Comment by the authors, May 23, 1995

We thank J. B. Shearer for pointing out that Theorem 3 of our paper is not new (see [1], [2], [3]), and that the argument given in our paper actually proves the result with c=1/64, which is weaker than the best known [3] value c = 1/7.

## References

- B. Bollobás and P. Erdős, Alternating Hamiltonian Cycles, Israel Journal of Mathematics, 23 (1976), pp. 126-131, (c=1/69).
- [2] C. C. Chen and D. E. Daykin, Graphs with Hamiltonian Cycles Having Adjacent Lines Different Colors, Journal of Combinatorial Theory B, 21 (1976), pp. 135-139, (c=1/17).
- [3] J. B. Shearer, A Property of the Complete Colored Graph, Discrete Mathematics, 25 (1979), pp. 175-178, (c=1/7).

Comment by Rachel Rue Dept. of Mathematics, Carnegie Mellon University (rue@cs.cmu.edu) September 19, 1995.

There is a mistake in Lemma 2 of the paper, which requires halving the constants in Theorems 1,2,and 4. The constant in Theorems 1 and 4 should be changed from c < 1/32 to c < 1/64, and the constant in Theorem 2 should be changed from c < 1/64 to c < 1/128.

As stated in the paper, Lemma 2 reads as follows:

Let e, f be edges of  $K_n$  and  $X \subseteq E(K_n)$  be such that no edge in X shares an endpoint with either e or f. Then we can find, for each Hamilton cycle C containing both e and f and no edges of X, a set S(C) of (n-6)(n-9)/2 Hamilton cycles containing neither e, f or any edge in X, in such a way that if  $C \neq C'$  then  $S(C) \cap S(C') = \emptyset$ .

We show with a counterexample that S(C) can be no greater than  $n^2/4(n-3)(n-4)$ , and then give an algorithm which associates (n-6)(n-7)/4 distinct cycles without e, f, or any edges of X with each distinct cycle containing e, f, and no edges of X.

## Counterexample:

We pick a set X such that the ratio of the number of Hamilton cycles not containing e, f, or any edges of X to the number of Hamilton cycles containing e and f and no edges of X is less than  $n^2: 4(n-3)(n-4)$ .

Let  $e_0, e_1$ , and  $f_0, f_1$  be the endpoints of e and f, respectively. Let S be a cycle on the n-4 nodes of  $K_n \setminus \{e_0, e_1, f_0, f_1\}$ . Let X contain all edges in  $K_n$  except those in S and those which share endpoints with e or f. Then the number of Hamilton cycles containing both e and f and no edges in X is 4(n-4)(n-3). Proof: In any such cycle, nodes other than endpoints of e and f may be adjacent only to their neighbors in S or to the endpoints of e and f. So all such cycles may be constructed by inserting e and f into S as follows.

First pick an edge (a,b) in S, and substitute one of the node sequences  $\langle a, e_0, e_1, b \rangle$ ,  $\langle a, e_1, e_0, b \rangle$  for the sequence  $\langle a, b \rangle$ . Call the resulting cycle S'. Then pick an edge (c, d) in  $S' \setminus \{e\}$ , and insert f by substituting one of the node sequences  $\langle c, f_o, f_1, d \rangle$ ,  $\langle c, f_1, f_0, d \rangle$  for  $\langle c, d \rangle$ . There are (n-4)places to insert e in S, 2 ways to orient e in that position, (n-3) places to insert f in  $S' \setminus \{e\}$ , and 2 ways to orient f, which makes 4(n-3)(n-4) cycles in all.

The number of cycles not containing e, f, or edges in f is less than  $n^4$ . Proof: The number of such cycles is just the number of ways to insert the four nodes  $e_0, e_1, f_0, f_1$  into S without making  $e_0$  and  $e_1$  or  $f_0$  and  $f_1$  adjacent. Let  $D = \{e_0, e_1, f_0, f_1\}$ . There are (n-4)(n-5)(n-6)(n-7) Hamilton cycles with none of the nodes in D adjacent to each other; if we add in the cycles with one or more pair of nodes in D adjacent, the number added is  $O(n^3)$ ; the total number will still turn out to be less than  $n^4$ . Thus there are fewer than  $n^4/4(n-3)(n-4)$ Hamilton cycles not containing e, f, or edges in X for each Hamilton cycle containing e, f, and no edges of X.

**Construction:** Let C be a Hamilton cycle containing eand f, and no edges in X. We associate with C a set T(C)of Hamilton cycles which don't contain e, f, or any edges of X, constructed as follows. The new cycles are created by taking one endpoint each of e and f, and moving them to new positions, leaving the order of all other nodes fixed.

- 1. Orient C so that the lower numbered of the two nodes adjacent to  $e_0$  and  $e_1$  follows e. Suppose without loss of generality that  $e_1$  then follows  $e_0$ .
- 2. Pick any two edges of C other than the 6 edges incident with  $e_0, e_1, f_0$ , and  $f_1$ . There are  $\binom{(n-6)}{2}$  ways to do this. Of the two chosen edges, let (a, b) be the first edge following  $e_0$  in the given orientation of C; let (c, d) be the other edge. (If  $e_1$  preceded  $e_0$  in C, we would let (a, b) be the first edge preceding  $e_0$ )
- 3. Remove  $e_1$  from its original position and insert it between a and b.
- 4. Choose one of  $f_0$  and  $f_1$  to move between c and d, choosing so as to preserve the order in which  $e_0$ ,  $f_0$ , and  $f_1$  appear in the cycle.

Call the new cycle C'. Notice that it is possible to reconstruct the original positions of both  $e_0$  and  $e_1$  by looking at C':  $e_0$  hasn't moved, and the original position of  $e_1$  is given

by sliding it back around the cycle toward  $e_0$ , in whichever direction passes neither  $f_0$  nor  $f_1$ . It is not possible, however, to reconstruct the original positions of  $f_0$  and  $f_1$ ; their relative positions are correct, but it is impossible to tell whether  $f_0$  or  $f_1$  has been moved. Thus each of the (n-6)(n-7)/2 cycles in T(C) is associated with two cycles: our original cycle C, and a cycle identical to C except that the edge f is positioned between nodes c and d (in the same orientation as in C). Thus for each distinct cycle C containing e, f, and no edges of X, there are (n-6)(n-7)/4 distinct cycles not containing e, f, for any edges of X.