

Graphs with Girth $2\ell + 1$ and Without Longer Odd Holes that Contain an Odd K_4 -Subdivision

Rong Chen^a Yidong Zhou^a

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Abstract

We say that a graph G has an *odd K_4 -subdivision* if some subgraph of G is isomorphic to a K_4 -subdivision which if embedded in the plane the boundary of each of its faces has odd length and is an induced cycle of G . For a number $\ell \geq 2$, let \mathcal{G}_ℓ denote the family of graphs which have girth $2\ell + 1$ and have no odd hole with length greater than $2\ell + 1$. Wu, Xu and Xu conjectured that every graph in $\bigcup_{\ell \geq 2} \mathcal{G}_\ell$ is 3-colorable. Recently, Chudnovsky et al. and Wu et al., respectively, proved that every graph in \mathcal{G}_2 and \mathcal{G}_3 is 3-colorable. In this paper, we prove that no 4-vertex-critical graph in $\bigcup_{\ell \geq 5} \mathcal{G}_\ell$ has an odd K_4 -subdivision. Using this result, Chen proved that all graphs in $\bigcup_{\ell \geq 5} \mathcal{G}_\ell$ are 3-colorable.

Keywords: chromatic number, odd holes

Mathematics Subject Classifications: 05C15, 05C17, 05C69

1 Introduction

All graphs considered in this paper are finite, simple, and undirected. A *proper coloring* of a graph G is an assignment of colors to the vertices of G such that no two adjacent vertices receive the same color. A graph is *k -colorable* if it has a proper coloring using at most k colors. The *chromatic number* of G , denoted by $\chi(G)$, is the minimum number k such that G is k -colorable.

The *girth* of a graph G , denoted by $g(G)$, is the minimum length of cycles in G . A *hole* in a graph is an induced cycle of length at least four. An *odd hole* means a hole of odd length. For any integer $\ell \geq 2$, let \mathcal{G}_ℓ be the family of graphs that have girth $2\ell + 1$ and have no odd holes of length at least $2\ell + 3$. Robertson conjectured in [4] that the Petersen graph is the only graph in \mathcal{G}_2 that is 3-connected and internally 4-connected. Plummer and Zha [5] disproved Robertson's conjecture and proposed the conjecture that all 3-connected and internally 4-connected graphs in \mathcal{G}_2 have bounded chromatic numbers, and proposed

^aCenter for Discrete Mathematics, Fuzhou University, Fuzhou, P. R. China (rongchen@fzu.edu.cn, zoed98@126.com).

the strong conjecture that such graphs are 3-colorable. The first was proved by Xu, Yu, and Zha [9], who proved that all graphs in \mathcal{G}_2 are 4-colorable. The strong conjecture proposed by Plummer and Zha in [5] was solved by Chudnovsky and Seymour [2]. Wu, Xu, and Xu [7] showed that graphs in $\bigcup_{\ell \geq 2} \mathcal{G}_\ell$ are 4-colorable and conjectured

Conjecture 1. ([7], Conjecture 6.1.) For each integer $\ell \geq 2$, every graph in \mathcal{G}_ℓ is 3-colorable.

Wu, Xu and Xu [8] recently proved that Conjecture 1 holds for $\ell = 3$.

We say that a graph G has an *odd K_4 -subdivision* if some subgraph of G is isomorphic to a K_4 -subdivision which if embedded in the plane the boundary of each of its faces has odd length and is an induced cycle of G . Note that an odd K_4 -subdivision of G maybe not induced. However, when $G \in \mathcal{G}_\ell$ for each integer $\ell \geq 2$, all odd K_4 -subdivisions of G are induced by Lemma 8 (2). In this paper, we prove the following theorem.

Theorem 2. No 4-vertex-critical graph in $\bigcup_{\ell \geq 5} \mathcal{G}_\ell$ has an odd K_4 -subdivision.

Using Theorem 2, Chen [1] proved that Conjecture 1 holds for all $\ell \geq 5$. Recently, following idea in this paper and [1], Wang and Wu [6] further proved that Conjecture 1 holds for $\ell = 4$.

2 Preliminaries

A *cycle* is a connected 2-regular graph. Let G be a graph. A vertex $v \in V(G)$ is called a *degree- k vertex* if it has exactly k neighbours. For any $U \subseteq V(G)$, let $G[U]$ be the subgraph of G induced on U . For subgraphs H and H' of G , set $|H| := |E(H)|$ and $H\Delta H' := E(H)\Delta E(H')$. Let $H \cup H'$ denote the subgraph of G whose vertex set is $V(H) \cup V(H')$ and edge set is $E(H) \cup E(H')$. Let $H \cap H'$ denote the subgraph of G with edge set $E(H) \cap E(H')$ and without isolated vertex. Let $N(H)$ be the set of vertices in $V(G) - V(H)$ that have a neighbour in H . Set $N[H] := N(H) \cup V(H)$.

Let P be an (x, y) -path and Q be a (y, z) -path. When P and Q are internally disjoint, let PQ denote the (x, z) -path $P \cup Q$. Evidently, PQ is a path when $x \neq z$, and PQ is a cycle when $x = z$. Let P^* denote the set of internal vertices of P . When $u, v \in V(P)$, let $P(u, v)$ denote the subpath of P with ends u, v . For simplicity, we will let $P^*(u, v)$ denote $(P(u, v))^*$.

A graph is *k -vertex-critical* if $\chi(G) = k$ but $\chi(G \setminus v) < k$ for each $v \in V(G)$. Dirac in [3] proved that every k -vertex-critical graph is $(k - 1)$ -edge-connected. Hence, we have

Lemma 3. For each integer $k \geq 4$, each k -vertex-critical graph G has no 2-edge-cut.

A *theta graph* is a graph that consists of a pair of distinct vertices joined by three internally disjoint paths. Let C be a hole of a graph G . A path P of G is a *chordal path* of C if $V(P^*) \cap V(C) = \emptyset$ and $C \cup P$ is an induced theta-subgraph of G . Lemma 4 will be frequently used.

Lemma 4. *Let $\ell \geq 2$ be an integer and C be an odd hole of a graph $G \in \mathcal{G}_\ell$. Let P be a chordal path of C , and P_1, P_2 be the internally disjoint paths of C that have the same ends as P . Assume that $|P|$ and $|P_1|$ have the same parity. If $|P_1| \neq 1$, then $|P_1| > |P_2|$ and all chordal paths of C with the same ends as P_1 have length $|P_1|$.*

Proof. Since $|C| = 2\ell + 1$, $|P_1| \neq 1$ and $|P|$ and $|P_1|$ have the same parity, $P \cup P_2$ is an odd hole. Moreover, since $g(G) = 2\ell + 1$ and all odd holes in G have length $2\ell + 1$, we have $\ell + 1 \leq |P_1| = |P|$ and $|P_2| \leq \ell$, so $|P_1| > |P_2|$ and all chordal paths of C with the same ends as P_1 have length $|P_1|$. \square

Let P be a path with i vertices. If $G - V(P)$ is disconnected, then we say that P is a P_i -cut. Usually, a P_2 -cut is also called a K_2 -cut. Evidently, every k -vertex-critical graph has no K_2 -cut. Chudnovsky and Seymour in [2] proved that every 4-vertex-critical graph G in \mathcal{G}_2 has no P_3 -cut. Using the same argument as [2], Wu et al. [8] extend this result to graphs in $\bigcup_{\ell \geq 2} \mathcal{G}_\ell$. Since the paper [8] does not include a proof of Lemma 5, we give a proof here for completeness.

Lemma 5. (*[8]*) *For any number $\ell \geq 2$, every 4-vertex-critical graph in \mathcal{G}_ℓ has neither a K_2 -cut nor a P_3 -cut.*

Proof. It is well-known that every k -vertex-critical graph has no clique as a cut. Hence, it suffice to show that every 4-vertex-critical graph in \mathcal{G}_ℓ has no P_3 -cut. Let $G \in \mathcal{G}_\ell$ be a 4-vertex-critical graph. Assume to the contrary that $P = v_1v_2v_3$ is a path such that $G \setminus \{v_1, v_2, v_3\}$ is disconnected. Since G has no K_3 as its cut, $v_1v_3 \notin E(G)$. Let A_1 be the a component of $G \setminus \{v_1, v_2, v_3\}$, and let A_2 be the union of all other components. Set $G_i := G[A_i \cup \{v_1, v_2, v_3\}]$ for $i = 1, 2$. Since G is 4-vertex-critical, both G_1 and G_2 are 3-colorable. Let $\phi_i : V(G_i) \rightarrow \{1, 2, 3\}$ be a 3-coloring for $i = 1, 2$. By symmetry we may assume that $\phi_i(v_1) = 1$ and $\phi_i(v_2) = 2$ for $i = 1, 2$. Thus $\phi_1(v_3), \phi_2(v_3) \in \{1, 3\}$. If $\phi_1(v_3) = \phi_2(v_3)$, then G is 3-colorable, which is a contradiction. Thus by symmetry we may assume that $\phi_1(v_3) = 1$ and $\phi_2(v_3) = 3$. Let H_1 be the subgraph of G_1 induced on the set of vertices $v \in V(G_1)$ with $\phi_1(v) \in \{1, 3\}$. If v_1, v_3 belong to different components of H_1 , then by exchanging colors in the component containing v_3 , we obtain another 3-coloring of G_1 that can be combined with ϕ_2 to show that G is 3-colorable. So v_1, v_3 belong to the same component of H_1 . Then there is an induced (v_1, v_3) -path P_1 in H_1 having even length as $\phi_1(v_1) = 1 = \phi_1(v_3)$. Similarly, there is an induced (v_1, v_3) -path P_2 in G_2 having odd length as $\phi_2(v_1) = 1$ and $\phi_2(v_3) = 3$. Moreover, since PP_1, PP_2 are cycles of G and $g(G) = 2\ell + 1$, we have $|P_1| \geq 2\ell - 1$ and $|P_2| \geq 2\ell$, so $P_1 \cup P_2$ is an odd hole of G of length at least $4\ell - 1$, which is a contradiction as $G \in \mathcal{G}_\ell$. \square

Lemma 6. *Let $\ell \geq 2$ be an integer and x, y be non-adjacent vertices of a graph $G \in \mathcal{G}_\ell$. Let P be an induced (x, y) -path of G . If $|P| \leq \ell$ and all induced (x, y) -paths have length $|P|$, then no block of G contains two non-adjacent vertices in $V(P)$. In particular, each vertex in P^* is a cut-vertex of G .*

Proof. Assume not. Then there is a block B of G containing two consecutive edges of P . Let Q be an induced path in B with ends in $V(P)$ and with $V(P) \cap V(Q^*) = \emptyset$.

Since every pair of edges in a 2-connected graph is contained in a cycle, such a Q exists. Without loss of generality we may further assume that Q is chosen with $|Q|$ as small as possible. Let C be the unique cycle in $P \cup Q$. Then $C \Delta P$ is an (x, y) -path. Since Q is induced, the ends of Q are not adjacent. Moreover, since Q is chosen with $|Q|$ as small as possible, $C \Delta P$ is an induced (x, y) -path, so $|C \Delta P| = |P| \leq \ell$ by the assumption of the lemma. Hence, $|C| \leq 2\ell$, contrary to the fact $g(G) = 2\ell + 1$. \square

Lemma 7. *Let $\ell \geq 4$ be an integer and x, y be non-adjacent vertices of a graph $G \in \mathcal{G}_\ell$. Let X be a vertex cut of G with $\{x, y\} \subseteq X \subseteq N[\{x, y\}]$, and G_1 be an induced subgraph of G whose vertex set consists of X and the vertex set of a component of $G - X$. If all induced (x, y) -paths in G_1 have length k with $4 \leq k \leq \ell$, then G has a degree-2 vertex, a K_1 -cut, or a K_2 -cut.*

Proof. Assume that G has no degree-2 vertices. Let P be an induced (x, y) -path in G_1 . Let uvw be a subpath of P^* . Such uvw exists as $k \geq 4$. By the definition of G_1 , we have $v \notin X$, so $N_G[v] = N_{G_1}[v]$, which implies $d_{G_1}(v) \geq 3$. By applying Lemma 6 to G_1 , there is a block B of G_1 such that either $V(B) \cap V(P) = \{v\}$, or B is not isomorphic to K_2 and $V(B) \cap V(P)$ is $\{u, v\}$ or $\{u, w\}$. When the first case happens, since $X \subseteq N[\{x, y\}]$, $x, y \notin B$ and $v \notin X$, we have $X \cap V(B) = \emptyset$, for otherwise $P(x, v)$ or $P(v, y)$ is contained in a cycle of $P \cup B$, so the vertex v is a cut-vertex of G as X is a vertex cut of G . When the latter case happens, by symmetry we may assume that $V(B) \cap V(P) = \{u, v\}$. Since B is a block of G_1 , $X \subseteq N[\{x, y\}]$ and uvw is a subpath of P^* , we have $V(B) \cap X = \{u\} \cap X$, so $\{u, v\}$ is a K_2 -cut of G_1 and G . \square

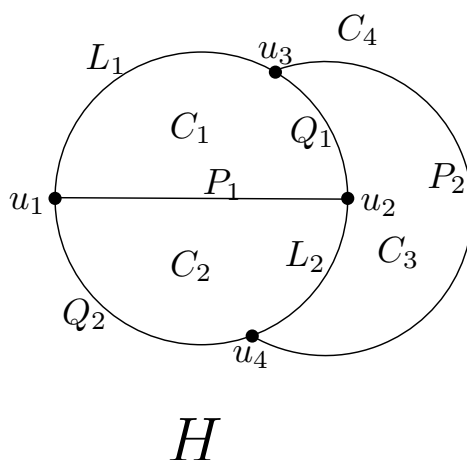


Figure 1: u_1, u_2, u_3, u_4 are the degree-3 vertices of H . All faces C_1, C_2, C_3, C_4 of H are odd holes. $\{P_1, P_2\}, \{Q_1, Q_2\}, \{L_1, L_2\}$ are the pairs of vertex disjoint arrises of H .

Let H be a graph that is isomorphic to a subdivision of K_4 , and let P be a path of H whose ends are degree-3 vertices in H . If P^* contains no degree-3 vertex of H , then we say that P is an *arris* of H . Evidently, there are exactly six arrises of H . See Figure 1.

Lemma 8. For any integer $\ell \geq 2$, if a graph $G \in \mathcal{G}_\ell$ has an odd K_4 -subdivision H , then the following statements hold.

- (1) Each pair of vertex disjoint arrises in H have the same length and their lengths are at most ℓ .
- (2) H is an induced subgraph of G .
- (3) When $\ell \geq 3$, every vertex in $V(G) - V(H)$ has at most one neighbour in $V(H)$.

Proof. Without loss of generality we may assume that H is pictured as the graph in Figure 1. First, we prove that (1) is true. Assume that $|P_1| > |P_2|$. Since C_1 and C_4 are odd holes, $|Q_1| < |Q_2|$. Hence, $|P_2 \cup Q_1 \cup L_2| < |P_1 \cup Q_2 \cup L_2|$, which is a contradiction to the fact that C_2 and C_3 are both odd holes. So $|P_1| = |P_2|$. By symmetry each pair of vertex-disjoint arrises have the same length. Moreover, since $C_1 \Delta C_2$ is an even cycle with length at least $2\ell + 2$, we have $|P_1| \leq \ell$. By symmetry we have $|Q_1|, |L_1| \leq \ell$. So (1) holds.

Secondly, we prove that (2) is true. Suppose not. Since odd holes have no chord, by symmetry we may assume that there is an edge st in G with $s \in V(P_1^*)$ and $t \in V(P_2^*)$. On one hand, since $P_1(u_1, s)stP_2(t, u_4)Q_2$ and $P_1(u_2, s)stP_2(t, u_3)Q_1$ are cycles, by (1) we have

$$|P_1| + |P_2| + |Q_1| + |Q_2| + 2 = 2(|P_1| + |Q_1| + 1) \geq 2(2\ell + 1).$$

On the other hand, since $|P_1|, |Q_1| \leq \ell$ by (1), we have $|P_1| = |Q_1| = \ell$, implying that $|L_1| = |L_2| = 1$. Moreover, by the symmetry between L_1, L_2 and Q_1, Q_2 , we have $|Q_1| = |Q_2| = 1$, which is a contradiction as $|Q_1| = \ell \geq 2$. So (2) holds.

Finally, we prove that (3) is true. Suppose to the contrary that some vertex $x \in V(G) - V(H)$ has at least two neighbours in $V(H)$. Since a vertex not in an odd hole can not have two neighbours in the odd hole, x has exactly two neighbours x_1, x_2 in $V(H)$. By symmetry we may further assume that $x_1 \in V(P_1^*)$ and $x_2 \in V(P_2^*)$. Since $C'_1 = P_1(u_1, x_1)x_1xx_2P_2(x_2, u_3)L_1$ and $C'_2 = P_1(u_1, x_1)x_1xx_2P_2(x_2, u_4)Q_2$ are cycles whose lengths have different parity,

$$|C'_1| + |C'_2| = 2\ell + 1 + 2(2 + |P_1(u_1, x_1)|) \geq 4\ell + 3.$$

Hence, $|P_1(u_1, x_1)| = \ell - 1$ and $x_1u_2 \in E(H)$ as $|P_1| \leq \ell$ by (1). This implies that $u_2x_1xx_2$ is a chordal path of C_3 with length 3, which is a contradiction to Lemma 4 as $\ell \geq 3$. \square

By Lemma 8 (1), all odd K_4 -subdivisions of a graph $G \in \mathcal{G}_\ell$ have exactly $4\ell + 2$ edges for each number $\ell \geq 2$.

3 Proof of Theorem 2

Let H_1, H_2 be vertex disjoint induced subgraphs of a graph G . An induced (v_1, v_2) -path P is a *direct connection* linking H_1 and H_2 if v_i is the only vertex in $V(P)$ having a neighbour in $V(H_i)$ for each $i \in \{1, 2\}$. Evidently, $V(P) \cap V(H_1 \cup H_2) = \emptyset$ and the set

of internal vertices of each shortest path joining H_1 and H_2 induces a direct connection linking H_1 and H_2 .

For convenience, Theorem 2 is restated here in another way.

Theorem 9. *Let $\ell \geq 5$ be an integer, and G be a graph in \mathcal{G}_ℓ . If G is 4-vertex-critical, then G has no odd K_4 -subdivisions.*

Proof. Suppose not. Let H be a subgraph of G that is isomorphic to an odd K_4 -subdivision and pictured as the graph in Figure 1. By Lemma 8 (2), H is an induced subgraph of G . By Lemma 8 (1), we have

$$|P_1| = |P_2| \leq \ell, \quad |Q_1| = |Q_2| \leq \ell, \quad \text{and} \quad |L_1| = |L_2| \leq \ell. \quad (3.1)$$

Without loss of generality we may assume that P_1, P_2 are longest arrises in H .

Let e, f be the edges of P_2 incident with u_3, u_4 , respectively. Since G is 4-vertex-critical, $\{e, f\}$ is not an edge-cut of G by Lemma 3, so there is a direct connection P in $G \setminus \{e, f\}$ linking P_2^* and $H - V(P_2^*)$. Let v_1, v_2 be the ends of P with v_2 having a neighbour in P_2^* and v_1 having a neighbour in $H - V(P_2^*)$. By Lemma 8 (3), both v_1 and v_2 have a unique neighbour in $V(H)$. Let x, y be the neighbours of v_1 and v_2 in $V(H)$, respectively. That is, $x \in V(H) - V(P_2^*)$ and $y \in V(P_2^*)$. Set $P' := xv_1Pv_2y$. Since H is an induced subgraph of G , so is $H \cup P'$.

9.1. $x \notin \{u_1, u_2\}$.

Subproof. Assume not. By symmetry we assume that $x = u_1$. Set $C'_4 = L_1P'P_2(y, u_3)$. Since C_4 is an odd hole, by symmetry we may assume that C'_4 is an even hole and $C_4\Delta C'_4$ is an odd hole. Since $P'P_2(y, u_3)$ is a chordal path of C_1 , by (3.1) and Lemma 4, we have $|L_1| = 1$. So $|P_1| = |Q_1| = \ell$ by (3.1) again. Since $P'P_2(y, u_4)$ is a chordal path of C_2 and $C_4\Delta C'_4$ is an odd hole, $|P'P_2(y, u_4)| = |P_1L_2| = \ell + 1$ by (3.1) and Lemma 4 again. Moreover, since $|P_2| = \ell$ and $|L_1| = 1$, we have $|C'_4| \leq 2\ell$, which is not possible. So $x \neq u_1$. \square

Set $d(H) := |P_1| - \min\{|Q_1|, |L_1|\}$. We say that $d(H)$ is the *difference* of H . Without loss of generality we may assume that among all odd K_4 -subdivisions, H is chosen with difference as small as possible.

9.2. $x \notin V(P_1)$.

Subproof. Suppose to the contrary that $x \in V(P_1)$. Then $x \in V(P_1^*)$ by 9.1. Without loss of generality we may assume that $|L_1| \geq |Q_1|$. Set $C'_2 = Q_2P_1(u_1, x)P'P_2(y, u_4)$. Since C_4 is an odd hole, either C'_2 or $C_4\Delta C'_2$ is an odd hole. Suppose that $C_4\Delta C'_2$ is an odd hole. Since $C_1 \cup C_3 \cup P'$ is an odd K_4 -subdivision, by Lemma 8 (1) and (3.1), we have $|P'| = |Q_1|$, $|P_1(u_1, x)| = |P_2(u_4, y)|$, and $|P_1(u_2, x)| = |P_2(u_3, y)|$. So C'_2 is an even hole of length $2(|Q_2| + |P_1(u_1, x)|)$ by (3.1) again, implying $|L_1| + |P_1(u_1, x)| \geq |Q_2| + |P_1(u_1, x)| \geq \ell + 1$ as $|L_1| \geq |Q_1|$. Then $|C_4\Delta C'_2\Delta C_1| = 2|P_1(x, u_2)Q_1| \leq 2\ell$, contrary to the fact $g(G) = 2\ell + 1$. So C'_2 is an odd hole.

Since $C_2 \cup C'_2 \cup C_3$ is an odd K_4 -subdivision, it follows from Lemma 8 (1) and (3.1) that

$$|P'| = |L_2|, |P_1(u_1, x)| = |P_2(u_3, y)|, \text{ and } |P_1(u_2, x)| = |P_2(u_4, y)|. \quad (3.2)$$

Then $|C_2 \Delta C'_2 \Delta C_1| = 2|L_1| + 2\ell + 1$. Since $C_2 \Delta C'_2 \Delta C_1$ is not an odd hole,

$$1 \in \{|Q_2|, |P_1(u_2, x)|, |P_2(u_3, y)|\}. \quad (3.3)$$

When $|P_1(u_2, x)| = 1$, since $|C_2 \Delta C'_2 \Delta C_1| = 2|L_1| + 2\ell + 1$ and $g(G) = 2\ell + 1$, we have $|L_1| = |P'| = \ell$ by (3.2), implying $|P_1| = \ell$ and $|Q_1| = 1$ as P_1, P_2 are longest arrises in H . Hence, $d(H) = \ell - 1$. Then $G[V(C_1 \cup C'_2 \cup P_2)]$ is an odd K_4 -subdivision with difference $\ell - 2$, which is a contradiction to the choice of H . So $|P_1(u_2, x)| \geq 2$. Assume that $|Q_2| = 1$. Then $|L_1| = |P_1| = \ell$ by (3.1). Since $|P_1(u_2, x)| \geq 2$, the graph $G[V(C'_2 \cup C_2 \cup C_3)]$ is an odd K_4 -subdivision whose difference is at most $\ell - 2$, which is a contradiction to the choice of H as $d(H) = \ell - 1$. So $|Q_2| \geq 2$. Then $yu_3 \in E(H)$ by (3.3), implying $xu_1 \in E(H)$ by (3.2). Hence, $|C_4 \Delta C'_2| = 2 + 2|L_1|$ by (3.1) and (3.2), and so $|L_1| = \ell$ by (3.1) again. Since $|P_1| \geq |L_1|$, we have $|P_1| = \ell$ and $|Q_1| = 1$ by (3.1), which is a contradiction as $|Q_2| \geq 2$. \square

9.3. When $x \in \{u_3, u_4\}$, the vertices x and y are adjacent, that is, $xy \in \{e, f\}$.

Subproof. By symmetry we may assume that $x = u_3$. Assume to the contrary that x, y are not adjacent. Set $C'_3 = P'P_2(y, u_3)$. Since P' is a chordal path of C_3 , we have that C'_3 is an odd hole by Lemma 4 and (3.1). Since $C'_3 \Delta C_3$ is an even hole, $|Q_1| = |L_2| = 1$ by (3.1) and Lemma 4 again. Then $|P_1| = 2\ell - 1 > \ell$, which is a contradiction to (3.1). So $e = xy$. \square

9.4. When $x \in V(L_1^*)$, we have that $|Q_1| = 1$, $|P_1| = |L_1| = \ell$, $|P'| = 2\ell - 1$, $xu_3, yu_3 \in E(H)$ and xu_3yP' is an odd hole.

Subproof. Set $C'_4 = L_1(x, u_1)Q_2P_2(u_4, y)P'$. We claim that $C_4 \Delta C'_4$ is an odd hole. Assume to the contrary that $C_4 \Delta C'_4$ is an even hole. Since $x \neq u_3$, the subgraph $C_1 \cup (C_4 \Delta C'_4)$ is an induced theta subgraph. Hence, $xu_3 \in E(H)$ by (3.1) and Lemma 4. Similarly, $yu_3 \in E(H)$. Since P' is a chordal path of C_4 , we get a contradiction to Lemma 4. So the claim holds, implying that C'_4 is an even hole.

Since $x \neq u_3$, the graph $C_2 \cup C'_4$ is an induced theta subgraph of G . Moreover, since C'_4 is an even hole, $|Q_2| = 1$ by (3.1) and Lemma 4. Hence, $|P_1| = |L_1| = \ell$ by (3.1) again. Assume that y, u_3 are not adjacent. Since $C_1 \cup C'_4$ is an induced theta subgraph of G , we have $xu_1 \in E(H)$, implying $|P(x, u_3)| = \ell - 1$. Since P' is a chordal path of C_4 and $C_4 \Delta C'_4$ is an odd hole, $yu_3 \in E(H)$ by Lemma 4, a contradiction. Hence, $yu_3 \in E(H)$. By symmetry we have $xu_3 \in E(H)$. This proves 9.4. \square

9.5. Assume that P' has the structure stated as in 9.4. Then no vertex in $V(G) - V(H \cup P')$ has two neighbours in $H \cup P'$.

Subproof. Assume to the contrary that some vertex $a \in V(G) - V(H \cup P')$ has two neighbours a_1, a_2 in $H \cup P'$. Since no vertex has two neighbours in an odd hole, it follows from Lemma 8 (3) and 9.4 that a has exactly two neighbours in $H \cup P'$ with $a_1 \in V(H) - \{x, y, u_3\}$ and $a_2 \in V(P)$. When $xa_2 \notin E(P')$, let Q be the unique (y, a_1) -path in $G[V(P) \cup \{a_1, y\}]$. Since Q^* is a direct connection in $G \setminus \{e, f\}$ linking P_2^* and $H - V(P_2^*)$, by 9.2-9.4 and the symmetry between P' and Q , we have $a_1 \in \{x, u_3\}$, contrary to the fact $a_1 \in V(H) - \{x, y, u_3\}$. So $xa_2 \in E(P')$. Moreover, since $|P_1| = |L_1| = \ell \geq 5$ and $g(G) = 2\ell + 1$, we have $a_1 \notin V(P_1)$. Let u'_1 be the neighbour of u_1 in L_1 . When $a_1 \in V(C_2) - V(P_1)$, since aa_2 is a direct connection in $G \setminus \{u_1u'_1, u_3x\}$ linking L_1^* and $H - V(L_1^*)$, which is not possible by 9.4 and the symmetry between P_2 and L_1 . So $a \in V(L_1^*)$. Then xa_2aa_1 is a chordal path of C_1 with length 3, contrary to Lemma 4. \square

9.6. $x \in \{u_3, u_4\}$ and $xy \in \{e, f\}$.

Subproof. By 9.1-9.3, it suffices to show that $x \notin V(L_1^* \cup L_2^* \cup Q_1^* \cup Q_2^*)$. Assume not. By symmetry we may assume that $x \in V(L_1^*)$. By 9.4, we have that

$$xu_3 \in E(L_1), e = yu_3, |P'| = 2\ell - 1, |P_1| = |L_1| = \ell, \text{ and } |Q_1| = 1.$$

Since no 4-vertex-critical graph has a P_3 -cut by Lemma 5, to prove that 9.6 is true, it suffice to show that $\{x, y, u_3\}$ is a P_3 -cut of G . Suppose not. Let R be a shortest induced path in $G - \{x, y, u_3\}$ linking P and $H - \{x, y, u_3\}$. Let s and t be the ends of R with $s \in V(P)$. By 9.5, $|R| \geq 3$ and no vertex in $V(H \cup P') - \{x, y, u_3, s, t\}$ has a neighbour in R^* .

We claim that $t \notin V(L_1 \cup P_2) - \{x, y, u_3\}$. Suppose not. By symmetry we may assume that $t \in V(L_1) - \{x, u_3\}$. Let R_1 be the induced (y, t) -path in $G[V(P' \cup R) - \{x\}]$. When u_3 has no neighbour in R_1^* , set $R_2 := R_1$ and $C := R_2L_1(t, u_3)u_3y$. When u_3 has a neighbour in R_1^* , let $t' \in V(R_1^*)$ be a neighbour of u_3 closest to t and set $R_2 := u_3t'R_1(t', t)$ and $C := R_2L_1(t, u_3)$. Note that $C_4\Delta C$ is a hole, although C maybe not a hole. Since $C\Delta C_1\Delta C_2$ is an odd hole with length at least $2\ell + 3$ when C is an odd cycle, to prove the claim, it suffices to show that $|C|$ is odd. When x has a neighbour in R_2^* , since $|R_2| \geq 2\ell$ by (3.1) and the fact that $g(G) = 2\ell + 1$, the subgraph $C_4\Delta C$ is an even hole, which implies that C is an odd cycle. So we may assume that x has no neighbour in R_2^* . When u_3 is an end of R_2 , since R_2 is a chordal path of C_1 , it follows from Lemma 4 and (3.1) that C is an odd hole. When y is an end of R_2 and u_3yR_2 is a chordal path of C_1 , for the similar reason, C is an odd hole. Hence, we may assume that y is an end of R_2 and u_3yR_2 is not a chordal path of C_1 , implying $xs \in E(P')$ and $s \in V(R_2)$. Since $P \subset R_2$, we have $|R_2| > 2\ell$, so $C_4\Delta C$ is an even hole, implying that C is an odd cycle. Hence, this proves the claim.

By symmetry we may therefore assume that $t \in V(P_1) - \{u_1\}$. Let R_1 be the induced (y, t) -path in $G[V(P' \cup R) - \{x\}]$. By 9.1 and 9.2, either $xs \in E(P')$ and y has no neighbour in R or some vertex in $\{x, u_3\}$ has a neighbour in R_1^* . No matter which case happens, we have $|R_1| \geq 2\ell$. That is, $R_1P_2(y, u_4)$ is a chordal path of C_2 with length at least $3\ell - 1$, which is a contradiction to Lemma 4 as t, u_4 are non-adjacent. Hence, $\{x, y, u_3\}$ is a P_3 -cut of G . \square

By 9.6, there is a minimal vertex cut X of G with $\{u_3, u_4\} \subseteq X \subset N[\{u_3, u_4\}]$ and $\{u_3, u_4\} = X \cap V(H)$. Let G_1 be the induced subgraph of G whose vertex set consists of X and the vertex set of the component of $G - X$ containing P_2^* . Since $\ell \geq 5$, we have $|P_2| \geq 4$ by (3.1). If all induced (u_3, u_4) -paths in G_1 have length $|P_2|$, by Lemma 7, G has a degree-2 vertex, a K_1 -cut or a K_2 -cut, which is not possible as G is 4-vertex-critical. Hence, to finish the proof of Theorem 9, it suffices to show that all induced (u_3, u_4) -paths in G_1 have length $|P_2|$.

Let Q be an arbitrary induced (u_3, u_4) -path in G_1 . When $|L_1| \geq 2$, since QQ_2 is a chordal path of C_1 by Lemma 8 (3) and the definition of G_1 , we have $|QQ_2| = |Q_1P_1|$ by Lemma 4, so $|Q| = |P_1|$ by (3.1). Hence, by (3.1) we may assume that $|L_1| = 1$ and $|Q_1| = |P_1| = \ell$. Since Q_1L_2 is an induced (u_3, u_4) -path of length $\ell + 1$, either $|Q| = |P_2| = \ell$ or $|Q| \geq \ell + 1$ and $|Q|$ has the same parity as $\ell + 1$. Assume that the latter case happens. Without loss of generality we may further assume that Q is chosen with length at least $\ell + 1$ and with $|P_2 \cup Q|$ as small as possible. Since $|Q|$ and $|P_2|$ have different parity, $P_2 \cup Q$ is not bipartite. Moreover, by the choice of Q , the subgraph $P_2 \cup Q$ contains a unique cycle C and $|C|$ is odd. Since $Q = P_2 \Delta C$ is an induced path of length at least $\ell + 1$, we have $|C \cap Q| > |C \cap P_2| \geq 2$. So $C \Delta C_3 \Delta C_1$ is an odd hole of length at least $2\ell + 3$, which is not possible. \square

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