)}$ and $W^{(k)}$. Studying these words, we characterize the palindromic factors of $W^{(k)}$ for any fixed integer $k \geq 3$. More precisely we show that the length of a palindromic factor of $W^{(k)}$ belongs to the set $L_k = \{2\} \cup \{2i - 1 : 2 \leq i \leq 3 \cdot 2^{k-2}\}$. Conversely, for each element ℓ of L_k we give the structure of palindromes of $W^{(k)}$ with length ℓ . We also enumerate the total number of palindromes of $W_n^{(k)}$.

2 Definitions and notation

In this paper, the alphabet, which can be a finite or a countable infinite set, is denoted as \mathcal{A} . When the alphabet is infinite, we simply take $\mathcal{A} = \mathbb{N}$. Each element of the alphabet \mathcal{A} is called a *letter*. When $\mathcal{A} = \mathbb{N}$, we equivalently use the term *digit* instead of letter. We denote by \mathcal{A}^* the set of finite words over \mathcal{A} and we let $\mathcal{A}^+ = \mathcal{A}^* \setminus \{\epsilon\}$, where ϵ is the empty word. We denote by \mathcal{A}^ω the set of all infinite words over \mathcal{A} and we let $\mathcal{A}^\infty = \mathcal{A}^* \cup \mathcal{A}^\omega$. If $a \in \mathcal{A}$ and $W \in \mathcal{A}^\infty$, then the symbols $|W|$ and $|W|_a$ denote respectively the length of W , and the number of occurrences of letter a in W (It is obvious that when $W \in \mathcal{A}^\omega$, $|W| = \infty$). For any word $W \in \mathcal{A}^\infty$, $\mathcal{Alph}(W)$ is defined to be the set of letters which have at least one occurrence in W , that is $\mathcal{Alph}(W) = \{a \in \mathcal{A} : |W|_a > 0\}$.

A word $V \in \mathcal{A}^*$ is a factor of a word $W \in \mathcal{A}^\infty$, denoted as $V \prec W$, if there exist $U \in \mathcal{A}^*$ and $U' \in \mathcal{A}^\infty$, such that $W = UVU'$. A word $V \in \mathcal{A}^*$ (resp. $V \in \mathcal{A}^\infty$) is said to be a *prefix* (resp. *suffix*) of a word $W \in \mathcal{A}^\infty$, denoted as $V \triangleleft W$ (resp. $V \triangleright W$), if there exists $U \in \mathcal{A}^\infty$ (resp. $U \in \mathcal{A}^*$) such that $W = VU$ (resp. $W = UV$). We denote the prefix (resp. suffix) V of length j of $W \in \mathcal{A}^+$ by $\text{Pref}_j(W)$ (resp. $\text{Suff}_j(W)$). If $W \in \mathcal{A}^*$ and $W = VU$ (resp. $W = UV$), we merely write $V = WU^{-1}$ (resp. $V = U^{-1}W$). For a finite word $W = w_1w_2 \dots w_n$, with $w_i \in \mathcal{A}$ and for $1 \leq j \leq j' \leq n$, we denote $W[j, j'] = w_j \dots w_{j'}$, and for simplicity we denote $W[j, j]$ by $W[j]$. The *reversal* of a finite word $W = w_1w_2 \dots w_n$, with $w_i \in \mathcal{A}$, is $W^R = w_nw_{n-1} \dots w_1$. A word $W \in \mathcal{A}^*$ is called *palindrome* if $W = W^R$. The set of all palindromic factors of the word $W \in \mathcal{A}^\infty$ is denoted by $\mathcal{Pal}(W)$. When the alphabet is finite, for any word $U \in \mathcal{A}^\infty$, the number of palindromic factors of length n of U , called the *palindrome complexity* of U , is denoted by $\text{pal}_U(n)$ (for more information about the palindrome complexity see [1, 2] and the references therein). When the alphabet is infinite (i.e. $\mathcal{A} = \mathbb{N}$), the definition of palindrome complexity can naturally be extended to all words $U \in \mathcal{A}^\infty$ with the same notation.

Let P be a palindrome of odd length $2k + 1$ and $W = \text{Pref}_k(P)$. Then the letter $a \in \mathcal{A}$ satisfying $P = WaW^R$, is called the *center* of the palindrome P . If P is a palindrome of even length, then the center of P is defined to be the empty word. For any occurrence of a palindromic factor $P \in \mathcal{A}^*$ in a word $W \in \mathcal{A}^\infty$ such as $W = UPV$ with $U \in \mathcal{A}^*$, $V \in \mathcal{A}^\infty$, the *center position* of this occurrence of P in W is denoted by $c_p(P, W)$ and defined to be $|U| + \frac{|P|+1}{2}$. We notice that when $|P|$ is even the center-position is a non-integer. A

palindromic factor $P \in \mathcal{A}^*$ of $W \in \mathcal{A}^\infty$ is called a *maximal palindromic factor* of W if there is no longer palindromic factor of W with the same center position. We note that a maximal palindromic factor of a word W , could be a factor of another palindromic factor P of W . For instance, the only maximal palindromic factors of $W = 1213121$ are 1, 121 and 1213121.

For $1 \leq i \leq n$, let $U_i \in \mathcal{A}^*$; then $\prod_{i=n}^1 U_i$ is defined to be $U_n U_{n-1} \dots U_1$. For a finite word W and an integer n , $n \oplus W$ denotes the word obtained by adding n to each letter of W . For example, let $W = 01023$ and $n = 5$, then $n \oplus W = 56578$. For a finite set $S = \{S_1, \dots, S_m\} \subset \mathcal{A}^+$, we define $n \oplus S$ to be the set $\{n \oplus S_1, \dots, n \oplus S_m\}$.

For any integer $k \geq 2$ the sequence of k -bonacci numbers, denoted by $\{f_n^{(k)}\}_{n \geq 0}$, is given as

$$f_n^{(k)} = \begin{cases} 0 & \text{if } n = 0, \dots, k-2, \\ 1 & \text{if } n = k-1, \\ \sum_{i=n-k}^{n-1} f_i^{(k)} & \text{if } n \geq k. \end{cases} \quad (2)$$

The last recurrence relation which holds eventually, states that the n -th term of the sequence is the summation of the k previous ones. This reminds the Fibonacci and Tribonacci recurrence relations in the special cases $k = 2$ and $k = 3$. In fact, $f_n^{(2)}$ and $f_n^{(3)}$ are essentially the well-known Fibonacci and Tribonacci numbers. It is easy to prove that regardless of the first $k-1$ zero terms, the above sequence of k -bonacci numbers can be given as

$$f_n^{(k)} = \begin{cases} 1 & \text{if } n = k-1, \\ 2^{n-k} & \text{if } k \leq n \leq 2k-1, \\ 2f_{n-1}^{(k)} - f_{n-k-1}^{(k)} & \text{if } n \geq 2k. \end{cases} \quad (3)$$

We define the finite (resp. infinite) k -bonacci words $W_n^{(k)}$ (resp. $W^{(k)}$) on the infinite alphabet \mathbb{N} , using the morphism φ_k given below for integers $i \geq 0$ and $0 \leq j < k$,

$$\varphi_k(ki + j) = \begin{cases} (ki)(ki + j + 1) & \text{if } j = 0, \dots, k-2 \\ (ki + j + 1) & \text{if } j = k-1. \end{cases}$$

More precisely, $W_n^{(k)} = \varphi_k^n(0)$ and $W^{(k)} = \varphi_k^\omega(0)$ (Note that $W_0^{(k)} = F_0^{(k)} = 0$). Consequently $F_n^{(k)} = W_n^{(k)} \pmod k$, that is for a fixed value of k , the k -bonacci words over the infinite alphabet are reduced to k -bonacci words over a finite alphabet when the digits are calculated $\pmod k$. It can be shown that for $n \geq 0$,

$$|F_n^{(k)}| = |W_n^{(k)}| = f_{n+k}^{(k)}. \quad (4)$$

We end this section with two examples giving some initial k -bonacci words on finite and infinite alphabets.

Example 1. In this example the 3-bonacci words on finite and infinite alphabet are given when $n \leq 5$. The infinite Tribonacci words on finite and infinite alphabet are also shown.

$$W_0^{(3)} = 0, \quad F_0^{(3)} = 0,$$

$$\begin{array}{ll}
W_1^{(3)} = 01, & F_1^{(3)} = 01, \\
W_2^{(3)} = 0102, & F_2^{(3)} = 0102, \\
W_3^{(3)} = 0102013, & F_3^{(3)} = 0102010, \\
W_4^{(3)} = 0102013010234, & F_4^{(3)} = 0102010010201, \\
W_5^{(3)} = 010201301023401020133435, & F_5^{(3)} = 010201001020101020100102,
\end{array}$$

The first terms of $W^{(3)}$ and $F_\infty^{(3)}$ are as below

$$\begin{array}{l}
W^{(3)} = 01020130102340102013343501020130102343435346 \dots \\
F_\infty^{(3)} = 0102010010201010201001020102010010201020100102010102010 \dots
\end{array}$$

Example 2. In this example we fix $n = 6$ and consider different cases $k = 3, 4, 5, 6$.

$$\begin{array}{l}
W_6^{(3)} = 01020130102340102013343501020130102343435346, \\
F_6^{(3)} = 01020100102010102010010201020100102010102010.
\end{array}$$

$$\begin{array}{l}
W_6^{(4)} = 01020103010201401020103010245010201030102014010201034546, \\
F_6^{(4)} = 01020103010201001020103010201010201030102010010201030102.
\end{array}$$

$$\begin{array}{l}
W_6^{(5)} = 0102010301020104010201030102015010201030102010401020103010256, \\
F_6^{(5)} = 0102010301020104010201030102010010201030102010401020103010201.
\end{array}$$

$$\begin{array}{l}
W_6^{(6)} = 010201030102010401020103010201050102010301020104010201030102016, \\
F_6^{(6)} = 010201030102010401020103010201050102010301020104010201030102010.
\end{array}$$

3 Some properties of $W_n^{(k)}$

In this section, we give some recursive identities which state the word $W_n^{(k)}$ as the concatenation of previous words of the same type. These identities will help us to discover the structure of palindromes in k -bonacci words in the future sections. First we present a simple lemma stating a property of the morphism φ_k which can be easily deduced from the definition.

Lemma 3. For any integer $i \geq 0$, $\varphi_k(k+i) = k \oplus \varphi_k(i)$.

Lemma 4. For $1 \leq n \leq k-1$,

$$W_n^{(k)} = \prod_{i=n-1}^0 W_i^{(k)} n. \tag{5}$$

Proof. We use induction on n . It is easy to check that the statement is true for $n = 1$. Suppose the lemma is true for all i with $1 \leq i \leq n < k - 1$. Then

$$\begin{aligned}
 W_{n+1}^{(k)} &= \varphi_k(W_n^{(k)}) \\
 &= \varphi_k\left(\prod_{i=n-1}^0 W_i^{(k)} n\right) \\
 &= \prod_{i=n-1}^0 \varphi_k(W_i^{(k)}) \varphi_k(n) \\
 &= \prod_{i=n}^1 W_i^{(k)} 0(n+1) \quad \text{since } n < k - 1 \text{ and } \varphi_k(n) = 0(n+1) \\
 &= \prod_{i=n}^0 W_i^{(k)} (n+1). \quad \square
 \end{aligned}$$

Lemma 5. For $1 \leq n \leq k - 1$,

$$W_n^{(k)} = W_{n-1}^{(k)} W_{n-1}^{(k)} (n-1)^{-1} n.$$

Proof. By Lemma 4, we have

$$\begin{aligned}
 W_n^{(k)} &= W_{n-1}^{(k)} W_{n-2}^{(k)} \cdots W_0^{(k)} n \\
 &= W_{n-1}^{(k)} W_{n-1}^{(k)} (n-1)^{-1} n \quad \square
 \end{aligned}$$

In the next lemma we give a recursive formula for $W_n^{(k)}$ when $n \geq k$.

Lemma 6. For $n \geq k$,

$$W_n^{(k)} = \prod_{i=n-1}^{n-k+1} W_i^{(k)} (k \oplus W_{n-k}^{(k)}). \quad (6)$$

Proof. We use induction on n . If $n = k$, the statement is true since

$$\begin{aligned}
 W_k^{(k)} &= \varphi_k(W_{k-1}^{(k)}) \\
 &= \varphi_k\left(\prod_{i=k-2}^0 W_i^{(k)} (k-1)\right) \\
 &= \prod_{i=k-2}^0 \varphi_k(W_i^{(k)}) \varphi_k(k-1) \\
 &= \prod_{i=k-1}^1 W_i^{(k)} k \\
 &= \prod_{i=k-1}^1 W_i^{(k)} (k \oplus W_0^{(k)}).
 \end{aligned}$$

Now, suppose Equation (6) holds for all n with $k \leq n \leq j$; we prove it below for $n = j + 1$.

$$\begin{aligned}
 W_{j+1}^{(k)} &= \varphi_k(W_j^{(k)}) \\
 &= \varphi_k\left(\prod_{i=j-1}^{j-k+1} W_i^{(k)}(k \oplus W_{j-k}^{(k)})\right) \\
 &= \prod_{i=j-1}^{j-k+1} \varphi_k(W_i^{(k)})\varphi_k(k \oplus W_{j-k}^{(k)}) \\
 &= \prod_{i=j}^{j-k+2} W_i^{(k)}(k \oplus \varphi_k(W_{j-k}^{(k)})) \text{ by Lemma 3} \\
 &= \prod_{i=j}^{j-k+2} W_i^{(k)}(k \oplus W_{j-k+1}^{(k)}). \quad \square
 \end{aligned}$$

Corollary 7. For integers $0 \leq i \leq n$, $W_i^{(k)}$ is a prefix of $W_n^{(k)}$.

Proof. By Lemmas 4 and 6, for any integer $n \geq 0$, $W_n^{(k)} \triangleleft W_{n+1}^{(k)}$. Hence, the result follows by induction on n . \square

Corollary 8. Let $k > 2$. For each $0 \leq n \leq k - 1$, $|W_n^{(k)}| = 2^n$ and for each nonnegative n , $|W_{n+1}^{(k)}| \leq 2|W_n^{(k)}|$. Furthermore, the following identity holds.

$$|W_{n+1}^{(k)}| = \begin{cases} 2|W_n^{(k)}| & \text{if } 0 \leq n \leq k - 2, \\ 2|W_n^{(k)}| - 1 & \text{if } n = k - 1, \\ 2|W_n^{(k)}| - |W_{n-k}^{(k)}| & \text{if } n > k - 1. \end{cases}$$

Proof. Considering the definition of k -bonacci numbers in Equation (2) and using Lemmas 4 and 6, we conclude that $|W_n^{(k)}| = f_{n+k}^{(k)}$ holds for all $n \geq 0$. The other statements follow from Equation (3). \square

Lemma 9. For any $n \geq 0$, the digit n is the largest one of $W_n^{(k)}$ and appears once at the end of this word.

Proof. We prove this using induction on n . When $n = 1$, $W_1^{(k)} = \varphi_k(0) = 01$ and it is obvious that the claim is true. Now suppose the result holds for $n \leq m$. We prove it for $n = m + 1$ below.

$$\begin{aligned}
 W_{m+1}^{(k)} &= \varphi_k(W_m^{(k)}) \\
 &= \varphi_k(W_m^{(k)}m^{-1}m) \text{ by the induction hypothesis} \\
 &= \varphi_k(W_m^{(k)}m^{-1})\varphi_k(m).
 \end{aligned}$$

Now, by induction hypothesis all digits of $W_m^{(k)}m^{-1}$ are less than m . Hence, by definition of φ_k , all digits of $\varphi_k(W_m^{(k)}m^{-1})$ are less than $m + 1$. Again using definition of φ_k , the largest digit of $\varphi_k(m)$ is $m + 1$ which occurs once at the end of $\varphi_k(m)$ and the proof is complete. \square

The following notation simplifies some definitions and proofs appearing in the rest of the paper. **Notation.** For any integer i , we let $(i)_* = \max\{i, 0\}$.

Definition 10. Let n be a positive integer. Considering factorizations of $W_n^{(k)}$ given in the Equations (5) and (6), we divide the set of factors of $W_n^{(k)}$ into three following types.

- **Included factors.** The digit n or the factors of $W_n^{(k)}$ which are included in any of the words $W_{n-1}^{(k)}, W_{n-2}^{(k)}, \dots, W_{(n-k+1)_*}^{(k)}$ or in $(k \oplus W_{n-k}^{(k)})$, if $n \geq k$.
- **Bordering factors.** Factors F which are of the form $F = X_j Y_j$ for some $(n - k + 1)_* + 1 \leq j \leq n - 1$, where $X_j \neq \epsilon$ is a suffix of $W_j^{(k)}$ and $Y_j \neq \epsilon$ is a prefix of $\prod_{i=j-1}^{(n-k+1)_*} W_i^{(k)}$. We call any such factor, a *bordering factor of type j* of $W_n^{(k)}$.
- **Straddling factors.** Factors F which are of the form $F = AB$, where $A \neq \epsilon$ is a suffix of $\prod_{i=n-1}^{(n-k+1)_*} W_i^{(k)}$ and $B = n$ if $n \leq k - 1$, and B is a prefix of $k \oplus W_{n-k}^{(k)}$ if $n \geq k$; these are called *(A, B) -straddling factors* (or straddling factors for short, if there is no danger of confusion) of $W_n^{(k)}$.

Definition 11. Let n be a positive integer. Considering Definition 10, we call a palindromic factor P of $W_n^{(k)}$ an *included* (resp. a *bordering*, a *straddling*) *palindrome* if P is an included (resp. a bordering, a straddling) factor of $W_n^{(k)}$.

The following definition helps us to detect factors of length two of $W_n^{(k)}$ and will be used in some lemmas.

Definition 12. Let $\mathcal{B}_1^{(k)} = \{(a.k) \oplus ki : i \geq 0, a \neq 0\}$, $\mathcal{B}_2^{(k)} = \{(0.b) \oplus ki : i \geq 0, 1 \leq b \leq k - 1\}$ and $\mathcal{B}_3^{(k)} = \{a.0 : a \neq 0\}$ and $\mathcal{B}^{(k)} = \mathcal{B}_1^{(k)} \cup \mathcal{B}_2^{(k)} \cup \mathcal{B}_3^{(k)}$.

Lemma 13. Let $B = s.t$ be a factor of length 2 of $W_n^{(k)}$, then $B \in \mathcal{B}^{(k)}$.

Proof. By Definition 12, it is easy to see that $\mathcal{B}^{(k)} \oplus k \subset \mathcal{B}^{(k)}$. We prove the lemma by induction on n . If $n = 1$, then $W_1^{(k)} = 01$ has just one factor of length 2, namely $01 \in \mathcal{B}_2^{(k)}$. Now let $m > 1$ and suppose that the lemma is true for all n , $1 \leq n < m$. To conclude its validity for m , let B be a factor of length 2 of $W_n^{(k)}$. Considering Definition 10, one of the following cases hold.

- **B is an included factor of $W_m^{(k)}$.** Then either $B \prec W_n^{(k)}$ for some $(m - k + 1)_* \leq n \leq m - 1$ or $B \prec (k \oplus W_{m-k}^{(k)})$ (this happens only if $m \geq k$). The former case leads to $B \in \mathcal{B}^{(k)}$ by the induction hypothesis, while the latter case leads to $B \ominus k \in \mathcal{B}^{(k)}$, whence by $\mathcal{B}^{(k)} \oplus k \subset \mathcal{B}^{(k)}$, we have $B \in \mathcal{B}^{(k)}$.
- **B is a bordering factor of $W_m^{(k)}$.** Then by Definition 10, $B = j0$, for some $(n - k + 1)_* + 1 \leq j \leq n - 1$. This means that $B \in \mathcal{B}_3^{(k)}$.
- **B is a straddling factor of $W_m^{(k)}$.** Then either $B = 0.m$, when $m \leq k - 1$, or $B = (m - k + 1)k$. The latter case leads to $B \in \mathcal{B}_1^{(k)}$ by Definition 12 while the former case leads to $B \in \mathcal{B}_2^{(k)}$. □

The next corollaries are direct consequences of Lemma 13.

Corollary 14. *Let $n \geq 0$ and $k > 2$. The finite word $W_n^{(k)}$ contains no factor 00.*

Corollary 15. *Let $s, t \in \mathbb{N}$ and let n, k be two positive integers with $k > 2$. If $st \prec W_n^{(k)}$, then either $k \mid t$ or $s < t$.*

Proof. Suppose that $k \nmid t$, then by Lemma 13, $st \in \mathcal{B}^{(2)}$. Hence, by Definition 12, $s < t$. \square

The following theorem gives us the suffix of length two of the word $W_n^{(k)}$.

Theorem 16. *Let $n \geq 1$ and $j = (n \bmod k)$. Then*

- $(n - j)n \triangleright W_n^{(k)}$ provided $j \neq 0$.
- $(n - k + 1)n \triangleright W_n^{(k)}$ provided $j = 0$.

Proof. Recall that $W_n^{(k)} = \varphi_k^n(0)$. We prove the theorem using induction on n . If $n = j = 1$, we have $01 \triangleright \varphi_k^1(0)$. Suppose that the result holds for all n , $n \leq m$. The validity of the result for $n = m + 1$ is then proved in the following two cases.

- If $m + 1 \equiv 0 \pmod{k}$, then $m \equiv k - 1 \pmod{k}$ and by induction hypothesis $(m - k + 1)(m) \triangleright \varphi_k^m(0)$. Hence,

$$\begin{aligned} \varphi_k((m - k + 1)(m)) &\triangleright \varphi_k^{m+1}(0), \\ \varphi_k(m - k + 1)\varphi_k(m) &\triangleright \varphi_k^{m+1}(0), \\ (m - k + 1)(m - k + 2)(m + 1) &\triangleright \varphi_k^{m+1}(0). \end{aligned}$$

- If $m + 1 \equiv j \pmod{k}$, and $0 < j < k$ then by Lemma 9, m is the last digit of $\varphi_k^m(0)$ or $m \triangleright \varphi_k^m(0)$. Hence,

$$\begin{aligned} \varphi_k(m) &\triangleright \varphi_k^{m+1}(0) \\ (m - j + 1)(m + 1) &\triangleright \varphi_k^{m+1}(0). \end{aligned}$$

So, the proof is complete. \square

4 The number of palindromes in $W_n^{(k)}$

In this section, we are going to count the total number of palindromes in $W_n^{(k)}$. Actually, the process of counting all palindromes in $W_n^{(k)}$ leads us to find all possible palindromic factors of $W_n^{(k)}$. Equations (5) and (6) are essential in the rest of this work. By Equation (5) and Lemma 5, we obtain an explicit formula for the number of palindromes in $W_n^{(k)}$ when $n \leq k - 1$; this is done in Section 4.1.

Let $P^{(k)}(n)$ denote the total number of palindromes in $W_n^{(k)}$ occurring in different positions and $B^{(k)}(n, j)$ and $S^{(k)}(n)$ denote the number of bordering palindromes of type j and straddling palindromes of $W_n^{(k)}$, respectively. Then by Definition 11, the following recurrence relation holds

$$P^{(k)}(n) = \sum_{i=n-k}^{n-1} P^{(k)}(i) + \sum_{i=n-k+2}^{n-1} B^{(k)}(n, i) + S^{(k)}(n). \quad (7)$$

The following theorem gives a recurrence formula for computing $P^{(k)}(n)$. The proof which is formally stated in Section 4.4, is based on considering several cases with respect to values of n and k . This is done in Sections 4.1-4.3.

Theorem 17. *Let $k > 2$ and $n \geq 1$ be given integers. Then the following holds*

(i) *If $0 \leq n \leq k - 1$, then $P^{(k)}(n) = n2^{n-1} + 1$,*

(ii) *If $n \geq k$, then $P^{(k)}(n) = \sum_{i=n-k}^{n-1} P^{(k)}(i) + \alpha^{(k)}(n)$, where*

$$\alpha^{(k)}(n) = \begin{cases} 2^k - k2^{n-k+2} + n2^{n-k+1} & \text{if } k \leq n \leq 2k - 3, \\ 0 & \text{if } n = 2k - 2, \\ 2^{n-2k+2} - 1 & \text{if } 2k - 1 \leq n \leq 3k - 3, \\ 2^k - 2 & \text{if } n = 3k - 2, \\ 0 & \text{if } n > 3k - 2. \end{cases}$$

To prove the theorem we divide the rest of this section into some subsections with respect to the values of n and k .

4.1 Palindromes in $W_n^{(k)}$ when $n \leq k - 1$

Lemma 18. *For $1 \leq n \leq k - 1$, $W_n^{(k)} n^{-1}$ is a palindromic word.*

Proof. We prove this by induction on n . Since for every $k > 2$ we have $W_1^{(k)} = 01$, the first step of the induction is true. Suppose $n = j < k - 1$, the word $W_j^{(k)} j^{-1}$ is a palindrome. Now using Lemma 5, we have

$$\begin{aligned} W_{j+1}^{(k)}(j+1)^{-1} &= W_j^{(k)} W_j^{(k)} j^{-1} \\ &= W_j^{(k)} j^{-1} j W_j^{(k)} j^{-1}, \end{aligned}$$

which is a palindromic word by induction hypothesis. □

Lemma 19. *For every $1 \leq n \leq k - 1$, $P^{(k)}(n) = 2P^{(k)}(n - 1) + 2^{n-1} - 1$ and $P^{(k)}(0) = 1$.*

Proof. Since the digit n just occurs in the last position of $W_n^{(k)}$, every palindromic factor of $W_n^{(k)}$ either equals to n or is a palindromic factor of $W_n^{(k)} n^{-1}$. By Lemma 5, for every $n \leq k - 1$, we have

$$W_n^{(k)} n^{-1} = W_{n-1}^{(k)} W_{n-1}^{(k)} (n-1)^{-1} = W_{n-1}^{(k)} (n-1)^{-1} (n-1) W_{n-1}^{(k)} (n-1)^{-1}. \quad (8)$$

From Equation (8), we conclude that $n - 1$ occurs once in $W_n^{(k)}$. Using Lemma 18 and again using Equation (8), we find that a factor P of $W_n^{(k)}$ is a palindromic word if and only if it is either a palindromic factor of $W_{n-1}^{(k)}$ or $P = a(n-1)a$, where a is a prefix of $W_{n-1}^{(k)}$. Therefore, $P^{(k)}(n) = 2(P^{(k)}(n-1) - 1) + |W_{n-1}^{(k)}| + 1$. Hence, using Corollary 8, we provide $P^{(k)}(n) = 2P^{(k)}(n-1) + 2^{n-1} - 1$. \square

Theorem 20. For every $0 \leq n \leq k - 1$, $P^{(k)}(n) = n2^{n-1} + 1$.

Proof. The proof is easy by induction on n . \square

4.2 Bordering Palindromes of $W_n^{(k)}$

In this section, we consider the bordering palindromes of $W_n^{(k)}$ when $n \geq k$. A bordering palindrome B of $W_n^{(k)}$ is called a *maximal bordering palindrome* if there is no longer bordering palindromic factor of $W_n^{(k)}$ with the same center position.

Lemma 21. Let B_j be one of the bordering palindromes of type j of $W_n^{(k)}$. Then

$$c_p(B_j, W_n^{(k)}) = |W_{n-1}^{(k)}| + \cdots + |W_j^{(k)}|$$

and j is the center of B_j .

Proof. By Definition 11, we know that B_j is a factor of $W = W_j^{(k)}W_{j-1}^{(k)} \cdots W_{n-k+1}^{(k)}$, and it contains the last digit of $W_j^{(k)}$, which is j . By Lemma 9, $|W|_j = 1$. Hence $|B_j|$ is an odd integer and j is the center of B_j . Moreover $c_p(B_j, W_n^{(k)}) = |W_{n-1}^{(k)}| + \cdots + |W_j^{(k)}|$. \square

Lemma 22. Let B_j be a bordering palindrome of type j of $W_n^{(k)}$. Then $n \leq 2k - 3$ and $n - k + 2 \leq j \leq k - 1$.

Proof. Let $c = c_p(B_j, W_n^{(k)})$ and $j' = (j \bmod k)$. By Theorem 16 and Lemma 21, we have $B_j[c] = j$ and

$$B_j[c-1] = \begin{cases} j - j' & \text{if } (j' \neq 0) \\ j - k + 1 & \text{otherwise} . \end{cases}$$

Using Definition 10 and Equation (6), we have $B_j[c+1] = 0$. If $j' \neq 0$, then by definition of c , we obtain $B_j[c+1] = B_j[c-1] = j - j'$. Since $B_j[c+1] = 0$, $j = j'$, which shows that $0 \leq j \leq k - 1$. Moreover, by Equation (6), $n - k + 2 \leq j \leq n - 1$. Hence $n - k + 2 \leq j \leq k - 1$, which shows that $n \leq 2k - 3$. If $j' = 0$, then $B_j[c-1] = j - k + 1$. Using $B_j[c+1] = B_j[c-1] = j - k + 1$, $B_j[c+1] = 0$ and $j \geq n - k + 2$, we have $j = k - 1$ and $n \leq 2k - 3$, as desired. \square

Lemma 23. Let $k \leq n \leq 2k - 3$ and $n - k + 2 \leq j \leq k - 1$. Then the maximal bordering palindrome of type j of $W_n^{(k)}$ is $B_j = (W_{j-1}^{(k)}W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R j (W_{j-1}^{(k)}W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})$.

Proof. By Definition 11, B_j is a palindromic factor of the following word

$$W_j^{(k)} j^{-1} j W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)},$$

and by Lemma 21, j is the center of B_j . Using Lemma 22, $j \leq k-1$ and $n-k+1 \leq j-1$. On the other hand by Lemma 4, we have $W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)} \triangleleft W_j^{(k)} j^{-1}$. Hence, using Lemma 18, $(W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R \triangleright W_j^{(k)} j^{-1}$. Therefore,

$$B_j = (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R j (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)}). \quad \square$$

Lemma 24. *Let n and j be two integers with $k \leq n \leq 2k-3$ and $n-k+2 \leq j \leq k-1$ and B_j be a maximal bordering palindrome of type j of $W_n^{(k)}$. Then $|B_j| = 2(|W_j^{(k)}| - |W_{n-k+1}^{(k)}|) + 1$.*

Proof. By Lemma 23, we have

$$B_j = (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R j (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)}).$$

Using Equation (5),

$$W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)} = W_j^{(k)} (W_{n-k}^{(k)} \cdots W_0^{(k)} j)^{-1}.$$

Hence,

$$\begin{aligned} |B_j| &= 2(|W_j^{(k)}| - |W_{n-k}^{(k)} \cdots W_0^{(k)} j|) + 1 \\ &= 2(|W_j^{(k)}| - |W_{n-k+1}^{(k)}|) + 1. \end{aligned} \quad \square$$

Lemma 25. *Let $n \geq k$, then*

$$B^{(k)}(n, j) = \begin{cases} 2^j - 2^{n-k+1} & \text{if } (k \leq n \leq 2k-3) \text{ and } (n-k+2 \leq j \leq k-1), \\ 0 & \text{otherwise.} \end{cases}$$

Proof. If $n \leq 2k-3$ and $n-k+2 \leq j \leq k-1$ and B_j be the maximal bordering palindrome of type j of $W_n^{(k)}$. then by Lemmas 23 and 24, we have

$$\begin{aligned} B^{(k)}(n, j) &= \frac{|B_j| - 1}{2} \\ &= |W_j^{(k)}| - |W_{n-k+1}^{(k)}| \\ &= 2^j - 2^{n-k+1}. \end{aligned}$$

The last equality holds using Corollary 8. In other cases, by Lemma 22, $B^{(k)}(n, j) = 0$, as desired. \square

4.3 Straddling Palindromes of $W_n^{(k)}$

In this subsection, we are going to count the number of straddling palindromes of $W_n^{(k)}$ when $n \geq k$. A straddling palindrome of $W_n^{(k)}$ is called a *maximal straddling palindrome* of $W_n^{(k)}$ if there is no longer straddling palindromic factor of $W_n^{(k)}$ with the same center position. Similarly, an (A, B) -*maximal straddling palindrome*, is an (A, B) -straddling palindrome which is a maximal straddling palindrome.

Lemma 26. *If S is a straddling palindrome of $W_n^{(k)}$, then $2k - 1 \leq n \leq 3k - 2$.*

Proof. First we note that $(k \oplus W_{n-k}^{(k)})$ starts with k and $W_{n-k+1}^{(k)}$ ends with $n - k + 1$. Hence, the sequence $(n - k + 1).k$ occurs in any straddling palindrome P of $W_n^{(k)}$. Since P is a palindrome, the sequence $B = k.(n - k + 1)$ also occurs in P . Hence, by Lemma 13, $B \in \mathcal{B}^{(k)}$. Since $n \geq k$, using Definition 12, we provide $B \in \mathcal{B}_1^{(k)} \cup \mathcal{B}_2^{(k)}$. If $B \in \mathcal{B}_1^{(k)}$, then $B = k.(n - k + 1) = (a.k) \oplus ki$, for some integers $a \neq 0$ and $i \geq 0$. Since $a \neq 0$, we conclude that $a = k$ and $i = 0$. So, in this case $n - k + 1 = k$ and $n = 2k - 1$.

If $B \in \mathcal{B}_2^{(k)}$, then $B = k.(n - k + 1) = (0.b) \oplus ki$ for some $i \geq 0$ and $1 \leq b \leq k - 1$. Hence, $i = 1$ and $k + 1 \leq n - k + 1 \leq 2k - 1$. Therefore, $2k \leq n \leq 3k - 2$, as desired. \square

Lemma 27. *The only straddling palindrome of $W_{2k-1}^{(k)}$ is $k.k$.*

Proof. By Lemma 6, $W_{2k-1}^{(k)} = W_{2k-2}^{(k)} \cdots W_k^{(k)}(k \oplus W_{k-1}^{(k)})$. It is clear that $k \triangleleft (k \oplus W_{k-1}^{(k)})$ and using Theorem 16, $(1.k) \triangleright W_k^{(k)}$. Hence, $k.k$ is a straddling palindrome of $W_{2k-1}^{(k)}$. Now, we are going to show that there is no other straddling palindrome in $W_{2k-1}^{(k)}$. Since there is no digit 1 in $(k \oplus W_{k-1}^{(k)})$, every straddling palindrome of $W_{2k-1}^{(k)}$ has the last digit of $W_k^{(k)}$ (i.e. k) as a prefix. By Corollary 14, 0.0 is not a factor of $W_{k-1}^{(k)}$ and hence $k.k$ is not a factor of $(k \oplus W_{k-1}^{(k)})$. Therefore, $W_{2k-1}^{(k)}$ could not have a straddling palindrome of length greater than 2 and $k.k$ is its only straddling palindrome. \square

Lemma 28. *Let $2k - 1 < n < 3k - 2$ and let W be an (A, B) -maximal straddling palindrome of $W_n^{(k)}$. Then $A \triangleright (k \oplus W_{n-2k+1}^{(k)})$ and hence $|A| \leq 2^{n-2k+1}$.*

Proof. Let $i = n - 2k$. Then, by Equation (6),

$$\begin{aligned} W_{2k+i}^{(k)} &= W_{2k+i-1}^{(k)} \cdots W_{k+i+1}^{(k)}(k \oplus W_{k+i}^{(k)}), \\ W_{k+i+1}^{(k)} &= W_{k+i}^{(k)} \cdots W_{i+2}^{(k)}(k \oplus W_{i+1}^{(k)}). \end{aligned}$$

Hence,

$$\begin{aligned} W_{k+i+1}^{(k)}(k \oplus W_{k+i}^{(k)}) &\triangleright W_{2k+i}^{(k)}, \\ W_{i+2}^{(k)}(k \oplus W_{i+1}^{(k)}) &\triangleright W_{k+i+1}^{(k)}. \end{aligned}$$

Therefore,

$$W_{i+2}^{(k)}(k \oplus W_{i+1}^{(k)})(k \oplus W_{k+i}^{(k)}) \triangleright W_{2k+i}^{(k)}. \tag{9}$$

Let $W_{2k+i}^{(k)} = VW_{i+2}^{(k)}(k \oplus W_{i+1}^{(k)})(k \oplus W_{k+i}^{(k)})$ and $\ell = |VW_{i+2}^{(k)}|$. By definition of the (A, B) -straddling palindrome of $W_{2k+i}^{(k)}$ and using Equation (9), we conclude that $A \triangleright VW_{i+2}^{(k)}(k \oplus W_{i+1}^{(k)})$. Let $A = W_{2k+i}^{(k)}[j, j']$, where by definition of A , $j' = |W_{2k+i}^{(k)}| - |W_{k+i}^{(k)}|$. We claim that $j > \ell$. For the contrary, let $j \leq \ell$. The letter following $i + 2$ in Equation (9) is k . Hence, $(i + 2)k \prec W$ and so $k(i + 2) \prec W$, but this is impossible by Corollary 15, since $1 < i + 2 < k$. Therefore, $A \triangleright (k \oplus W_{i+1}^{(k)})$. Hence, $|A| \leq |W_{i+1}^{(k)}|$, so by Corollary 8, $|A| \leq 2^{i+1}$. \square

The next lemma is useful to give an upper bound for the size of the word B in any (A, B) -maximal straddling palindrome of $W_n^{(k)}$, $2k \leq n \leq 3k - 2$.

Lemma 29. *Let $i \geq 0$ and P be palindromic prefix of $W = (i + 1)W_{k+i}^{(k)}$. Then the largest digit of P is $i + 1$.*

Proof. Let ℓ be the largest digit of P and for contrary suppose that $\ell > i + 1$. Using Lemma 9, we have $\ell < i + k$. Hence, $\ell + 1 \leq i + k$ which yields $W_{\ell+1}^{(k)} \triangleleft W_{i+k}^{(k)}$. We note that $\ell + 1$, which is the last digit of $W_{\ell+1}^{(k)}$, does not appear in P . Hence,

$$|P| \leq |W_{\ell+1}^{(k)}|. \quad (10)$$

On the other hand $W_{\ell}^{(k)} \triangleleft W_{\ell+1}^{(k)}$ and since $i + 1 < \ell$ and using Lemma 9, we conclude that the first place that ℓ occurs in W is the last digit of $W_{\ell}^{(k)}$. In other words if we let $m = |W_{\ell}^{(k)}| + 1$, then $W[m] = \ell$ and for any integer $j < m$ we have $W[j] < \ell$. By our assumption $\ell \in \text{Alph}(P)$, therefore $c_p(P, W) \geq m = |W_{\ell}^{(k)}| + 1$. Thus $|P| \geq 2|W_{\ell}^{(k)}| + 1$ and using Equation (10), we have $|P| \leq |W_{\ell+1}^{(k)}|$. Hence, $2|W_{\ell}^{(k)}| + 1 \leq |W_{\ell+1}^{(k)}|$ but this contradicts with Corollary 8. Hence our assumption is not true and $\ell \leq i + 1$, as desired. \square

Lemma 30. *Let $2k - 1 < n < 3k - 2$ and let W be an (A, B) -maximal straddling palindrome of $W_n^{(k)}$. Then $B \triangleleft k \oplus (W_{n-2k+2}^{(k)}(n - 2k + 2)^{-1})$ and hence $|B| \leq 2^{n-2k+2} - 1$.*

Proof. Let $i = n - 2k$. Then, using Equation (6) we have

$$W_{2k+i}^{(k)} = W_{2k+i-1}^{(k)} \cdots W_{k+i+1}^{(k)}(k \oplus W_{k+i}^{(k)}).$$

So we have

$$B \triangleleft (k \oplus W_{k+i}^{(k)}).$$

By Lemma 28,

$$A \triangleright (k \oplus W_{i+1}^{(k)}) \quad (11)$$

Hence, $k + i + 2 \notin \text{Alph}(A)$. We claim that $k + i + 2 \notin \text{Alph}(B)$. For contrary suppose that $k + i + 2 \in \text{Alph}(B)$. We note that $(k \oplus W_{i+2}^{(k)}) \triangleleft (k \oplus W_{k+i}^{(k)})$. Now, let $m = |W_{i+2}^{(k)}|$, then by Lemma 9, $B[m] = k + i + 2$ and for all $m' < m$, $B[m'] < k + i + 2$. On the other hand by Equation (11), $k + i + 2 \notin \text{Alph}(A)$, hence $c_p(W, W) \geq |A| + m$. Therefore, $(k + i + 1)B$ contains a palindromic prefix P with $k + i + 2 \in \text{Alph}(P)$, which is impossible by Lemma 29. Hence, $k + i + 2 \notin \text{Alph}(B)$, which implies that $B \triangleleft k \oplus (W_{i+2}^{(k)}(i + 2)^{-1})$ and by Corollary 8, $|B| \leq 2^{i+2} - 1 = 2^{n-2k+2} - 1$. \square

Now, we are ready to prove the following lemma.

Lemma 31. *Let $2k - 1 < n < 3k - 2$. Then $W_n^{(k)}$ has exactly two maximal straddling palindromes, one of which is the $(k \oplus W_{n-2k+1}^{(k)}, k \oplus (W_{n-2k+1}^{(k)}(n - 2k + 1)^{-1}))$ -straddling palindrome and the other is the $(k \oplus W_{n-2k+1}^{(k)}, k \oplus (W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)}(n - 2k + 1)^{-1}))$ -straddling palindrome.*

Proof. Let $S = AB$ be an (A, B) -maximal straddling palindrome of $W_n^{(k)}$. Let $i = n - 2k$ and $W = k \oplus (W_{i+1}^{(k)} W_{i+2}^{(k)}(i + 2)^{-1})$. Then by Lemmas 28 and 30, we conclude that $S \prec W$. Moreover, by Equation (5) we have

$$\begin{aligned} W &= k \oplus (W_{i+1}^{(k)} W_{i+2}^{(k)}(i + 2)^{-1}) \\ &= k \oplus (W_{i+1}^{(k)} W_{i+1}^{(k)} W_i^{(k)} \cdots W_0^{(k)}) \\ &= k \oplus (W_{i+1}^{(k)} W_{i+1}^{(k)} W_{i+1}^{(k)}(i + 1)^{-1}) \\ &= k \oplus (W_{i+1}^{(k)}(i + 1)^{-1}(i + 1)W_{i+1}^{(k)}(i + 1)^{-1}(i + 1)W_{i+1}^{(k)}(i + 1)^{-1}). \end{aligned} \tag{12}$$

By the last equation and using Lemma 9, we conclude that $|W|_{k+i+1} = 2$. Let c_1 and c_2 be two integers with $c_1 < c_2 \leq |W|$ such that $W[c_1] = W[c_2] = k + i + 1$. Using Equation (6), we have

$$\begin{aligned} W_{2k+i}^{(k)} &= W_{2k+i-1}^{(k)} \cdots W_{k+i+2}^{(k)} W_{k+i+1}^{(k)}(k \oplus W_{k+i}^{(k)}) \\ &= W_{2k+i-1}^{(k)} \cdots W_{k+i+2}^{(k)}(W_{k+i}^{(k)} \cdots W_{i+2}^{(k)}(k \oplus W_{i+1}^{(k)}))(k \oplus W_{k+i}^{(k)}) \end{aligned}$$

W starts by $(k \oplus W_{k+i}^{(k)})$ in the previous expression. Thus, the first occurrence of $k + i + 1$ in W is the last letter of $(k \oplus W_{k+i}^{(k)})$. Since S is a straddling palindrome of $W_{2k+i}^{(k)}$, the first occurrence of the digit has to be in S . Therefore, either $c_p(S, W) = c_1$ or $c_p(S, W) = \frac{c_1+c_2}{2}$. In other words, the center position of S is either the corresponding position of the first occurrence of $k + i + 1$, or is the position of the middle digit between the two occurrences of $k + i + 1$ in W . So, we have the following two cases

- S has only one digit $k + i + 1$: in this case we show that $A = k \oplus W_{i+1}^{(k)}$ and $B = k \oplus (W_{i+1}^{(k)}(i + 1)^{-1})$. By Lemma 28, it suffices to show that $k \oplus (W_{i+1}^{(k)} W_{i+1}^{(k)}(i + 1)^{-1})$ is a palindrome and this is true by Lemma 18.
- S has exactly two digits $k + i + 1$: in this case we show that $A = k \oplus W_{i+1}^{(k)}$ and $B = k \oplus (W_{i+1}^{(k)} W_{i+1}^{(k)}(i + 1)^{-1})$. By Lemma 18, $k \oplus (W_{i+1}^{(k)} W_{i+1}^{(k)} W_{i+1}^{(k)}(i + 1)^{-1})$ is a palindrome and using Lemmas 28 and 30 and Equation (12), $S = AB$ is a maximal straddling palindrome, as desired. \square

Lemma 32. *Let $n = 3k - 2$ and let W be an (A, B) -maximal straddling palindrome of $W_n^{(k)}$. Then $A \triangleright (k \oplus W_{k-1}^{(k)})$ and hence $|A| \leq 2^{k-1}$.*

Proof. By Lemma 6, we have

$$\begin{aligned} W_{3k-2}^{(k)} &= W_{3k-3}^{(k)} \cdots W_{2k-1}^{(k)}(k \oplus W_{2k-2}^{(k)}), \\ W_{2k-1}^{(k)} &= W_{2k-2}^{(k)} \cdots W_k^{(k)}(k \oplus W_{k-1}^{(k)}). \end{aligned}$$

Hence, $W_k^{(k)}(k \oplus W_{k-1}^{(k)})(k \oplus W_{2k-2}^{(k)}) \triangleright W_{3k-2}^{(k)}$. Since $1k \triangleright W_k^{(k)}$ and by Corollary 15, $k1 \not\prec W_{3k-2}^{(k)}$, we conclude that $|A| \leq |k(k \oplus W_{k-1}^{(k)})|$. Hence, using Corollary 8, we have $|A| \leq 2^{k-1} + 1$. Moreover, $kk \triangleleft k(k \oplus W_{k-1}^{(k)})$, but kk could not appear in $k \oplus (W_{k-1}^{(k)} W_{2k-2}^{(k)})$, because otherwise $00 \prec (W_{k-1}^{(k)} W_{2k-2}^{(k)})$ which is impossible by Corollary 14. Therefore, $A \triangleright (k \oplus W_{k-1}^{(k)})$ and using Corollary 8, $|A| \leq 2^{k-1}$. \square

Lemma 33. *Let $n = 3k - 2$ and let W be an (A, B) -maximal straddling palindrome of $W_n^{(k)}$. Then $B \triangleleft k \oplus (W_k^{(k)} k^{-1})$, hence, $|B| \leq 2^k - 1$.*

Proof. Using Equation (6) we have

$$\begin{aligned} W_{3k-2}^{(k)} &= W_{3k-3}^{(k)} \cdots W_{2k-1}^{(k)}(k \oplus W_{2k-2}^{(k)}), \\ W_{2k-1}^{(k)} &= W_{2k-2}^{(k)} \cdots W_k^{(k)}(k \oplus W_{k-1}^{(k)}). \end{aligned}$$

So we have

$$B \triangleleft (k \oplus W_{2k-2}^{(k)}).$$

By Lemma 32,

$$A \triangleright (k \oplus W_{k-1}^{(k)})$$

Hence, $2k \notin \text{Alph}(A)$. We claim that $2k \notin \text{Alph}(B)$ as well. For contrary suppose that $2k \in \text{Alph}(B)$. We note that $(k \oplus W_k^{(k)}) \triangleleft (k \oplus W_{2k-2}^{(k)})$. Now, let $m = |W_k^{(k)}|$, then by Lemma 9, $B[m] = 2k$ and for all $m' < m$, $B[m'] < 2k$. On the other hand, $2k \notin \text{Alph}(A)$, hence $c_p(W, W) \geq |A| + m$. Therefore, $(2k - 1)B$ contains a palindrome prefix P with $2k \in \text{Alph}(P)$, which is impossible by Lemma 29. Hence, $2k \notin \text{Alph}(B)$, which implies that $B \triangleleft k \oplus (W_k^{(k)}(k)^{-1})$ and by Corollary 8 $|B| \leq 2^k - 1$. \square

Lemma 34. *Let $n = 3k - 2$. Then $W_n^{(k)}$ has exactly two maximal straddling palindromes which are respectively the $(k \oplus W_{k-1}^{(k)}, k \oplus (W_{k-1}^{(k)}(k - 1)^{-1}))$ -straddling palindrome and the $(k \oplus (0^{-1}W_{k-1}^{(k)}), k \oplus (W_{k-1}^{(k)}W_{k-1}^{(k)}(0(k - 1))^{-1}))$ -straddling palindrome in $W_n^{(k)}$.*

Proof. Let W be an (A, B) -maximal straddling palindrome of $W_{3k-2}^{(k)}$. Using Equation (6) we have

$$\begin{aligned} k \oplus (W_k^{(k)} k^{-1}) &= k \oplus (W_{k-1}^{(k)} W_{k-2}^{(k)} \cdots W_1^{(k)} k \cdot k^{-1}) \\ &= k \oplus (W_{k-1}^{(k)} W_{k-2}^{(k)} \cdots W_1^{(k)}) \\ &= k \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k - 1))^{-1}). \end{aligned}$$

Therefore, using Lemmas 33 and 34 every (A, B) -straddling palindrome S of $W_{3k-2}^{(k)}$, we have $S = AB \prec W$, where

$$W = k \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k-1))^{-1}) \quad (13)$$

$$= k \oplus (W_{k-1}^{(k)} (k-1)^{-1} (k-1) W_{k-1}^{(k)} (k-1)^{-1} (k-1) W_{k-1}^{(k)} (0(k-1))^{-1}). \quad (14)$$

By Lemma 9 and Equation (14) the digit $2k-1$ occurs twice in W . Let c_1 and c_2 be two integers with $c_1 < c_2 \leq |W|$ such that $W[c_1] = W[c_2] = 2k-1$. Since S is a straddling palindrome of $W_{3k-2}^{(k)}$, S should contain the first occurrence of $2k-1$ in W . Therefore, either we have $c_p(S, W) = c_1$ or $c_p(S, W) = \frac{c_1+c_2}{2}$. In other words the center position of S is either the position of the first occurrence of $2k-1$ in W or exactly the position of the digit in the middle of the two occurrences of $2k-1$ in W . So, we have the following two cases

- $|S|_{2k-1} = 1$: In this case we show that $A = k \oplus W_{k-1}^{(k)}$ and $B = k \oplus (W_{k-1}^{(k)} (k-1)^{-1})$. By Lemma 32, it suffices to show that $k \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} (k-1)^{-1})$ is a palindrome and this is true by Lemma 18.
- $|S|_{2k-1} = 2$: In this case we show that $A = k \oplus (0^{-1} W_{k-1}^{(k)})$ and $B = k \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k-1))^{-1})$. By Lemma 18, $k \oplus (0^{-1} W_{k-1}^{(k)} W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k-1))^{-1})$ is a palindrome and using Lemmas 32 and 33, $S = AB$ is a maximal straddling palindrome as desired. \square

Theorem 35. Let $S^{(k)}(n)$ be the number of straddling palindromes of $W_n^{(k)}$. Then

$$S^{(k)}(n) = \begin{cases} 2^{n-2k+2} - 1 & \text{if } 2k-1 \leq n < 3k-2, \\ 2^k - 2 & \text{if } n = 3k-2, \\ 0 & \text{otherwise.} \end{cases}$$

Proof. If $n = 2k-1$, then by Lemma 27, the only straddling palindrome of $W_n^{(k)}$ is kk . Hence, $S^{(k)}(2k-1) = 1$.

If $2k \leq n \leq 3k-3$, then by Lemma 31, the only maximal straddling palindromes of $W_n^{(k)}$ are $S_1 = (k \oplus W_{i+1}^{(k)})(k \oplus (W_{i+1}^{(k)} (i+1))^{-1})$ and $S_2 = (k \oplus W_{i+1}^{(k)})(k \oplus (W_{i+1}^{(k)} W_{i+1}^{(k)} (i+1))^{-1})$, where $i = n - 2k$. According to the proof of Lemma 31, the center of S_1 is $k + i + 1$ and $c_p(S_1, S_1) = |W_{i+1}^{(k)}|$. Therefore, the number of straddling palindromes with the same center position as S_1 , is $|W_{i+1}^{(k)}| - 1$, which equals to $2^{i+1} - 1 = 2^{n-2k+1} - 1$, using Corollary 8. Now, again by the proof of Lemma 31, $c_p(S_2, S_2) > |W_{i+1}^{(k)}|$. Hence, by Corollary 8, the number of straddling palindromes with the same center position as S_2 , is $|W_{i+1}^{(k)}| = 2^{i+1} = 2^{n-2k+1}$. Therefore, for $2k \leq n \leq 3k-3$, $S^{(k)}(n) = 2^{n-2k+1} + 2^{n-2k+1} - 1 = 2^{n-2k+2} - 1$.

If $n = 3k-2$, then by Lemma 34, the only maximal straddling palindromes of $W_n^{(k)}$ are $S_1 = k \oplus W_{k-1}^{(k)} k \oplus (W_{k-1}^{(k)} (k-1)^{-1})$ and $S_2 = k \oplus (0^{-1} W_{k-1}^{(k)}) k \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k-1))^{-1})$. According to the proof of Lemma 34, $c_p(S_1, S_1) = |W_{k-1}^{(k)}|$. Therefore, the number of straddling palindromes with the same center position as S_1 is $|W_{k-1}^{(k)}| - 1$, which equals to

$2^{k-1} - 1$ by Corollary 8. Now, again by the proof of Lemma 34, $c_p(S_2, S_2) > |W_{k-1}^{(k)}|$. Hence, using Corollary 8, the number of straddling palindrome with the same center position as S_2 , is $|0^{-1}W_{k-1}^{(k)}| = 2^{k-1} - 1$. Therefore, $S^{(k)}(3k - 2) = 2^k - 2$.

Finally, $S^{(k)}(n) = 0$, when $n < 2k - 1$ or $n > 3k - 2$, using by Lemma 26. \square

4.4 Proof of Theorem 17

Theorem 12. *Let $k > 2$ and $n \geq 0$ be given integers. Then the following holds*

(i) *If $0 \leq n \leq k - 1$, then $P^{(k)}(n) = n2^{n-1} + 1$,*

(ii) *If $n \geq k$, then $P^{(k)}(n) = \sum_{i=n-k}^{n-1} P^{(k)}(i) + \alpha^{(k)}(n)$, where*

$$\alpha^{(k)}(n) = \begin{cases} 2^k - k2^{n-k+2} + n2^{n-k+1} & \text{if } k \leq n \leq 2k - 3, \\ 0 & \text{if } n = 2k - 2, \\ 2^{n-2k+2} - 1 & \text{if } 2k - 1 \leq n \leq 3k - 3, \\ 2^k - 2 & \text{if } n = 3k - 2, \\ 0 & \text{if } n > 3k - 2. \end{cases}$$

Proof. (i) This part is true according to Theorem 20.

(ii) To prove this part by Equation (7) we have

$$P^{(k)}(n) = \sum_{i=n-k}^{n-1} P^{(k)}(i) + \sum_{j=n-k+2}^{n-1} B^{(k)}(n, j) + S^{(k)}(n).$$

Therefore,

$$P^{(k)}(n) = \sum_{i=n-k}^{n-1} P^{(k)}(i) + \alpha^{(k)}(n),$$

where

$$\alpha^{(k)}(n) = \sum_{j=n-k+2}^{n-1} B^{(k)}(n, j) + S^{(k)}(n). \quad (15)$$

By Theorem 35, when $k \leq n \leq 2k - 3$, $S^{(k)}(n) = 0$. Therefore, $\alpha^{(k)}(n) = \sum_{j=n-k+2}^{n-1} B^{(k)}(n, j)$, when $k \leq n \leq 2k - 3$. Now, by Lemmas 22 and 25 we have

$$\begin{aligned} \alpha^{(k)}(n) &= \sum_{j=n-k+2}^{n-1} B^{(k)}(n, j) \\ &= \sum_{j=n-k+2}^{k-1} B^{(k)}(n, j) \text{ by Lemma 22} \\ &= \sum_{j=n-k+2}^{k-1} (2^j - 2^{n-k+1}) \\ &= 2^k - k2^{n-k+2} + n2^{n-k+1}. \end{aligned}$$

If $2k - 1 \leq n \leq 3k - 2$, then by Lemma 25 and Equation (15), $\alpha^{(k)}(n) = S^{(k)}(n)$. Hence using Theorem 35, we have

$$\alpha^{(k)}(n) = \begin{cases} 2^{n-2k+2} - 1 & \text{if } 2k - 1 \leq n \leq 3k - 3, \\ 2^k - 2 & \text{if } n = 3k - 2. \end{cases}$$

Finally, using Theorem 35 and Lemma 25, $\alpha^{(k)}(n) = 0$, if either $n = 2k - 2$ or $n > 3k - 2$. \square

4.5 Examples

In the following example in the case $k = 4$ for some different values of n we give all of the maximal straddling palindromes and the maximal bordering palindromes of $W_n^{(k)}$, if there exists any.

Example 36. Let $k = 4$. Then by Lemma 23, the word $W_4^{(4)}$ has two maximal bordering palindromes $B_2 = (W_1^{(4)})^R 2 (W_1^{(4)})$ and $B_3 = (W_2^{(4)} W_1^{(4)})^R 3 (W_2^{(4)} W_1^{(4)})$, which are shown below. We also notice that by Theorem 35, $W_4^{(4)}$ has no straddling palindrome.

$$\begin{aligned} W_4^{(4)} &= 01020103.0102.01.4 \\ W_4^{(4)} &= 01020103.0102.01.4. \end{aligned}$$

In the case $n = 5 = 2k - 3$, by Lemma 23, the word $W_5^{(4)}$ has a maximal bordering palindrome $B_3 = (W_2^{(4)})^R 3 (W_2^{(4)})$, which is shown below. We also notice that by Theorem 35, $W_5^{(4)}$ has no straddling palindrome.

$$W_5^{(4)} = 010201030102014.01020103.0102.45 \tag{16}$$

In the case $n = 6$, by Theorem 35 and Lemma 25, $W_6^{(4)}$ contains neither a straddling palindrome nor a bordering palindrome.

$$W_6^{(4)} = 01020103010201401020103010245.010201030102014.01020103.4546$$

Example 37. Let $k = 4$. Again by the same reasoning as the previous example the word $W_7^{(4)}$ has one straddling palindrome and no bordering palindrome.

$$\begin{aligned} W_7^{(4)} &= 01020103010201401020103010245010201030102014010201034546.010201030102014 \\ &01020103010245.010201030102014.45464547 \end{aligned}$$

By Lemma 25, $W_9^{(4)}$ has no bordering palindrome and by Lemma 31 it has two maximal straddling palindromes which are given by $(4 \oplus W_2^{(4)}, 4 \oplus (W_2^{(4)} 2^{-1}))$ and $(4 \oplus W_2^{(4)}, 4 \oplus (W_2^{(4)} W_2^{(4)} 2^{-1}))$ as shown below.

$$\begin{aligned} W_9^{(4)} &= 01020103010201401020103010245010201030102014010201034546010201030102014 \\ &01020103010245010201030102014454645470102010301020140102010301024501020 \\ &103010201401020103454601020103010201401020103010245454645474546458.0102 \\ &01030102014010201030102450102010301020140102010345460102010301020140102 \\ &010301024501020103010201445464547.0102010301020140102010301024501020103 \\ &0102014010201034546.45464547454645845464547454689 \end{aligned}$$

$$W_9^{(4)} = 01020103010201401020103010245010201030102014010201034546010201030102014$$

01020103010245010201030102014454645470102010301020140102010301024501020
103010201401020103454601020103010201401020103010245454645474546458.0102
01030102014010201030102450102010301020140102010345460102010301020140102
010301024501020103010201445464547.0102010301020140102010301024501020103
0102014010201034546.45464547454645845464547454689.

5 Palindrome Structure

In this section, based on finding the structure of all palindromic factors of $W^{(k)}$, we compute its palindrome complexity, $\text{pal}_{W^{(k)}}(n)$. We recall that for a fixed word $U \in \mathcal{A}^\infty$, $\text{pal}_U(n)$ is the number of palindromic factors of length n of U . Hence, $\text{pal}_U(n)$ is a function from \mathbb{N} to $\mathbb{N} \cup \{\infty\}$.

Definition 38. For a set P of palindromic words, we define

$$\mathcal{CPal}(P) = \{W \in \mathcal{A}^* \mid \text{There exist } U \in P \text{ and integers } i, j \text{ with } \\ 1 \leq i \leq j \leq |U|, \ i + j = |U| + 1 \text{ and } W = U[i, j]\}.$$

Remark 39. It is obvious from Definition 38 that any element of $\mathcal{CPal}(P)$ is a palindromic factor of some word of P , but there may exist other palindromic factors of words of P which do not belong to $\mathcal{CPal}(P)$ as is seen from the following example.

Example 40. Let $P = \{1213121, 33433\}$. Then by Definition 38,

$$\mathcal{CPal}(P) = \{1213121, 21312, 131, 3, 33433, 343, 4\}.$$

Note that the words 1, 2, 33 and 121 are palindromic factors of some words of P but they are not elements of $\mathcal{CPal}(P)$.

Lemma 41. Let $U \in \mathcal{A}^*$ and $P \subseteq \text{Pal}(U)$. Then $\mathcal{CPal}(P) \subseteq \text{Pal}(U)$.

Proof. Since any element of P is a palindrome, by Definition 38, any element $W \in \mathcal{CPal}(P)$ is also a palindromic factor of U , whence the result follows. \square

Lemma 42. For any integer $k > 2$ we have

$$\text{Pal}(W^{(k)}) = \bigcup_{i \geq 0} (ki \oplus (\bigcup_{j=1}^{3k-2} \text{Pal}(W_j^{(k)}))).$$

Proof. Let $1 \leq j \leq 3k - 2$. By Equation (6), we have $(k \oplus W_j^{(k)}) \triangleright W_{j+k}^{(k)}$. Hence, using induction on i , for every nonnegative integer i we provide

$$(ki \oplus W_j^{(k)}) \triangleright W_{j+ki}^{(k)}.$$

Now, since $W_{j+ki}^{(k)} \prec W^{(k)}$ we conclude that

$$\bigcup_{i \geq 0} (ki \oplus (\bigcup_{j=1}^{3k-2} \mathcal{P}al(W_j^{(k)}))) \subseteq \mathcal{P}al(W^{(k)}).$$

To complete the proof, we need to show that

$$\mathcal{P}al(W_n^{(k)}) \subseteq \bigcup_{i \geq 0} (ki \oplus (\bigcup_{j=1}^{3k-2} \mathcal{P}al(W_j^{(k)}))), \quad (17)$$

holds for all integers n . But since $m < n$ implies that $W_m^{(k)} \prec W_n^{(k)}$ and $\mathcal{P}al(W_m^{(k)}) \subset \mathcal{P}al(W_n^{(k)})$, it suffices to prove that Equation (17) holds from a point on, say that it holds for all $n > 3k - 3$. To prove this, we use strong induction on n . The basis step $n = 3k - 2$ is obviously true because $\mathcal{P}al(W_{3k-2}^{(k)})$ appears in the right side of Equation (17). For the inductive step, let $n > 3k - 2$ and assume that Equation (17) holds for all integers j , $3k - 2 \leq j < n$. Since $3k - 2 < n$, using Lemma 25 and Theorem 35, $W_n^{(k)}$ has neither a straddling palindrome nor a bordering palindrome. Hence, by Equation (6), we have

$$\mathcal{P}al(W_n^{(k)}) = \bigcup_{j=n-k+1}^{n-1} \mathcal{P}al(W_j^{(k)}) \bigcup (k \oplus (\mathcal{P}al(W_{n-k}^{(k)}))).$$

By the induction hypothesis we have

$$\bigcup_{j=n-k+1}^{n-1} \mathcal{P}al(W_j^{(k)}) \bigcup (k \oplus (\mathcal{P}al(W_{n-k}^{(k)}))) \subseteq \bigcup_{i \geq 0} (ki \oplus (\bigcup_{j=1}^{3k-2} \mathcal{P}al(W_j^{(k)}))).$$

Therefore, Equation (17) holds for $j = n$, as desired. \square

To present the next results we need the following definition.

Definition 43. Let $k > 2$. We define the following sets of words:

$$\begin{aligned} \mathcal{P}_1^{(k)} &:= \{ki \oplus (W_n^{(k)} n^{-1}) : 1 \leq n \leq k - 1, i \geq 0\}, \\ \mathcal{P}_2^{(k)} &:= \{ki \oplus ((W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R j (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})) : \\ &\quad k \leq n \leq 2k - 3, n - k + 2 \leq j \leq k - 1, i \geq 0\}, \\ \mathcal{P}_3^{(k)} &:= \{ki \oplus (W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)} (n - 2k + 1)^{-1}), ki \oplus (W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)} (n - 2k + 1)^{-1}) : \\ &\quad 2k \leq n \leq 3k - 3, i \geq 1\}, \\ \mathcal{P}_4^{(k)} &:= \{ki \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} (k - 1)^{-1}), ki \oplus (0^{-1} W_{k-1}^{(k)} W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k - 1))^{-1}), ki \oplus (00) : i \geq 1\}. \end{aligned}$$

Lemma 44. Let $k > 2$. Then $\mathbb{N} \subseteq \mathcal{C}Pal(\mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)})$.

Proof. Let $m \in \mathbb{N}$ and $m = (j \bmod k)$. If $j = 0$ or $j = 1$, then by Definition 43, we have $0, 010 \in \mathcal{P}_1^{(k)}$ and hence $m \in \mathcal{C}Pal(\mathcal{P}_1^{(k)})$. Otherwise $2 \leq j \leq k - 1$ and using Lemma 23, for every integer n satisfying $k \leq n \leq k + j - 2$ we have

$$(W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R j (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)}) \in \mathcal{P}al(W_n^{(k)}).$$

Hence, $j \in \mathcal{C}Pal(\mathcal{P}_2^{(k)})$ and by Definition 43, $m \in \mathcal{C}Pal(\mathcal{P}_2^{(k)})$ as well. \square

Lemma 45. Let $k > 2$. Then $\bigcup_{i=1}^4 \mathcal{P}_i^{(k)} \subseteq \mathcal{Pal}(W^{(k)})$.

Proof. Using Lemma 18,

$$\{W_n^{(k)} n^{-1} : 1 \leq n \leq k-1\} \subseteq \mathcal{Pal}(W^{(k)}). \quad (18)$$

By Lemma 23, the set

$$\{(W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)})^R j (W_{j-1}^{(k)} W_{j-2}^{(k)} \cdots W_{n-k+1}^{(k)}) : \\ k \leq n \leq 2k-3, n-k+2 \leq j \leq k-1\} \subseteq \mathcal{Pal}(W^{(k)}). \quad (19)$$

By Lemma 31,

$$\{k \oplus (W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)} (n-2k+1)^{-1}), k \oplus (W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)} W_{n-2k+1}^{(k)} (n-2k+1)^{-1}) : \\ 2k \leq n \leq 3k-3\} \subseteq \mathcal{Pal}(W^{(k)}). \quad (20)$$

Finally, by Lemmas 27 and 34,

$$\{k \oplus (W_{k-1}^{(k)} W_{k-1}^{(k)} (k-1)^{-1}), k \oplus (0^{-1} W_{k-1}^{(k)} W_{k-1}^{(k)} W_{k-1}^{(k)} (0(k-1))^{-1}), k.k\} \subseteq \mathcal{Pal}(W^{(k)}). \quad (21)$$

On the other hand, using Equation (6), we conclude that if $P \in \mathcal{Pal}(W^{(k)})$, then for every integer $i \geq 0$, $ki \oplus P \in \mathcal{Pal}(W^{(k)})$. Therefore, using Equations (18)-(21), we conclude. \square

The following six lemmas give the structure of the palindromes of $W_n^{(k)}$, when $n \leq 3k-2$.

Lemma 46. Let $k > 2$, $1 \leq n < k$ and P be a maximal palindromic factor of $W_n^{(k)}$ of length at least 2. Then $P \in \mathcal{P}_1^{(k)}$.

Proof. By Lemma 18, $W_n^{(k)} n^{-1} = W_{n-1}^{(k)} (n-1)^{-1} (n-1) W_{n-1}^{(k)} (n-1)^{-1}$ is a maximal palindrome with center $n-1$. Therefore, it is easy to see that the maximal palindromes appearing in $W_n^{(k)}$ are either equal to $W_n^{(k)} n^{-1}$ or are a maximal palindrome of $W_{n-1}^{(k)} (n-1)^{-1}$. Hence, using induction we can see that the set of all maximal palindromes of $W_n^{(k)}$ is $\{W_i^{(k)} i^{-1} : 2 \leq i \leq n\} \subset \mathcal{P}_1^{(k)}$. \square

Lemma 47. Let $2 < k \leq n \leq 2k-3$ and P be a maximal palindromic factor of $W_n^{(k)}$ of length at least 2, then $P \in \mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)}$.

Proof. We prove the result using induction on n . For the basis step let $n = k$ and $P \in \mathcal{P}_2^{(k)} \setminus \mathcal{P}_1^{(k)}$ be a maximal palindromic factor of $W_n^{(k)}$. Since $P \notin \mathcal{P}_1^{(k)}$, using Lemma 46, we conclude that P is not an included palindromic factor of $W_n^{(k)}$. On the other hand by Theorem 35, P is not a straddling factor of $W_n^{(k)}$. Hence, P is a maximal bordering factor of $W_n^{(k)}$. Now, using Lemma 23, we provide $P \in \mathcal{P}_2^{(k)}$.

For the inductive step, let $k < n \leq 3k - 3$ and suppose that the result is true for all j , $k \leq j < n$. Let P be a maximal palindromic factor of $W_n^{(k)}$. Then by Theorem 35, either P is an included maximal palindromic factor of $W_n^{(k)}$ or it is a bordering maximal palindromic factor of $W_n^{(k)}$. In the former case, using Lemma 46 and induction hypothesis we conclude that $P \in \mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)}$. In the latter case, by Lemma 23 we have $P \in \mathcal{P}_2^{(k)}$, as desired. \square

Lemma 48. *Let $k > 2$ and $n = 2k - 2$ and P be a maximal palindromic factor of $W_n^{(k)}$ of length at least 2. Then $P \in \mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)}$.*

Proof. By Lemma 25 and Theorem 35, $W_n^{(k)}$ has neither a straddling palindrome nor a bordering palindrome. Hence, using Equation (6), Lemmas 46 and 47, $\text{Pal}(W_n^{(k)}) \subseteq (\mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)})$, as desired. \square

Lemma 49. *Let $k > 2$ and $n = 2k - 1$ and P be a maximal palindromic factor of $W_n^{(k)}$ of length at least 2. Then either $P = k.k$ or $P \in \mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)}$.*

Proof. Let $P \notin \mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)}$. Then by Lemmas 46-48, P is not a maximal palindromic factor of any $W_j^{(k)}$, $j < n$. Hence, P is either a straddling palindrome or a bordering palindrome. By Lemmas 22 and 27, P is straddling and $P = k.k$. \square

Lemma 50. *Let $k > 2$ and $2k \leq n \leq 3k - 3$ and P be a maximal palindromic factor of $W_n^{(k)}$ of length at least 2. Then $P \in (\mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)} \cup \mathcal{P}_3^{(k)} \cup \{ki.ki : i \geq 1\})$.*

Proof. We use induction on n . First, let $n = 2k$ and let P be a maximal palindromic factor of $W_{2k}^{(k)}$. Then by Lemma 25, P is not a bordering factor, hence, it is either an included palindromic factor or a straddling palindromic factor of $W_{2k}^{(k)}$. In the former case, by Lemmas 46-49, $P \in (\mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)} \cup \{ki.ki : i \geq 1\})$, while in the latter case, by Lemma 31, $P \in \mathcal{P}_3^{(k)}$. This terminates the basis step of the induction.

For the inductive step, consider any integer n , $2k < n \leq 3k - 3$, and assume that the lemma holds for all integers j , $2k \leq j < n$. To conclude the validity of lemma for n , note that a maximal palindromic factor P of $W_n^{(k)}$ is either an included or a straddling palindrome, by Lemma 25. In the former case, by induction hypothesis, $P \in (\mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)} \cup \mathcal{P}_3^{(k)} \cup \{ki.ki : i \geq 1\})$, while in the latter case, by Lemma 31, $P \in \mathcal{P}_3^{(k)}$. Hence, we are done. \square

Lemma 51. *Let $k > 2$, $n = 3k - 2$ and P be a maximal palindromic factor of $W_n^{(k)}$ of length at least 2. Then $P \in \bigcup_{i=1}^4 \mathcal{P}_i^{(k)}$.*

Proof. Let $P \notin (\mathcal{P}_1^{(k)} \cup \mathcal{P}_2^{(k)} \cup \mathcal{P}_3^{(k)} \cup \{ki.ki : i \geq 1\})$. Thus, by Lemmas 46-50, P is not a maximal palindromic factor of any $W_j^{(k)}$, $j < n$. Hence, P is either a straddling palindrome or a bordering palindrome. By Lemma 25, $W_n^{(k)}$ has no bordering palindrome and hence it is a maximal straddling palindrome of $W_n^{(k)}$. By Lemma 34, we conclude that $P \in \mathcal{P}_4^{(k)}$. \square

Lemma 52. Let $k > 2$ and let P be a maximal palindromic factor of $W^{(k)}$ of length at least 2. Then $P \in \bigcup_{i=1}^4 \mathcal{P}_i^{(k)}$.

Proof. By Lemma 42, there exist $i \geq 0$ and $1 \leq j \leq 3k - 2$ such that P' is a maximal palindrome of $W_j^{(k)}$ and $P = ki \oplus P'$. Now, using Lemmas 46-51, $P' \in \bigcup_{i=1}^4 \mathcal{P}_i^{(k)}$. Since $ki \oplus (\bigcup_{i=1}^4 \mathcal{P}_i^{(k)}) = \bigcup_{i=1}^4 \mathcal{P}_i^{(k)}$, we have $P \in \bigcup_{i=1}^4 \mathcal{P}_i^{(k)}$. \square

Theorem 53. For any integer $k > 2$ we have

$$\mathcal{P}al(W^{(k)}) = \mathcal{C}P\mathcal{a}l\left(\bigcup_{i=1}^4 \mathcal{P}_i^{(k)}\right). \quad (22)$$

Proof. First we prove that the left side of Equation (22) is a subset of its right side. For this, consider any element $P \in \mathcal{P}al(W^{(k)})$. If $|P| = 1$, then by Lemma 44, $P \in \bigcup_{i=1}^4 \mathcal{C}P\mathcal{a}l(\mathcal{P}_i^{(k)})$. Otherwise, consider a maximal palindromic factor U of $W^{(k)}$ such that $P \in \mathcal{C}P\mathcal{a}l(\{U\})$. By Lemma 52, $U \in \bigcup_{i=1}^4 \mathcal{P}_i^{(k)}$. therefore, using Lemma 41, $P \in \bigcup_{i=1}^4 \mathcal{C}P\mathcal{a}l(\mathcal{P}_i^{(k)})$ as required.

To prove that the right side of Equation (22) is a subset of its left side, note that by Lemma 45, $\bigcup_{i=1}^4 \mathcal{P}_i^{(k)} \subseteq \mathcal{P}al(W^{(k)})$. Thus using Lemma 41, $\mathcal{C}P\mathcal{a}l(\bigcup_{i=1}^4 \mathcal{P}_i^{(k)}) \subseteq \mathcal{P}al(W^{(k)})$. \square

Example 54. Consider the word $W^{(3)}$. Then by Definition 43, the sets $\mathcal{P}_1^{(3)}, \mathcal{P}_2^{(3)}, \mathcal{P}_3^{(3)}, \mathcal{P}_4^{(3)}$ are as follows:

$$\begin{aligned} \mathcal{P}_1^{(3)} &= \{3i, 3i \oplus (010) : i \geq 0\}, \\ \mathcal{P}_2^{(3)} &= \{3i \oplus (10201) : i \geq 0\}, \\ \mathcal{P}_3^{(3)} &= \{3i \oplus (010), 3i \oplus (01010) : i \geq 1\}, \\ \mathcal{P}_4^{(3)} &= \{3i \oplus (0102010), 3i \oplus (102010201), 3i \oplus (00) : i \geq 1\}. \end{aligned}$$

By Theorem 53, $\mathcal{P}al(W^{(3)}) = \mathcal{C}P\mathcal{a}l\left(\bigcup_{i=1}^4 \mathcal{P}_i^{(3)}\right)$.

Example 55. Consider the word $W^{(4)}$. Then by Definition 43, the sets $\mathcal{P}_1^{(4)}, \mathcal{P}_2^{(4)}, \mathcal{P}_3^{(4)}, \mathcal{P}_4^{(4)}$ are as follows:

$$\mathcal{P}_1^{(4)} = \{4i, 4i \oplus (010), 4i \oplus (0102010) : i \geq 0\},$$

$$\begin{aligned} \mathcal{P}_2^{(4)} &= \{4i \oplus (10201), 4i \oplus (1020103010201), 4i \oplus (201030102) : i \geq 0\}, \\ \mathcal{P}_3^{(4)} &= \{4i \oplus (010), 4i \oplus (01010), 4i \oplus (0102010), 4i \oplus (01020102010) : i \geq 1\}, \\ \mathcal{P}_4^{(4)} &= \{4i \oplus (010201030102010), 4i \oplus (102010301020103010201), 4i \oplus (00) : i \geq 1\}. \end{aligned}$$

By Theorem 53, $\mathcal{P}al(W^{(4)}) = \mathcal{CP}al(\bigcup_{i=1}^4 \mathcal{P}_i^{(4)})$.

Example 56. Consider the word $W^{(5)}$. Then by Definition 43, the sets $\mathcal{P}_1^{(5)}, \mathcal{P}_2^{(5)}, \mathcal{P}_3^{(5)}, \mathcal{P}_4^{(5)}$ are as follows:

$$\begin{aligned} \mathcal{P}_1^{(5)} &= \{5i, 5i \oplus (010), 5i \oplus (0102010), 5i \oplus (010201030102010) : i \geq 0\}, \\ \mathcal{P}_2^{(5)} &= \{5i \oplus (10201), 5i \oplus (1020103010201), 5i \oplus (10201030102010401020103010201), \\ &\quad 5i \oplus (201030102), 5i \oplus (2010301020104010201030102), 5i \oplus (30102010401020103) : i \geq 0\}, \\ \mathcal{P}_3^{(5)} &= \{5i \oplus (010), 5i \oplus (01010), 5i \oplus (0102010), 5i \oplus (01020102010), 5i \oplus (010201030102010), \\ &\quad 5i \oplus (01020103010201030102010) : i \geq 1\}, \\ \mathcal{P}_4^{(5)} &= \{5i \oplus (00), 5i \oplus (102010301020104010201030102010401020103010201), \\ &\quad 5i \oplus (0102010301020104010201030102010) : i \geq 1\}. \end{aligned}$$

By Theorem 53, $\mathcal{P}al(W^{(5)}) = \mathcal{CP}al(\bigcup_{i=1}^4 \mathcal{P}_i^{(5)})$.

5.1 Length of Palindromes in $W^{(k)}$

In this section, we want to compute all possible values for the lengths of palindromes in $W^{(k)}$.

Definition 57. Let $k > 2$, and for $1 \leq i \leq 4$. Let $\mathcal{P}_i^{(k)}$ be the sets given in Definition 43. For $1 \leq i \leq 4$, we define $L(\mathcal{P}_i^{(k)}) := \{|P| : UPU^R \in \mathcal{P}_i^{(k)}, |P| > 1 \text{ and } U \in \mathbb{N}^*\}$.

Lemma 58. For each integer $k \geq 3$

$$\begin{aligned} L(\mathcal{P}_1^{(k)}) &:= \{2i - 1 : 1 \leq i \leq 2^{k-2}\}, \\ L(\mathcal{P}_2^{(k)}) &:= \{2i - 1 : 2 \leq i \leq 2^{k-1} - 1\}, \\ L(\mathcal{P}_3^{(k)}) &:= \{2i - 1 : 2 \leq i \leq 3 \cdot 2^{k-3}\}, \\ L(\mathcal{P}_4^{(k)}) &:= \{2, 2i - 1 : 2 \leq i \leq 3 \cdot 2^{k-2} - 1\}. \end{aligned}$$

Proof. We just prove $L(\mathcal{P}_1^{(k)}) = \{2i - 1 : 1 \leq i \leq 2^{k-2}\}$, the proof of the rest parts are similar to this case. By Definition 43, it is clear that the set $L(\mathcal{P}_1^{(k)})$ just contains odd integers. Again by Definition 43, it can be seen that if $t \in L(\mathcal{P}_1^{(k)})$ is an odd number greater than 2, then $t - 2 \in L(\mathcal{P}_1^{(k)})$. So if ℓ_1 is the maximum integer in $L(\mathcal{P}_1^{(k)})$,

then $L(\mathcal{P}_1^{(k)}) = \{1, 3, 5, \dots, \ell_1\}$. Therefore, to prove this part it suffices to show that $\ell_1 = 2^{k-1} - 1$. By Definition 43 and Corollary 8, we have

$$\ell_1 = |W_{k-1}^{(k)}| - 1 = 2^{k-1} - 1. \quad \square$$

Theorem 59. *For every integer $k \geq 3$, the palindrome complexity of $W^{(k)}$ is given by*

$$\text{pal}_{W^{(k)}}(n) = \begin{cases} \infty & \text{if } n \in \{2\} \cup \{2i - 1 : 1 \leq i \leq 3 \cdot 2^{k-2} - 1\}, \\ 0 & \text{otherwise.} \end{cases}$$

Proof. Let $A = \{2\} \cup \{2i - 1 : 1 \leq i \leq 3 \cdot 2^{k-2} - 1\}$. Using Theorem 53 and Lemma 58, we find that if $n \notin A$, then there is no palindromic factor in $W^{(k)}$ of length n , in other words $\text{pal}_{W^{(k)}}(n) = 0$. If $n \in A$, then by Theorem 53 and Lemma 58, $W^{(k)}$ has at least one palindromic factor of length n or equivalently $\text{pal}_{W^{(k)}}(n) \neq 0$. By Definitions 43 and 57, it is easy to see that in this case $\text{pal}_{W^{(k)}}(n) = \infty$. \square

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