

# A generalization of the Bollobás set pairs inequality

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## Abstract

The Bollobás set pairs inequality is a fundamental result in extremal set theory with many applications. In this paper, for  $n \geq k \geq t \geq 2$ , we consider a collection of  $k$  families  $\mathcal{A}_i : 1 \leq i \leq k$  where  $\mathcal{A}_i = \{A_{i,j} \subset [n] : j \in [n]\}$  so that  $A_{1,i_1} \cap \dots \cap A_{k,i_k} \neq \emptyset$  if and only if there are at least  $t$  distinct indices  $i_1, i_2, \dots, i_k$ . Via a natural connection to a hypergraph covering problem, we give bounds on the maximum size  $\beta_{k,t}(n)$  of the families with ground set  $[n]$ .

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## 1 Introduction

A central topic of study in extremal set theory is the maximum size of a family of subsets of an  $n$ -element set subject to restrictions on their intersections. Classical theorems in the area are discussed in Bollobás [2]. In this paper, we generalize one such theorem, known as the Bollobás set pairs inequality or two families theorem [3]:

**Theorem 1.** (Bollobás) *Let  $\mathcal{A} = \{A_1, A_2, \dots, A_m\}$  and  $\mathcal{B} = \{B_1, B_2, \dots, B_m\}$  be families of finite sets, such that  $A_i \cap B_j \neq \emptyset$  if and only if  $i, j \in [m]$  are distinct. Then*

$$\sum_{i=1}^m \binom{|A_i \cup B_i|}{|A_i|}^{-1} \leq 1. \quad (1)$$

For convenience, we refer to a pair of families  $\mathcal{A}$  and  $\mathcal{B}$  satisfying the conditions of Theorem 1 as a *Bollobás set pair*. The inequality above is tight, as we may take the pairs  $(A_i, B_i)$  to be distinct partitions of a set of size  $a + b$  with  $|A_i| = a$  and  $|B_i| = b$  for  $1 \leq i \leq \binom{a+b}{a}$ .

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The latter inequality was proved for  $a = 2$  by Erdős, Hajnal and Moon [5], and in general has a number of different proofs [11, 12, 14, 17, 18]. A geometric version was proved by Lovász [17, 18], who showed that if  $A_1, A_2, \dots, A_m$  and  $B_1, B_2, \dots, B_m$  are respectively  $a$ -dimensional and  $b$ -dimensional subspaces of a linear space and  $\dim(A_i \cap B_j) = 0$  if and only if  $i, j \in [m]$  are distinct, then  $m \leq \binom{a+b}{a}$ .

## 1.1 Main Theorem

Theorem 1 has been generalized in a number of different directions in the literature [6, 9, 13, 16, 21, 24]. In this paper, we give a generalization of Theorem 1 from the case of two families to  $k \geq 3$  families of sets with conditions on the  $k$ -wise intersections. For  $2 \leq t \leq k$ , a *Bollobás  $(k, t)$ -tuple* is a sequence  $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$  of set families  $\mathcal{A}_j = \{A_{j,i} : 1 \leq i \leq m\}$  where  $\bigcap_{j=1}^k A_{j,i_j} \neq \emptyset$  if and only if at least  $t$  of the indices  $i_1, i_2, \dots, i_k$  are distinct. We refer to  $m$  as the *size* of the Bollobás  $(k, t)$ -tuple. Let  $[m]_{(t)}$  denote the set of sequences of  $t$  distinct elements of  $[m]$  and fix a surjection  $\phi : [k] \rightarrow [t]$ . For  $\sigma \in [m]_{(t-1)}$ , set  $\sigma(t) = \sigma(1)$  and define  $A_{1,\sigma}(\phi) = \bigcap_{j:\phi(j)=1} A_{j,\sigma(1)}$  and, for  $2 \leq j \leq t$ , we define

$$A_{j,\sigma}(\phi) = \bigcap_{h:\phi(h)=j} A_{h,\sigma(j)} \setminus \bigcup_{h=1}^{j-1} A_{h,\sigma}(\phi).$$

Using this notation, we generalize (1) as follows:

**Theorem 2.** *Let  $k \geq t \geq 2$  and  $m \geq t$ , let  $\phi : [k] \rightarrow [t]$  be a surjection, and let  $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$  be a Bollobás  $(k, t)$ -tuple of size  $m$ . Then*

$$\sum_{\sigma \in [m]_{(t-1)}} \left( \frac{|\bigcup_{j=1}^t A_{j,\sigma}(\phi)|}{|A_{1,\sigma}(\phi)| |A_{2,\sigma}(\phi)| \cdots |A_{t,\sigma}(\phi)|} \right)^{-1} \leq 1. \quad (2)$$

We show in Section 2.1 that this inequality is tight for all  $k \geq t = 2$ , but do not have an example to show that this inequality is tight for any  $t > 2$ .

For  $n \geq k \geq t \geq 2$ , let  $\beta_{k,t}(n)$  denote the maximum  $m$  such that there exists a Bollobás  $(k, t)$ -tuple of size  $m$  consisting of subsets of  $[n]$ . Then (1) gives  $\beta_{2,2}(n) \leq \binom{n}{\lfloor n/2 \rfloor}$  which is tight for all  $n \geq 2$ . Letting  $H(q) = -q \log_2 q - (1-q) \log_2(1-q)$  denote the standard binary entropy function, we prove the following theorem:

**Theorem 3.** *For  $k \geq 3$  and large enough  $n$ ,*

$$\frac{1}{k} \leq \frac{\log_2 \beta_{k,2}(n)}{n} \leq H\left(\frac{1}{k}\right) \leq \frac{\log_2(ke)}{k}. \quad (3)$$

*For  $k \geq t \geq 3$  and large enough  $n$ ,*

$$\frac{\log_2 e}{\binom{k}{t-1}(t+1)t^{t-1}} \leq \frac{\log_2 \beta_{k,t}(n)}{n} \leq \frac{2}{\binom{k}{t-1}(t-1)^{t-3}}. \quad (4)$$

This determines  $\log_2 \beta_{k,2}(n)$  up to a factor of order  $\log_2 k$  and  $\log_2 \beta_{k,t}(n)$  up to a factor of order  $t^3$ . We leave it as an open problem to determine the asymptotic value of  $(\log_2 \beta_{k,t}(n))/n$  as  $n \rightarrow \infty$  for any  $k \geq 3$  and  $t \geq 2$ . A natural source for lower bounds on  $\beta_{k,t}(n)$  comes from the probabilistic method – see the random constructions in Section 3.1 which establish the lower bounds in Theorem 3. To prove Theorem 3, we use a natural connection to hypergraph covering problems.

## 1.2 Covering hypergraphs

Theorem 1 has a wide variety of applications, from saturation problems [3, 19] to covering problems for graphs [11, 20], complexity of 0-1 matrices [23], geometric problems [1], counting cross-intersecting families [7], crosscuts and transversals of hypergraphs [24, 25, 26], hypergraph entropy [15, 22], and perfect hashing [8, 10]. In this section, we give an application of our main results to hypergraph covering problems. For a  $k$ -uniform hypergraph  $H$ , let  $f(H)$  denote the minimum number of complete  $k$ -partite  $k$ -uniform hypergraphs whose union is  $H$ . In the case of graph covering, a simple connection to the Bollobás set pairs inequality (1) may be described as follows. Let  $K_{n,n} \setminus M$  denote the complement of a perfect matching  $M = \{x_i y_i : 1 \leq i \leq n\}$  in the complete bipartite graph  $K_{n,n}$  with parts  $X = \{x_1, x_2, \dots, x_n\}$  and  $Y = \{y_1, y_2, \dots, y_n\}$ . If  $H_1, H_2, \dots, H_m$  are complete bipartite graphs in a minimum covering of  $K_{n,n} \setminus M$ , then let  $A_i = \{j : x_j \in V(H_i)\}$  and  $B_i = \{j : y_j \in V(H_i)\}$ . Setting  $\mathcal{A} = \{A_i\}_{i \in [m]}$  and  $\mathcal{B} = \{B_i\}_{i \in [m]}$ , it is straightforward to check that  $(\mathcal{A}, \mathcal{B})$  is a Bollobás set pair, and Theorem 1 applies to give

$$f(K_{n,n} \setminus M) = \min\left\{m : \binom{m}{\lceil m/2 \rceil} \geq n\right\}. \quad (5)$$

In a similar way, Theorem 2 applies to covering complete  $k$ -partite  $k$ -uniform hypergraphs. Let  $K_{n,n,\dots,n}$  denote the complete  $k$ -partite  $k$ -uniform hypergraph with parts  $X_i = \{x_{ij} : j \in [n]\}$  for  $i \in [k]$ . Let  $H_{k,t}(n)$  denote the subhypergraph consisting of hyperedges  $\{x_{1,i_1}, x_{2,i_2}, \dots, x_{k,i_k}\}$  such that at least  $t$  of the indices  $i_1, i_2, \dots, i_k$  are distinct, and set  $f_{k,t}(n) = f(H_{k,t}(n))$ . Then there is a one-to-one correspondence between Bollobás  $(k, t)$ -tuples of subsets of  $[m]$  and coverings of  $H_{k,t}(n)$  with  $m$  complete  $k$ -partite  $k$ -graphs. We let  $\beta_{k,t}(m)$  be the maximum size of a Bollobás  $(k, t)$ -tuple of subsets of  $[m]$ , so that

$$f_{k,t}(n) = \min\{m : \beta_{k,t}(m) \geq n\}. \quad (6)$$

This correspondence together with Theorem 2 will be exploited to prove

$$f_{k,2}(n) \geq \min\left\{m : \binom{m}{\lceil m/k \rceil} \geq n\right\} \quad (7)$$

which is partly an analog of (5). More generally, we prove the following theorem:

**Theorem 4.** *For  $k \geq 3$  and large enough  $n$ ,*

$$\frac{k}{\log_2(ke)} \leq \frac{1}{H(\frac{1}{k})} \leq \frac{f_{k,2}(n)}{\log_2 n} \leq k. \quad (8)$$

For  $k \geq t \geq 3$  and large enough  $n$ ,

$$\binom{k}{t-1} \frac{(t-1)^{t-3}}{2} \leq \frac{f_{k,t}(n)}{\log_2 n} \leq \frac{(t+1)t^{t-1}}{\log_2 e} \binom{k}{t-1}. \quad (9)$$

The bounds on  $\beta_{k,t}(n)$  in Theorem (3) follow immediately from this theorem and (6). Equation (9) gives the order of magnitude for each  $t \geq 3$  as  $k \rightarrow \infty$ , but for  $t = 2$ , Equation (8) has a gap of order  $\log_2 k$ . From (7), we obtain  $\beta_{k,2}(n) \leq \binom{n}{\lfloor n/k \rfloor}$ . It is perhaps unsurprising that the asymptotic value of  $f_{k,t}(n)/\log_2 n$  as  $n \rightarrow \infty$  is not known for any  $k > 2$ , since a limiting value of  $f(K_n^k)/\log_2 n$  is not known for any  $k > 2$  – see Körner and Marston [15] and Guruswami and Riazanov [10].

### 1.3 Organization and notation

Given a subset  $A \subset [n]$ , let  $A^c := [n] \setminus A$  be the complement of  $A$  in  $[n]$ . For positive integers  $k \leq n$ , let  $(n)_{(k)} = (n)(n-1)\cdots(n-k+1)$  denote the falling factorial. This paper is organized as follows. In Section 2, we prove Theorem 2. In Section 2.1, we construct a Bollobás  $(k, 2)$ -tuple which achieves equality in Theorem 2 and in Section 2.2, we construct a Bollobás  $(k, 2)$ -tuple which gives the lower bound in Equation (3). The upper bound on  $f_{k,t}(n)$  in Theorem 4 comes from a probabilistic construction in Section 3.1, and the proof of the lower bound on  $f_{k,t}(n)$  is given in Section 3.3; we prove (7) in Section 3.2.

## 2 Proof of Theorem 2

Given a Bollobás set  $(k, t)$ -tuple  $(\mathcal{A}_1, \dots, \mathcal{A}_k)$  with  $\mathcal{A}_j = \{A_{j,i} : 1 \leq i \leq m\}$  and a surjection  $\phi : [k] \rightarrow [t]$ , consider  $\mathcal{A}_\ell(\phi) : 1 \leq \ell \leq t$  where  $\mathcal{A}_\ell(\phi) = \{A_{\ell,i}(\phi) : 1 \leq i \leq m\}$  and

$$A_{\ell,i}(\phi) = \bigcap_{h:\phi(h)=\ell} A_{h,i}.$$

It follows that  $(\mathcal{A}_1(\phi), \dots, \mathcal{A}_t(\phi))$  is a Bollobás set  $(t, t)$ -tuple and hence it suffices to prove Theorem 2 in the case where  $t = k$ . In this setting, surjections  $\phi : [k] \rightarrow [k]$  simply permute the  $k$  families and as such we suppress the notation of  $\phi$  for the remainder of this section. One of the proofs of Theorem 1, given a Bollobás set pair, defines a collection of chains  $\mathcal{C}_i$  for  $i \in [m]$  and shows that these chains are necessarily disjoint. Similarly, given a Bollobás set  $(k, k)$ -tuple, we will define a collection of chains  $\mathcal{C}_\sigma$  for every ordered collection  $\sigma$  of  $(k-1)$  distinct indices of  $[m]$  and show these chains are pairwise disjoint.

Let  $(\mathcal{A}_1, \dots, \mathcal{A}_k)$  with  $\mathcal{A}_j = \{A_{j,i} : 1 \leq i \leq m\}$  be a Bollobás set  $(k, k)$ -tuple, and set

$$X = \bigcup_{i=1}^m (A_{1,i} \cup A_{2,i} \cup \dots \cup A_{k,i})$$

with  $|X| = n$ . For  $\sigma \in [m]_{(k-1)}$ , define a subset  $\mathcal{C}_\sigma$  of permutations  $\pi : X \rightarrow [n]$  by

$$\mathcal{C}_\sigma := \left\{ \pi : X \rightarrow [n] : \max_{x \in A_{1,\sigma}} \pi(x) < \min_{y \in A_{2,\sigma}} \pi(y) \leq \max_{y \in A_{2,\sigma}} \pi(y) < \cdots < \min_{z \in A_{k,\sigma}} \pi(z) \right\}.$$

Letting  $U_\sigma := A_{1,\sigma} \cup \cdots \cup A_{k,\sigma}$ , elementary counting methods give

$$|\mathcal{C}_\sigma| = \binom{n}{|U_\sigma|} |A_{1,\sigma}|! \cdots |A_{k,\sigma}|! (n - |U_\sigma|)! = n! \cdot \binom{|U_\sigma|}{|A_{1,\sigma}| \cdots |A_{k,\sigma}|}^{-1}. \quad (10)$$

We will now prove a lemma which states that  $\{\mathcal{C}_\sigma\}_{\sigma \in [m]_{(k-1)}}$  forms a disjoint collection of a permutations. The general proof only works for  $k \geq 4$ , so we first consider  $k = 3$ .

**Lemma 5.** *If  $\sigma_1, \sigma_2 \in [m]_{(2)}$  are distinct, then  $\mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2} = \emptyset$ .*

*Proof.* Seeking a contradiction, suppose there exists  $\pi \in \mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2}$ . After relabeling, it suffices to consider the following five cases.

- (1)  $\sigma_1 = \{1, 3\}$  and  $\sigma_2 = \{2, 4\}$
- (2)  $\sigma_1 = \{1, 3\}$  and  $\sigma_2 = \{2, 3\}$
- (3)  $\sigma_1 = \{1, 2\}$  and  $\sigma_2 = \{1, 3\}$
- (4)  $\sigma_1 = \{1, 2\}$  and  $\sigma_2 = \{2, 3\}$
- (5)  $\sigma_1 = \{1, 2\}$  and  $\sigma_2 = \{3, 1\}$ .

In case (1), without loss of generality,  $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,2}\}$  and thus  $\pi \in \mathcal{C}_{\sigma_2}$  yields

$$\max_{x \in A_{1,1}} \pi(x) \leq \max_{x \in A_{1,2}} \pi(x) < \min_{y \in A_{2,4} \setminus A_{1,2}} \pi(y).$$

Then as  $A_{1,1} \cap A_{2,4} \cap A_{3,2} \neq \emptyset$ , there exists  $w \in A_{1,1} \cap A_{2,4} \cap A_{3,2}$ . It follows that  $w \notin A_{1,2}$  since if  $w \in A_{1,2}$ , then  $w \in A_{1,2} \cap A_{2,4} \cap A_{3,2} \neq \emptyset$ ; a contradiction. But this yields a contradiction as

$$\pi(w) \leq \max_{x \in A_{1,1}} \pi(x) \leq \max_{x \in A_{1,2}} \pi(x) < \min_{y \in A_{2,4} \setminus A_{1,2}} \pi(y) \leq \pi(w).$$

In case (2), without loss of generality,  $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,2}\}$  and we recover a similar contradiction as case (1) by noting that there exists  $w \in A_{1,1} \cap A_{2,3} \cap A_{3,2}$  with  $w \notin A_{1,2}$ .

In case (3) we may assume  $\max\{\pi(x) : x \in A_{2,2} \setminus A_{1,1}\} \leq \max\{\pi(x) : x \in A_{2,3} \setminus A_{1,1}\}$  and  $\pi \in \mathcal{C}_{1,3}$  yields  $\max\{\pi(x) : x \in A_{2,3} \setminus A_{1,1}\} < \min\{\pi(x) : x \in A_{3,1} \setminus (A_{1,1} \cup A_{2,3})\}$ . Thus

$$\max\{\pi(x) : x \in A_{2,2} \setminus A_{1,1}\} < \min\{\pi(x) : x \in A_{3,1} \setminus (A_{1,1} \cup A_{2,3})\}$$

and there exists  $w \in A_{1,3} \cap A_{2,2} \cap A_{3,1}$  with  $w \notin A_{1,1}$  and  $w \notin A_{2,3}$ . It follows that  $\pi(w) < \pi(w)$ , a contradiction.

In case (4), if  $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,2}\}$ , then using  $w \in A_{1,1} \cap A_{2,3} \cap A_{3,2}$  and noting  $w \notin A_{1,2}$ , we get a contradiction. Thus, we may assume otherwise and  $\pi \in \mathcal{C}_{1,2}$  gives

$$\max_{x \in A_{1,2}} \pi(x) < \max_{x \in A_{1,1}} \pi(x) < \min_{z \in A_{3,1} \setminus (A_{1,1} \cup A_{2,2})} \pi(z).$$

This is a contradiction as there exists  $w \in A_{1,2} \cap A_{2,3} \cap A_{3,1}$  with  $w \notin A_{1,1}$  and  $w \notin A_{2,2}$ . In case (5), if  $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,3}\}$ , then we may proceed as in the latter part of case (4) using  $w \in A_{1,1} \cap A_{2,2} \cap A_{3,3}$  and  $w \notin A_{2,1}$  and  $w \notin A_{1,3}$  to get a contradiction. Otherwise, proceeding as in case (1) and noting there exists  $w \in A_{1,3} \cap A_{2,2} \cap A_{3,1}$ , but  $w \notin A_{1,1}$  yields a contradiction.  $\square$

A similar argument yields the analog of Lemma 5 to the case where  $k \geq 4$ .

**Lemma 6.** *Let  $k \geq 4$ . If  $\sigma_1, \sigma_2 \in [m]_{(k-1)}$  are distinct, then  $\mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2} = \emptyset$ .*

*Proof.* Since  $\sigma_1, \sigma_2 \in [m]_{(k-1)}$  are distinct, there exists minimal  $h \in [k-1]$  so that  $\sigma_1(h) \neq \sigma_2(h)$ . Seeking a contradiction, suppose there exists a  $\pi \in \mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2}$ . Without loss of generality,

$$\max\{\pi(x) : x \in A_{h,\sigma_1}\} \leq \max\{\pi(x) : x \in A_{h,\sigma_2}\} < \min\{\pi(z) : z \in A_{k,\sigma_2}\}.$$

Now, consider a bijection  $\tau : [k-1] \setminus \{h\} \rightarrow [k-1] \setminus \{1\}$  which has no fixed points. As in Lemma 5, we want to show that there exists a  $w \in A_{h,\sigma_1} \cap A_{k,\sigma_2}$  and consider two separate cases.

First, suppose that  $\sigma_1(h) \notin \sigma_2([k-1])$ . As  $|\{\sigma_1(h), \sigma_2(1), \dots, \sigma_2(k-1)\}| = k$ , there exists

$$w \in A_{h,\sigma_1(h)} \cap A_{k,\sigma_2(1)} \cap \bigcap_{l \in [k-1] \setminus \{h\}} A_{l,\sigma_2(\tau(l))}. \quad (11)$$

Next, suppose that  $\sigma_1(h) = \sigma_2(x)$  for some  $x$ . We now claim that  $x \neq 1$ . If  $h = 1$ , then this is trivial. If  $h > 1$ , then  $\sigma_1(1) = \sigma_2(1)$ , so  $\sigma_1(h) \neq \sigma_2(1)$  since  $\sigma_1(h) \neq \sigma_1(1)$ . For  $\tau$  as above, there exists  $y \in [k-1] \setminus \{h\}$  so that  $\tau(y) = x$ . Taking  $\gamma$  distinct from  $\{\sigma_2(1), \dots, \sigma_2(k-1)\} \setminus \{\sigma_2(x)\}$ ,  $|\{\sigma_1(h), \gamma, \sigma_2(1), \dots, \sigma_2(k-1)\} \setminus \{\sigma_2(x)\}| = k$  and hence there exists

$$w \in A_{h,\sigma_1(h)} \cap A_{k,\sigma_2(1)} \cap A_{y,\gamma} \cap \bigcap_{l \in [k-1] \setminus \{y,h\}} A_{l,\sigma_2(\tau(l))}. \quad (12)$$

By construction,  $w \in A_{h,\sigma_1(h)} \cap A_{k,\sigma_2(1)}$ . Suppose there exists a  $t \in [k-1] \setminus \{h\}$  so that  $w \in A_{t,\sigma_2(t)}$ . As  $\tau$  has no fixed points, replacing the set in the  $k$ -wise intersection corresponding to  $A_t$  with  $A_{t,\sigma_2(t)}$  in either (11) or (12),  $w$  is an element of this new  $k$ -wise intersection with  $(k-1)$  distinct indices; a contradiction. If  $w \in A_{h,\sigma_2(h)}$ , then we may similarly replace  $A_{h,\sigma_1(h)}$  with  $A_{h,\sigma_2(h)}$  in the  $k$ -wise intersection in either (11) or (12) to get a contradiction. Thus,  $w \notin A_{1,\sigma_2(1)} \cup \dots \cup A_{k-1,\sigma_2(k-1)}$  and hence  $w \in A_{h,\sigma_1} \cap A_{k,\sigma_2}$  so that  $\pi(w) < \pi(w)$ ; a contradiction.  $\square$

Using Equation (10), Lemma 5, and Lemma 6, we are now able to prove Theorem 2 in the case where  $t = k$ . There are  $n!$  total permutations, and Lemma 5 and Lemma 6 yield that each of which appears in at most one of the sets  $\mathcal{C}_\sigma$  for  $\sigma \in [m]_{(k-1)}$ . Hence, using  $|\mathcal{C}_\sigma|$  in Equation (10),

$$\sum_{\sigma \in [m]_{(k-1)}} |\mathcal{C}_\sigma| = \sum_{\sigma \in [m]_{(k-1)}} n! \cdot \left( \frac{|A_{1,\sigma} \cup \dots \cup A_{k,\sigma}|}{|A_{1,\sigma}| \cdots |A_{k,\sigma}|} \right)^{-1} \leq n!$$

and thus the result follows by dividing through by  $n!$ .

## 2.1 Sharpness of Theorem 2

We give a simple construction establishing the sharpness of Theorem 2 for  $k \geq t = 2$ . Let  $n \geq 4k$  and using addition modulo  $n$ , define  $A_{1,i} = \{i\}^c$ ,  $A_{j,i} = \{i - (j - 1), i + (j - 1)\}^c$  for  $j \in [2, k - 1]$ , and  $A_{k,i} = \{i - k + 2, i - k + 3, \dots, i + k - 2\}$ . Letting  $\mathcal{A}_j = \{A_{j,i}\}_{i \in [n]}$  for all  $j \in [k]$ , we will show  $(\mathcal{A}_1, \dots, \mathcal{A}_k)$  is a Bollobás  $(k, 2)$ -tuple. Since  $|A_{1,i}| = n - 1$  and  $|A_{2,i} \cap \dots \cap A_{k,i}| = 1$ , Theorem 2 with  $t = 2$  and surjection  $\phi : [k] \rightarrow [2]$  with  $\phi(1) = 1$  and  $\phi(i) = 2$  for  $i \neq 1$  gives

$$1 \geq \sum_{i=1}^n \left( \frac{|A_{1,i}| + |A_{2,i} \cap \dots \cap A_{k,i}|}{|A_{1,i}|} \right)^{-1} = \sum_{i=1}^n \frac{1}{n} = 1.$$

By construction, for all  $i \in [n]$ ,  $A_{1,i} \cap A_{2,i} \cap \dots \cap A_{k,i} = \emptyset$ . It thus suffices to show these are the only empty  $k$ -wise intersections. To this end, for  $\mathbf{i} = (i_1, \dots, i_{k-1})$ , define

$$A(\mathbf{i}) := A_{1,i_1} \cap \dots \cap A_{k-1,i_{k-1}}.$$

**Lemma 7.** *Let  $\mathbf{i} = (i_1, \dots, i_{k-1})$ . If  $A(\mathbf{i})^c = A_{k,i_k}$ , then  $i_1 = \dots = i_k$ .*

*Proof.* We proceed by induction on  $k$  where the result is trivial when  $k = 2$ . In the case where  $k > 2$ ,  $i_{k-1} - k + 2 = i_k + x$  for some  $x$  such that  $-(k - 2) \leq x \leq (k - 2)$  and thus  $i_{k-1} + (k - 2) = i_{k-1} - (k - 2) + (2k - 4) = i_k + x + (2k - 4)$ .

Next, there is a  $y$  such that  $-(k - 2) \leq y \leq (k - 2)$  with  $i_{k-1} + (k - 2) = i_k + y$ , and since  $n \geq 4k$ ,  $x + 2k - 4 = y$  with equality over  $\mathbb{Z}$  and moreover  $i_{k-1} + (k - 2) = i_k + (k - 2)$  over  $\mathbb{Z}$  and hence  $i_k = i_{k-1}$ . Removing these elements from each set, the result then follows by induction.  $\square$

If  $A_{1,i_1} \cap \dots \cap A_{k,i_k} = \emptyset$ , then as  $A(\mathbf{i}) = A_{1,i_1} \cap A_{2,i_2} \cap \dots \cap A_{k-1,i_{k-1}}$ ,

$$\emptyset = A_{1,i_1} \cap A_{2,i_2} \cap \dots \cap A_{k-1,i_{k-1}} \cap A_{k,i_k} = A(\mathbf{i}) \cap A_{k,i_k}.$$

The result follows by noting  $|A(\mathbf{i})| \geq n - (2k - 3)$ ,  $|A_{k,i_k}| = 2k - 3$ , and using Lemma 7.

## 2.2 An Explicit Construction

Let  $k \geq 3$ . An explicit construction of a Bollobás  $(k, 2)$ -tuple  $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$  where  $|\mathcal{A}_i| = 2^n$  and each  $\mathcal{A}_i$  consists of subsets of  $X$  for  $|X| = kn$  may be described as follows. Let  $I_j := \{x_{j,1}, x_{j,2}, \dots, x_{j,k}\}$  and consider  $X = I_1 \sqcup \dots \sqcup I_n$ . Now, for each  $f : [n] \rightarrow [2]$  and  $j \in [k]$ , define

$$A_{j,f} := \{x_{1,f(1)+j-1}, \dots, x_{n,f(n)+j-1}\}^c$$

where we work modulo  $k$  within the subscripts of  $I_j$ . It is straightforward to check that  $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$  is a Bollobás  $(k, 2)$ -tuple. This establishes the lower bound on  $\beta_{k,2}(n)$  in Equation (3) and hence the upper bound on  $f_{k,2}(n)$  in Equation (8).

### 3 Proof of Theorem 4

#### 3.1 Upper bound on $f_{k,t}(n)$

We wish to find a covering of  $H_{k,t}(n)$  with complete  $k$ -partite  $k$ -graphs and assume the parts of  $H_{k,t}(n)$  are  $X_1, X_2, \dots, X_k$ . For each subset  $T$  of  $[k]$  of size  $t$ , consider the uniformly random coloring  $\chi_T : [n] \rightarrow T$ . Given such a  $\chi_T$ , let  $Y_i \subset X_i$  be the vertices of color  $i$  for  $i \in T$ ; that is  $Y_i := \{x_{ij} : \chi(j) = i\}$  and  $Y_i = X_i$  for  $i \notin T$ . Denote by  $H(T, \chi)$  the (random) complete  $k$ -partite hypergraph with parts  $Y_1, Y_2, \dots, Y_k$ , and note that  $H(T, \chi) \subset H_{k,t}(n)$ . We place each  $H(T, \chi)$  a total of  $N$  times independently and randomly where

$$N = \left\lfloor \frac{(t+1)t^t \log_2 n}{(k-t+1) \log_2 e} \right\rfloor$$

and produce  $\binom{k}{t}N$  random subgraphs  $H(T, \chi)$ . For a set partition  $\pi$  of  $[k]$ , let  $|\pi|$  denote the number of parts in the partition and index the parts by  $[\pi]$ . Given a set partition  $\pi = (P_1, P_2, \dots, P_s)$ , let

$$f(\pi, t) = \sum_{T \in \binom{[s]}{t}} \prod_{i \in T} |P_i|.$$

If  $U$  is the number of edges of  $H_{k,t}(n)$  not in any of these subgraphs, then

$$\mathbb{E}(U) \leq \sum_{|\pi| \geq t} n^{|\pi|} (1-t^{-t})^{Nf(\pi,t)} = \sum_{t \leq s \leq k} n^s \sum_{|\pi|=s} (1-t^{-t})^{Nf(\pi,t)}. \quad (13)$$

For sufficiently large  $n$ , we claim that  $\mathbb{E}(U) < 1$ , which implies there exists a covering of  $H_{k,t}(n)$  with at most  $\binom{k}{t}N$  complete  $k$ -partite  $k$ -graphs, as required. The following technical lemma states that  $f$  is a decreasing function in the set partition lattice, and that  $f(\pi, t)$  increases when we merge all but one element of a smaller part of  $\pi$  with a larger part of  $\pi$ :

**Lemma 8.** *Let  $k \geq s \geq t \geq 2$ , and let  $\pi = (P_1, P_2, \dots, P_s)$  be a partition of  $[k]$ .*

- (i) *If  $\pi'$  is a refinement of  $\pi$  with  $|\pi'| = s+1$ , then  $f(\pi, t) \leq f(\pi', t)$ .*
- (ii) *If  $|P_1| \geq |P_2| \geq 2$  and  $a \in P_2$ , and  $\pi'$  is the partition  $(P'_1, P'_2, \dots, P'_s)$  of  $[k]$  with  $P'_1 = P_1 \cup P_2 \setminus \{a\}$  and  $P'_2 = \{a\}$  and with  $P'_i = P_i$  for  $3 \leq i \leq s$ , then  $f(\pi', t) \leq f(\pi, t)$ .*

The proof of Lemma 8 part (i) is in Appendix A and the proof of (ii) is similar to the proof of (i). By Lemma 8, a set partition of  $[k]$  into  $s$  parts which minimizes  $f(\pi, t)$  consists of one part of size  $k-s+1$  and  $s-1$  singleton parts and hence

$$\min\{f(\pi, t) : |\pi| = s\} = (k-s+1) \binom{s-1}{t-1} + \binom{s-1}{t}. \quad (14)$$

In what follows, we denote a set partition of  $[k]$  into  $s$  parts which minimizes  $f(\pi, t)$  by  $\pi_s$ .



For  $n$  large enough, and all  $s$  where  $t \leq s \leq k$ , we will show

$$\frac{\sum_{|\pi|=t}(1-t^{-t})^{Nf(\pi,t)}}{\sum_{|\pi|=s}(1-t^{-t})^{Nf(\pi,t)}} \geq n^{s-t}.$$

Replacing the numerator with its largest term and each term in denominator with its largest term,

$$\frac{\sum_{|\pi|=t}(1-t^{-t})^{Nf(\pi,t)}}{\sum_{|\pi|=s}(1-t^{-t})^{Nf(\pi,t)}} \geq \frac{(1-t^{-t})^{Nf(\pi_t,t)}}{S(k,s)(1-t^{-t})^{Nf(\pi_s,t)}} = \frac{1}{S(k,s)}(1-t^{-t})^{N(f(\pi_t,t)-f(\pi_s,t))}$$

where  $S(k,s)$  is the Stirling number of the second kind. Taking  $n \geq S(k,s)$ , we will show in Appendix B that

$$\frac{1}{S(k,s)}(1-t^{-t})^{N(f(\pi_t,t)-f(\pi_s,t))} \geq n^{s-t}. \quad (15)$$

Therefore, the index  $s = t$  maximizes the right hand side of Equation (13), and hence

$$\mathbb{E}[U] \leq (k-t+1)(n^t) \sum_{|\pi|=t} (1-t^{-t})^{Nf(\pi,t)} < (k-t+1)n^t S(k,t)(1-t^{-t})^{N(k-t+1)} < 1$$

for our choice of  $N$  provided  $n \geq kS(k,t)$ . Thus,

$$f_{k,t}(n) \leq \binom{k}{t} \frac{(t+1)t^t \log_2 n}{(k-t+1) \log_2 e} = \frac{(t+1)t^{t-1}}{\log_2 e} \binom{k}{t-1} \log_2 n.$$

### 3.2 Lower bound on $f_{k,2}(n)$

In this section, we show

$$f_{k,2}(n) \geq \min\{m : \binom{m}{\lceil m/k \rceil} \geq n\}. \quad (16)$$

Let  $\{H_1, H_2, \dots, H_m\}$  be a covering of  $H_{k,2}(n)$  with  $m = f_{k,2}(n)$  complete  $k$ -partite  $k$ -graphs. We recall  $H_{k,2}(n) = K_{n,n,\dots,n} \setminus M$ , where  $M$  is a perfect matching of  $K_{n,n,\dots,n}$ . For  $i \in [k]$  and  $j \in [n]$ , define  $A_{i,j} = \{H_r : x_{ij} \in V(H_r)\}$  and  $\mathcal{A}_i = \{A_{i,j} : 1 \leq j \leq n\}$ . As in (6),  $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$  is a Bollobás  $(k, 2)$ -tuple of size  $n$ . For convenience, for each  $i \in [k]$ , let  $\phi_i : [k] \rightarrow [2]$  be so that  $\phi_i^{-1}(1) = \{i\}$ . Taking the sum of inequality from Theorem 2 with  $t = 2$  over all  $i \in [k]$ ,

$$\sum_{i=1}^k \sum_{j=1}^n \left( \frac{|A_{1,j}(\phi_i) \cup A_{2,j}(\phi_i)|}{|A_{1,j}(\phi_i)|} \right)^{-1} \leq k. \quad (17)$$

We use this inequality to give a lower bound on  $f_{k,2}(n) = m$ . First we observe

$$\sum_{r=1}^m |V(H_r)| = \sum_{j=1}^n \sum_{i=1}^k |A_{i,j}| = \sum_{j=1}^n \sum_{i=1}^k |A_{1,j}(\phi_i)|. \quad (18)$$

Let  $\partial H$  denote the set of  $(k-1)$ -tuples of vertices contained in some edge of a hypergraph  $H$ . Then

$$\sum_{r=1}^m |\partial H_r \cap \partial M| = \sum_{j=1}^n \sum_{i=1}^k |A_{2,j}(\phi_i)|. \quad (19)$$

Putting the above identities together,

$$\sum_{r=1}^m |V(H_r)| + \sum_{r=1}^m |\partial H_r \cap \partial M| = \sum_{j=1}^n \sum_{i=1}^k (|A_{1,j}(\phi_i)| + |A_{2,j}(\phi_i)|). \quad (20)$$

We note  $|\partial H_r \cap \partial M| \leq |V(H_r)|/(k-1)$ , and therefore

$$\sum_{r=1}^m |\partial H_r \cap \partial M| \leq \frac{1}{k-1} \sum_{r=1}^m |V(H_r)|. \quad (21)$$

It follows that

$$\sum_{j=1}^n \sum_{i=1}^k (|A_{1,j}(\phi_i)| + |A_{2,j}(\phi_i)|) \leq \frac{k}{k-1} \sum_{r=1}^m |V(H_r)|. \quad (22)$$

Subject to the linear inequalities (18) and (22), the left side of (17) is minimized when  $kn|A_{1,j}(\phi_i)| = \sum_{r=1}^m |V(H_r)|$  and  $kn(|A_{1,j}(\phi_i)| + |A_{2,j}(\phi_i)|) = (k-1)|A_{1,j}(\phi_i)|$ . Since  $|V(H_r)| \leq (k-1)n$  for all  $r \in [m]$ , (17) implies  $\binom{m}{\lceil m/k \rceil} \geq n$ , which gives (16).  $\square$

### 3.3 Lower bound on $f_{k,k}(n)$

Let  $\mathcal{H} = \{H_1, H_2, \dots, H_m\}$  be a minimal covering of  $H_{k,k}(n)$  with complete  $k$ -partite  $k$ -graphs, so  $m = f(H_{k,k}(n))$ . Given a  $k$ -partite  $k$ -graph  $H$ , consider its 2-shadow  $\partial_2(H) = \{R \subset V(H) : |R| = k-2, R \subset e \text{ for some } e \in H\}$ . Let  $\partial_2(\mathcal{H}) = \bigcup_{i=1}^m \partial_2(H_i)$ .

Given  $R \in \partial_2(\mathcal{H})$  and  $H_i \in \mathcal{H}$ , let  $H_i(R) := \{e \in \binom{V(H_i)}{2} : e \cup R \in H_i\}$  be the possibly empty link graph of the edge  $R$  in the hypergraph  $H_i$  and let  $V(H_i(R))$  be the set of vertices in the link graph. Observe that double counting yields

$$\sum_{R \in \partial_2(\mathcal{H})} \left( \sum_{i=1}^m |V(H_i(R))| \right) = \sum_{i=1}^m \left( \sum_{R \in \partial_2(H_i)} |V(H_i(R))| \right). \quad (23)$$

An optimization argument yields  $|\partial_2(H_i)|$  is maximized when the parts of  $H_i$  are of equal or nearly equal maximal size. Since  $|V(H_i(R))| \leq 2(n-k+2)$ , the right hand side of Equation (23) is bounded above by

$$\sum_{i=1}^m \left( \sum_{R \in \partial_2(H_i)} |V(H_i(R))| \right) \leq m \cdot \binom{k}{2} \cdot \left( \frac{n}{k} \right)^{k-2} \cdot 2(n-k+2). \quad (24)$$

For a lower bound on the left hand side of Equation (23), fix  $R \in \partial_2(\mathcal{H})$  and without loss of generality suppose that  $R = \{x_{1,1}, \dots, x_{k-2,k-2}\}$ . Let  $Y = [k-1, n]$ . Let  $K_{Y,Y}$  be the complete bipartite graph with two distinct copies of  $Y$  and  $\mathcal{M} = \{(x_{k-1,i}, x_{k,i} : i \in Y)\}$  be a perfect matching in  $K_{Y,Y}$ . Then,  $\{H_1(R), \dots, H_m(R)\}$  forms a biclique cover of  $K_{Y,Y} \setminus \mathcal{M}$ . Applying the convexity result of Tarjan [23, Lemma 5],

$$\sum_{i=1}^m |V(H_i(R))| \geq (n - k + 2) \log_2(n - k + 2).$$

Noting that  $|\partial_2(\mathcal{H})| = \binom{k}{2} \binom{n}{k-2}$ , the left hand side of Equation (23) is bounded below by

$$\sum_{R \in \partial_2(\mathcal{H})} \left( \sum_{i=1}^m |V(H_i(R))| \right) \geq \binom{k}{2} \binom{n}{k-2} (n - k + 2) \log_2(n - k + 2). \quad (25)$$

Comparing the bounds from Equation (24) and Equation (25),

$$m \geq \frac{\binom{n}{k-2} \log_2(n - k + 2)}{2 \left(\frac{n}{k}\right)^{k-2}} \geq \frac{k^{k-2}}{2} \log_2 n$$

provided that  $n$  is large enough.

For  $t \geq 3$  and  $t < k$ , the lower bound on  $f_{k,t}(n)$  in Theorem 4 is obtained from the lower bounds on  $f_{t-1,t-1}(n-1)$  as follows: Let  $\mathcal{H} = \{H_1, H_2, \dots, H_m\}$  be a minimal covering of  $H_{k,t}(n)$  with complete  $k$ -partite  $k$ -graphs, so  $m = f(H_{k,t}(n))$ . Given  $T \in \binom{[k]}{k-t+1}$ , define  $H_T \subset H_{k,t}(n)$  by

$$H_T := \{ \{x_{1,i_1}, \dots, x_{k,i_k}\} \in H_{k,t}(n) : i_j = 1 \forall j \in T \}.$$

It follows that at least  $f_{t-1,t-1}(n-1)$  of the complete  $k$ -partite  $k$ -graphs in  $\mathcal{H}$  are needed to cover  $H_T$ . Moreover, for distinct  $T, T' \in \binom{[k]}{k-t+1}$ , the corresponding complete  $k$ -partite  $k$ -graphs from  $\mathcal{H}$  are necessarily pairwise disjoint and hence

$$f_{k,t}(n) \geq \binom{k}{k-t+1} f_{t-1,t-1}(n-1) \geq \binom{k}{t-1} \frac{(t-1)^{t-3}}{2} \log_2 n$$

provided that  $n$  is large enough.

## 4 Concluding remarks

- Our main theorem, Theorem 2 is tight for  $t = 2$  and  $k \geq 2$ , as shown in Section 2.1. It would be interesting to generalize this example to  $2 < t \leq k$  to determine whether Theorem 2 is tight in general. The first open case is  $t = k = 3$ .
- A particular case of the Bollobás set pairs inequality occurs when every set in  $\mathcal{A}$  has size  $a$  and every set in  $\mathcal{B}$  has size  $b$ , and one obtains the tight bound  $|\mathcal{A}| \leq \binom{a+b}{b}$ . The

generalization to Bollobás  $(k, t)$ -tuples for  $k \geq 3$  is equally interesting but wide open, as are potential generalizations to vector spaces – see Lovász [17, 18].

• Orlin [20] proved that the clique cover number  $cc(K_n \setminus M)$  of a complete graph  $K_n$  minus a perfect matching  $M$  is precisely  $\min\{m : 2^{\binom{m-1}{\lfloor m/2 \rfloor}} \geq n\}$ . Theorem 4 yields lower bounds on the clique cover number of the complement of a perfect matching  $M$  in the complete  $k$ -uniform hypergraph  $K_n^k$ :

**Corollary 9.** *Let  $K_n^k \setminus M$  be the complement of a perfect matching in  $K_n^k$ . Then*

$$cc(K_n^k \setminus M) \geq \frac{\log_2 \frac{n}{k}}{H(\frac{1}{k})} \geq \frac{k \log_2 \frac{n}{k}}{\log_2(ke)}.$$

• It would be interesting to prove an analog of Equation (16) for  $t \geq 3$ . That is,

$$f_{k,t}(n) \geq \min\{m : \binom{m}{\alpha_1, \dots, \alpha_t} \geq n_{(t-1)}\}$$

for some optimal  $\alpha_1, \dots, \alpha_t$ . The difficulty here lies in determining effective bounds on  $|A_{i,\sigma}(\phi)|$ .

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## A Proof of Lemma 8(i)

Let  $k \geq s \geq t \geq 2$ , and let  $\pi = (P_1, P_2, \dots, P_s)$  be a partition of  $[k]$ . In this section, we will show that if  $\pi'$  is a refinement of  $\pi$  with  $|\pi'| = s + 1$ , then  $f(\pi, t) \leq f(\pi', t)$ .

*Proof.* Let  $\pi = P_1|P_2|\dots|P_s$  and without loss of generality,  $\pi' = P_x|P_y|P_2|\dots|P_s$ . Setting  $\mathcal{T}(\bar{1}) = \{T \in [s]^{(t)} : 1 \notin T\}$  and  $\mathcal{T}'(\bar{x}, \bar{y}) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x, y \notin T\}$ , it follows that

$$\sum_{T \in \mathcal{T}(\bar{1})} \prod_{i \in T} |P_i| = \sum_{T \in \mathcal{T}'(\bar{x}, \bar{y})} \prod_{i \in T} |P_i|.$$

Now, letting  $\mathcal{T}(1) = \{T \in [s]^{(t)} : 1 \in T\}$  and  $\mathcal{T}'(x, \bar{y}) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x \in T, y \notin T\}$  and  $\mathcal{T}'(\bar{x}, y) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x \notin T, y \in T\}$ , we see that

$$\sum_{T \in \mathcal{T}(1)} \prod_{i \in T} |P_i| = \sum_{T \in \mathcal{T}'(\bar{x}, y)} \prod_{i \in T} |P_i| + \sum_{T \in \mathcal{T}'(x, \bar{y})} \prod_{i \in T} |P_i|$$

since  $|P_1| = |P_x| + |P_y|$ . Thus letting  $\mathcal{T}'(x, y) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x \in T, y \in T\}$ ,

$$f(\pi', t) - f(\pi, t) = \sum_{T \in \mathcal{T}'(x, y)} \prod_{i \in T} |P'_i|$$

and in particular  $f(\pi, t) \leq f(\pi', t)$ . □

## B Proof of Equation (15)

Let  $S(k, s)$  be the Stirling number of the second kind and  $f(\pi)$  be as in Section 3. In this section we will show

$$\frac{1}{S(k, s)} (1 - t^{-t})^{N(f(\pi_t, t) - f(\pi_s, t))} \geq n^{s-t}.$$

*Proof.* First, we recall that

$$N = \left\lfloor \frac{(t+1)t^t \log_2 n}{(k-t+1) \log_2 e} \right\rfloor \quad \text{and} \quad f(\pi_s, t) = (k-s+1) \binom{s-1}{t-1} + \binom{s-1}{t}.$$

As a result, when  $t \leq s < k$ , a calculation yields that

$$f(\pi_{s+1}, t) - f(\pi_s, t) = (k-s) \binom{s-1}{t-2}. \tag{26}$$

Letting  $n \geq S(k, t)$ , after taking  $\log_2(\cdot)$  on both sides of (15), it suffices to prove that

$$N \cdot \frac{f(\pi_s, t) - f(\pi_t, t)}{t^t} \left( -t^t \log_2(1 - t^{-t}) \right) \geq (s-t+1) \log_2(n). \tag{27}$$

Using the fact that  $(1 - t^{-t})^{t^t} \leq e^{-1}$  and our choice of  $N$ , it suffices to show that

$$f(\pi_s, t) - f(\pi_t, t) \geq \frac{(s-t+1)(k-t+1)}{t+1}. \tag{28}$$

The inequality in (28) holds for all  $k \geq s > t \geq 3$  by using (26). □