

# Hardinian Arrays

Robert Dougherty-Bliss<sup>a</sup>      Manuel Kauers<sup>b</sup>

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## Abstract

In 2014, R.H. Hardin contributed a family of sequences about king-moves on an array to the On-Line Encyclopedia of Integer Sequences (OEIS). The sequences were recently noticed in an automated search of the OEIS by Kauers and Koutschan, who conjectured a recurrence for one of them. We prove their conjecture as well as some older conjectures stated in the OEIS entries. We also have some new conjectures for the asymptotics of Hardin's sequences.

**Mathematics Subject Classifications:** 05A15, 33F10

## 1 Introduction

The On-Line Encyclopedia of Integer sequences [15] contains over 350,000 sequences and perhaps tens of thousands of conjectures about them. Here we resolve some of these conjectures related to a family of sequences due to R.H. Hardin.

For any positive integer  $r$ , let  $H_r(n, k)$  be the number of  $n \times k$  arrays which obey the following rules:

- The entry in position  $(1, 1)$  is 0, and the entry in position  $(n, k)$  is  $\max(n, k) - r - 1$ .
- The entry in position  $(i, j)$  must equal or be one more than each of the entries in positions  $(i - 1, j)$ ,  $(i, j - 1)$ , and  $(i - 1, j - 1)$ .
- The entry in position  $(i, j)$  must be within  $r$  of  $\max(i, j) - 1$ .

We call such an arrangement of numbers a Hardinian array.

Equivalently, Hardinian arrays can be defined in terms of the king-distance between entries, i.e., the length of the shortest path that a king on a chessboard can take to get from one entry to the other. Using this notion, we can say that a Hardinian array has a 0 at position at the top-left entry, the entry at the bottom-right entry is king-distance between the top-left corner and the bottom-right corner minus  $r$ , every entry increases

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<sup>a</sup>Rutgers University, New Brunswick NJ, USA ([robert.w.bliss@gmail.com](mailto:robert.w.bliss@gmail.com)).

<sup>b</sup>Institute for Algebra, J. Kepler University Linz, Austria ([manuel.kauers@jku.at](mailto:manuel.kauers@jku.at)).

by either 0 or 1 taking a single step towards left, down, or south-east, and the entry at every position is within  $r$  of the position's king-distance from the top-left corner.

For  $r = 1, 2, 3$ , Hardinian arrays are counted by the tables [A253026](#) [4], [A253223](#) [5], and [A253004](#) [3], respectively. Below is an example for  $r = 1$ ,  $n = 6$ , and  $k = 5$ .

$$\begin{bmatrix} 0 & 1 & 2 & 2 & 3 \\ 1 & 1 & 2 & 2 & 3 \\ 2 & 2 & 2 & 3 & 3 \\ 3 & 3 & 3 & 3 & 4 \\ 4 & 4 & 4 & 4 & 4 \\ 4 & 4 & 4 & 4 & 4 \end{bmatrix}$$

Hardin noticed several interesting patterns. For example, for every fixed  $r$  and  $k$ , the sequence  $H_r(n, k)$  seems to be a polynomial in  $n$  of degree  $r$  for sufficiently large  $n$ . He also conjectured an evaluation of the diagonal for  $r = 1$ , namely

$$H_1(n, n) = \frac{1}{3}(4^{n-1} - 1).$$

More recently, Kauers and Koutschan [8] performed an automated search for sequences in the OEIS which satisfy linear recurrences with polynomial coefficients. Hardin happened to submit the diagonal of  $r = 2$  as its own sequence, which led Kauers and Koutschan to conjecture a recurrence for  $f(n) = H_2(n, n)$ , namely

$$\begin{aligned} & 32(n+1)(2n+1)^2(1575n^6 + 21285n^5 + 117954n^4 + 343020n^3 + 551943n^2 \\ & \quad + 465785n + 161046)f(n) \\ & - 8(121275n^9 + 1933470n^8 + 13267683n^7 + 51280818n^6 + 122556360n^5 + 186866686n^4 \\ & \quad + 180574335n^3 + 105734340n^2 + 33718283n + 4443102)f(n+1) \\ & + 2(294525n^9 + 4763070n^8 + 33170868n^7 + 130145646n^6 + 315713355n^5 + 488415476n^4 \\ & \quad + 478464380n^3 + 283626704n^2 + 91378536n + 12137328)f(n+2) \\ & + (294525n^9 + 4668570n^8 + 31877118n^7 + 122735586n^6 + 292620525n^5 + 445804136n^4 \\ & \quad + 431097970n^3 + 252913504n^2 + 80866406n + 10688508)f(n+3) \\ & - (121275n^9 + 1961820n^8 + 13655808n^7 + 53503836n^6 + 129484209n^5 + 199650088n^4 \\ & \quad + 194784258n^3 + 114948300n^2 + 36871922n + 4877748)f(n+4) \\ & + 2(2n+7)(1575n^6 + 11835n^5 + 35154n^4 + 52554n^3 + 41382n^2 \\ & \quad + 16118n + 2428)(n+3)^2f(n+5) = 0. \end{aligned}$$

Our main results are that many of these conjectures are correct. In Section 2 we will prove Hardin's conjectured closed form for  $H_1(n, n)$  and extend this to a closed form for the rectangular case  $H_1(n, k)$ . In Section 3 we will prove that the conjectured recurrence of Kauers and Koutschan for  $H_2(n, n)$  is correct, and in fact that *every*  $H_r(n, n)$  satisfies a linear recurrence with polynomial coefficients. We conjecture asymptotic estimates for  $H_r(n, n)$  for all  $r \geq 2$ .

## 2 The case $r = 1$

This case can be settled by an elementary combinatorial argument. Let us first consider the diagonal and confirm the closed form representation conjectured by Hardin. In the following proof we index our arrays beginning from 0 rather than 1.

**Theorem 1.**  $H_1(n, n) = \frac{1}{3}(4^{n-1} - 1)$  for all  $n \geq 1$ .

*Proof.* Consider a valid  $n \times n$  array. Above the upper diagonal, draw a dividing path between row entries which are equal to their king-distance and less than their king-distance. Draw the same path below the diagonal, but make it with respect to columns. See Figure 1 for an example.

By the monotonicity rule, the upper path can only move down and to the right. Further, if the first entry to its right in row  $i$  is  $(i, j)$ , then the first entry to its right in row  $i + 1$  is either  $(i + 1, j)$  or  $(i + 1, j + 1)$ . Thus the upper-path essentially consists of two kinds of steps: down and right-down. The situation is mirrored in the lower path.

If the upper path does not divide row  $i$  just after the row's entry on the main diagonal, then the row is determined from the diagonal to the right endpoint. Entries between the diagonal and the path equal their king-distance, entries after the path equal one less than their king-distance, and the diagonal must equal  $i$  as its king-distance is  $i$  and to its right is an  $i + 1$ . The analogous statement is true for the lower path with respect to columns. Thus every entry is determined except for when both paths divide the  $i$ th row and column just after the diagonal. In fact, the *first* time this happens, the diagonal entry is still determined, as one of the entries above or to the left of the diagonal entry equals  $i$ .

In summary, the only entries not determined by these paths are the diagonal entries which both paths are adjacent to, except the first one and last one (by rule). If one path first touches the diagonal at position  $i$ , and the other at position  $j > i$ , then there are  $n - j - 2$  diagonal entries not determined. Of these entries, we may choose at most one to be the first less than its king-distance. After this choice all later entries must do the same. Thus each such pair of paths generates  $n - j - 1$  valid arrays.

If  $C(k)$  is the number of paths which are first adjacent to the diagonal at position  $k$ , then

$$H_1(n, n) = 2 \sum_{j=0}^{n-1} \sum_{i=0}^{j-1} C(i)C(j)(n - j - 1) + \sum_{j=0}^{n-1} C(j)^2(n - j - 1).$$

Because each path essentially has two steps to choose from, both of them moving one step closer to their end, we have  $C(k) = 2^{k-1}$  if  $k > 0$  and  $C(0) = 1$ . Evaluating the above summations and simplifying produces  $H_1(n, n) = (4^{n-1} - 1)/3$ .  $\square$

The double-path idea used in the proof above extends to the case of rectangular Hardinian arrays. The closed form expression for  $H_1(n, k)$  shown next confirms conjectures stated by Hardin for  $H_1(n, 1)$ ,  $H_1(n, 2)$ ,  $\dots$ ,  $H_1(n, 7)$ .

**Theorem 2.**  $H_1(n, k) = 4^{k-1}(n - k) + \frac{1}{3}(4^{k-1} - 1)$  for all  $n \geq k \geq 1$ .

$$\begin{bmatrix} 0 & 1 & 1 & 2 & 3 \\ 1 & 1 & 2 & 2 & 3 \\ 1 & 1 & 2 & 2 & 3 \\ 2 & 2 & 2 & * & 3 \\ 3 & 3 & 3 & 3 & 3 \end{bmatrix}$$

Figure 1: A generic  $5 \times 5$  matrix with two specific paths as constructed in the combinatorial proof of Theorem 1. Every entry is determined by the paths except the one labeled  $*$ , which may be 3 or 2.

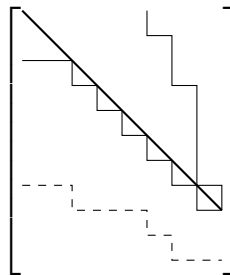


Figure 2: The generic picture for paths in the proof of Theorem 2. The lower two paths are examples of the two possible cases.

*Proof.* Draw the same paths indicated in the proof of Theorem 1. See Figure 2 for an example.

A lower path now is either adjacent to the diagonal at some point or not. The number of valid arrays where the lower path is adjacent to the diagonal at some point is  $H_1(k, k)$ . All other pairs of paths contribute only one valid array. There are  $n - k$  possible ending positions for a lower path which is never adjacent to the diagonal and  $2^{k-1}$  paths originating from each. Thus this case contributes  $(n - k)4^{k-1}$  valid arrays. Together this yields  $H_1(n, k) = 4^{k-1}(n - k) + H_1(k, k)$ .  $\square$

As these combinatorial arguments do not seem to extend to  $r > 1$ , we give some alternative proofs of Theorem 1. They all rely on the theorem of Gessel and Viennot [11, Theorem 10.13.1], which translates the counting problem into a determinant evaluation problem. We will evaluate the determinant in three different ways. The following notation will be used.

**Definition 3.** 1. For each positive integer  $n$ , let  $M(n)$  be the  $n \times n$  matrix of binomial coefficients

$$\left\{ \binom{u+v}{u} \right\}_{0 \leq u, v < n}.$$

Observe that rows and columns are indexed starting from zero.

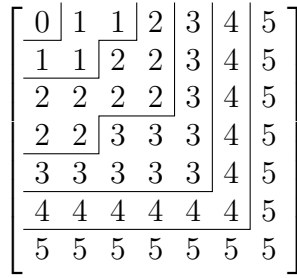


Figure 3: The contiguous regions of a Hardinian array are separated by a tuple of nonintersecting lattice walks starting on the left and ending at the top.

2. For any  $n \times n$  matrix  $A$ , any distinct row indices  $i_1, i_2, \dots, i_r \in \{0, \dots, n-1\}$  and distinct column indices  $j_1, j_2, \dots, j_r \in \{0, \dots, n-1\}$ , let  $A_{i_1, i_2, \dots, i_r}^{j_1, j_2, \dots, j_r}$  be the  $(n-r) \times (n-r)$  matrix obtained from  $A$  by deleting rows  $i_1, \dots, i_r$  and columns  $j_1, \dots, j_r$ .
3. For every  $n \geq 1$ , define

$$\Delta(n) = \det M(n)$$

$$\Delta(n)_{i_1, i_2, \dots, i_r}^{j_1, j_2, \dots, j_r} = \det M(n)_{i_1, i_2, \dots, i_r}^{j_1, j_2, \dots, j_r}.$$

**Lemma 4.**  $\Delta(n) = 1$  for all  $n$ .

*Proof.* Observe that  $M(n) = AB$  where  $A$  is the matrix whose entry at  $(u, v)$  is  $\binom{u}{v}$  and  $B$  is the matrix whose entry at  $(u, v)$  is  $\binom{v}{u}$ . This follows from Vandermonde's identity  $\binom{u+v}{v} = \sum_k \binom{u}{k} \binom{v}{k}$ . As  $A$  and  $B$  are triangular matrices with 1's on the diagonal, the claim follows from  $\Delta(n) = \det(M(n)) = \det(A) \det(B)$ .  $\square$

The key observation is that the valid  $n \times n$  arrays can be partitioned into contiguous regions, as shown in Figure 3. There is a region for 0, a region for 1, a region for 2, and so on. In the  $n \times n$  case, the region corresponding to  $k$  is obtained by beginning at the lowest occurrence of  $k$  in the first column, moving as far right as possible while only passing  $k$ 's, and moving up when stuck. For an  $n \times n$  Hardinian array this process always terminates in the first row.

**Proposition 5.**  $H_1(n, n) = \sum_{i=0}^{n-2} \sum_{j=0}^{n-2} \Delta(n-1)_i^j$  for all  $n \geq 1$ .

*Proof.* The  $n-1$  contiguous regions in a Hardinian array of size  $n \times n$  are separated by  $n-2$  nonintersecting lattice paths. These paths begin on one of the  $n-1$  edges between entries in the first column and end on one of the  $n-1$  edges between entries in the first row, using only steps to the right ( $\rightarrow$ ) and upwards ( $\uparrow$ ). Each Hardinian array corresponds to exactly one such set of paths.

In the other direction, each such set of paths corresponds to a Hardinian array. Given such a set, assign the induced regions the values  $0, 1, \dots, n-2$  in order from the top-left to the bottom-right. The top left will contain a 0, the bottom right will contain an

$n - 2$ , and adjacent entries differ by no more than 1. To see that the king-distance rule is not violated, note that it is not violated at the entries before the boundaries on the first column and first row—because at most one entry does not have a path just before it—and that these points have the largest king-distance of any entry reached using the available steps.

It follows that the number of Hardinian arrays of size  $n \times n$  equals the number of sets of nonintersecting lattice paths we have described. If we label the possible starting and ending positions  $0, 1, \dots, n - 2$ , then there are altogether  $\binom{u+v}{v}$  paths from  $u$  to  $v$ , for any  $u$  and  $v$ .

Consider the set of paths where  $i$  is the unique unchosen startpoint and  $j$  the unique unchosen endpoint. In this case the  $k$ th path ( $k = 0, \dots, n - 3$ ) starts at  $k + [i \leq k]$  and ends at  $k + [j \leq k]$ . By the theorem of Gessel and Viennot, the number of such sets of paths is the determinant of the  $(n - 2) \times (n - 2)$  matrix whose entry at position  $(u, v)$  is  $\binom{u+v+[i \leq u]+[j \leq v]}{v+[j \leq v]}$ . This determinant equals  $\Delta(n - 1)_i^j$ . It follows that  $H_1(n, n)$  is the sum of  $\Delta(n - 1)_i^j$  over all possible rows  $i$  and columns  $j$ .  $\square$

The proposition reduces the enumeration problem to the problem of evaluating a sum of determinants. This can be done as follows.

*Second proof of Theorem 1.* Let  $\tilde{M}(n)$  be the  $(n + 1) \times (n + 1)$  matrix obtained from  $M(n)$  by first attaching an additional row  $1, -1, 1, -1, \dots$  at the top and then an additional column  $0, -1, 1, -1, 1, \dots$  at the left, e.g.,

$$\tilde{M}(5) = \begin{vmatrix} 0 & 1 & -1 & 1 & -1 & 1 \\ -1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 1 & 2 & 3 & 4 & 5 \\ -1 & 1 & 3 & 6 & 10 & 15 \\ 1 & 1 & 4 & 10 & 20 & 35 \\ -1 & 1 & 5 & 15 & 35 & 70 \end{vmatrix}.$$

By expanding along the first row and then along the first column, we have  $\det \tilde{M}(n) = \sum_{i=0}^{n-1} \sum_{j=0}^{n-1} \Delta(n)_i^j$ .

It remains to determine the determinant of  $\tilde{M}(n)$ .

Subtract the  $(n - 2)$ nd row from the  $(n - 1)$ st, then the  $(n - 3)$ rd row from the  $(n - 2)$ nd,



satisfies the recurrence

$$s(n + 2) = 5s(n + 1) - 4s(n)$$

for all  $n \geq 1$ . Together with the initial values  $s(1) = 1$  and  $s(2) = 5$ , the claimed closed form expression now follows again by induction.  $\square$

While the sum  $\Delta(n)_i^j = \sum_{\ell=0}^{n-1} \binom{\ell}{i} \binom{\ell}{j}$  does not have a hypergeometric closed form, it does simplify in the special case  $j = n - 1$ . For this special case, it is obvious that  $\Delta(n)_i^{n-1} = \sum_{\ell=0}^{n-1} \binom{\ell}{i} \binom{\ell}{n-1} = \binom{n-1}{i}$ , because only the summand with  $\ell = n - 1$  is nonzero. Taking the knowledge of this special case for granted, we can give a fourth proof of Theorem 1.

*Fourth proof of Theorem 1.* Dodgson's identity (cf. Prop. 10 of Krattenthaler's tutorial on evaluating determinants [10]) says that

$$\det(A) \det(A_{i,n-1}^{j,n-1}) = \det(A_i^j) \det(A_{n-1}^{n-1}) - \det(A_i^{n-1}) \det(A_{n-1}^j)$$

for every  $n \times n$  matrix  $A$ . (Actually, Krattenthaler states the equation for  $i = j = 0$ , but it is easily seen that it holds for arbitrary  $i$  and  $j$ , because we can multiply  $A$  with suitable permutation matrices from the left and the right in order to reduce to the case  $i = j = 0$ .)

Consider  $A = M(n)$  and observe that  $A_{n-1}^{n-1} = M(n - 1)$ . Then, because of  $\Delta(n) = \Delta(n - 1) = 1$  it follows that

$$\Delta(n - 1)_i^j = \Delta(n)_i^j - \Delta(n)_i^{n-1} \Delta(n)_{n-1}^j.$$

Using the identities  $\Delta(n)_i^{n-1} = \binom{n-1}{i}$  and  $\Delta(n)_{n-1}^j = \binom{n-1}{j}$  observed just before this proof, it follows that

$$\Delta(n)_i^j = \Delta(n - 1)_i^j + \binom{n - 1}{i} \binom{n - 1}{j}.$$

Summing over all  $i$  and  $j$  gives

$$s(n) = s(n - 1) + 4^{n-1},$$

and with  $s(1) = 1$ , the claim follows again by induction.  $\square$

### 3 The case $r \geq 2$

Via the theorem of Gessel and Viennot, we also have access to the sequences  $H_r(n, n)$  for  $r > 1$ . The argument is the same as for  $r = 1$ , except that now a Hardinian array of size  $n \times n$  consists of  $n - r$  contiguous regions, separated by  $n - r - 1$  nonintersecting lattice paths, whose start points and end points are taken from the set  $\{0, \dots, n - 2\}$ . According to Gessel and Viennot,  $\Delta(n - 1)_{i_1, \dots, i_r}^{j_1, \dots, j_r}$  is the number of sets of  $n - r - 1$  nonintersecting lattice walks whose start points are  $\{0, \dots, n - 2\} \setminus \{i_1, \dots, i_r\}$  and whose end points are  $\{0, \dots, n - 2\} \setminus \{j_1, \dots, j_r\}$ .



In order to deal with these determinants, it helps to observe that Dodgson's identity quoted in the fourth proof of Theorem 1 is a special case of a more general identity due to Jacobi [6, 14, 1]: For an  $n \times n$  matrix  $A$  and two choices  $0 \leq i_1 < i_2 < \dots < i_r < n$  and  $0 \leq j_1 < j_2 < \dots < j_r < n$  of indices, form the  $r \times r$  matrix  $B$  whose entry at  $(u, v)$  is defined as  $\det(A_{i_u}^{j_v})$ . Then Jacobi's identity says

$$\det(A)^{r-1} \det(A_{i_1, \dots, i_r}^{j_1, \dots, j_r}) = \det(B).$$

For example, for  $r = 2$  we obtain

$$\det(A) \det(A_{i_1, i_2}^{j_1, j_2}) = \begin{vmatrix} \det(A_{i_2}^{j_2}) & \det(A_{i_2}^{j_1}) \\ \det(A_{i_1}^{j_2}) & \det(A_{i_1}^{j_1}) \end{vmatrix} = \det(A_{i_1}^{j_1}) \det(A_{i_2}^{j_2}) - \det(A_{i_1}^{j_2}) \det(A_{i_2}^{j_1}),$$

and setting  $i_2 = j_2 = n - 1$  gives Dodgson's version.

**Theorem 6.** *For every  $r \geq 2$ , the sequence  $H_r(n, n)$  is D-finite. In particular, the sequences [A253217](#) ( $r = 2$ ) and [A252998](#) ( $r = 3$ ) are D-finite.*

*Proof.* For  $A = M(n)$ , Jacobi's identity implies

$$\Delta(n)_{i_1, \dots, i_r}^{j_1, \dots, j_r} = \begin{vmatrix} \Delta(n)_{i_1}^{j_1} & \dots & \Delta(n)_{i_1}^{j_r} \\ \vdots & \ddots & \vdots \\ \Delta(n)_{i_r}^{j_1} & \dots & \Delta(n)_{i_r}^{j_r} \end{vmatrix}$$

For every fixed  $r$ , the determinant on the right is D-finite because it depends polynomially on quantities which we have recognized in the previous section as being D-finite. It follows that the left hand side is D-finite, and consequently,

$$H_r(n, n) = \sum_{0 \leq i_1 < \dots < i_r \leq n-2} \sum_{0 \leq j_1 < \dots < j_r \leq n-2} \Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r}$$

is D-finite, too. □

Theorem 6 is not quite enough to confirm the correctness of the recurrence equation Kauers and Koutschan obtained for  $H_2(n, n)$  via guessing [8]. The theorem only implies that the sequence satisfies *some* recurrence. In order to explicitly construct a recurrence,

we have to evaluate the two 6-fold sums

$$\begin{aligned}
 S_1(n) &= \sum_{i_1 \geq 0} \sum_{i_2 > i_1} \sum_{j_1 \geq 0} \sum_{j_2 > j_1} \sum_{u=0}^n \sum_{v=0}^n \binom{u}{i_1} \binom{u}{j_1} \binom{v}{i_2} \binom{v}{j_2} \\
 &= \sum_{u=0}^n \sum_{v=0}^n \underbrace{\left( \sum_{i_1 \geq 0} \sum_{i_2 > i_1} \binom{u}{i_1} \binom{v}{i_2} \right)}_{=:s(u,v)} \underbrace{\left( \sum_{j_1 \geq 0} \sum_{j_2 > j_1} \binom{u}{j_1} \binom{v}{j_2} \right)}_{=:s(u,v)} \text{ and} \\
 S_2(n) &= \sum_{i_1 \geq 0} \sum_{i_2 > i_1} \sum_{j_1 \geq 0} \sum_{j_2 > j_1} \sum_{u=0}^n \sum_{v=0}^n \binom{u}{i_1} \binom{u}{j_2} \binom{v}{i_2} \binom{v}{j_1} \\
 &= \sum_{u=0}^n \sum_{v=0}^n \underbrace{\left( \sum_{i_1 \geq 0} \sum_{i_2 > i_1} \binom{u}{i_1} \binom{v}{i_2} \right)}_{=:s(u,v)} \underbrace{\left( \sum_{j_1 \geq 0} \sum_{j_2 > j_1} \binom{v}{j_1} \binom{u}{j_2} \right)}_{=:s(v,u)}.
 \end{aligned}$$

It seems best to do this using generating functions. We have

$$\sum_{u=0}^{\infty} \sum_{v=0}^{\infty} s(u,v) x^u y^v = \frac{y}{(1-x-y)(1-2y)}.$$

The generating functions of  $s(u,v)^2$  and  $s(u,v)s(v,u)$  can be expressed as Hadamard products. As explained in [2], Hadamard products can be rephrased as residues, and residues can be computed via creative telescoping [16]. Using Koutschan's implementation [9], it is easy to prove

$$\begin{aligned}
 &\frac{y}{(1-x-y)(1-2y)} \odot_{x,y} \frac{y}{(1-x-y)(1-2y)} \\
 &= \frac{y}{2x+2y-1} \left( \frac{1}{\sqrt{x^2-2x(y+1)+(y-1)^2}} + \frac{2}{4y-1} \right) \\
 &\frac{y}{(1-x-y)(1-2y)} \odot_{x,y} \frac{x}{(1-x-y)(1-2x)} \\
 &= \frac{1}{2(2x+2y-1)} \left( \frac{x+y-1}{\sqrt{x^2-2x(y+1)+(y-1)^2}} + 1 \right),
 \end{aligned}$$

respectively. Summing  $u$  from 0 to  $n$  and  $v$  from 0 to  $m$  amounts to multiplying these series by  $\frac{1}{(1-x)(1-y)}$ , and setting  $m$  to  $n$  amounts to taking the diagonals of the resulting bivariate series:

$$\begin{aligned}
 &\text{diag} \frac{1}{(1-x)(1-y)} \frac{y}{2x+2y-1} \left( \frac{1}{\sqrt{x^2-2x(y+1)+(y-1)^2}} + \frac{2}{4y-1} \right), \\
 &\text{diag} \frac{1}{(1-x)(1-y)} \frac{1}{2(2x+2y-1)} \left( \frac{x+y-1}{\sqrt{x^2-2x(y+1)+(y-1)^2}} + 1 \right),
 \end{aligned}$$

respectively. As diagonals can also be rephrased as residues (cf. again [2] for a detailed discussion), we can apply creative telescoping to obtain linear differential operators annihilating these series. Their least common left multiple is an annihilator of the generating function of  $H_2(n, n)$ .

In the end, we obtained a linear differential operator of order 10 with polynomial coefficients of degree 43. With this certified operator at hand, we can prove that the guessed recurrence of Kauers and Koutschan is correct.

In principle, we could derive a recurrence for  $H_r(n, n)$  for any  $r \geq 2$  in the same way, but already for  $r = 3$  the computations become too costly. We can however use the formula

$$H_r(n, n) = \sum_{0 \leq i_1 < \dots < i_r \leq n-2} \sum_{0 \leq j_1 < \dots < j_r \leq n-2} \Delta(n-1)_{i_1, \dots, i_r}^{j_1, \dots, j_r}$$

to compute some more terms of the sequences. In order to do this efficiently, we can recycle the idea of the second proof of Theorem 1 and translate some of the summation signs into additional rows and columns of the determinant. For example, for  $r = 3$  we have

$$H_r(n, n) = \sum_{i=0}^{n-2} \sum_{j=0}^{n-2} |\det(A_{i,j})|$$

where  $A_{i,j}$  is the matrix obtained from  $M(n-1)$  by removing the  $i$ th row and the  $j$ th column and adding a row with alternating signs in the column range  $0 \dots j-1$  followed by zeros and an additional row with zeros in the column range  $0 \dots j-1$  followed by alternating signs; and similarly two additional columns. For example, for  $n = 8, i = 4, j = 5$  we have

$$A_{i,j} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 1 & -1 \\ 0 & 0 & -1 & 1 & -1 & 1 & 0 & 0 \\ 0 & -1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 2 & 3 & 4 & 6 & 7 \\ 0 & -1 & 1 & 3 & 6 & 10 & 21 & 28 \\ -1 & 0 & 1 & 5 & 15 & 35 & 126 & 210 \\ 1 & 0 & 1 & 6 & 21 & 56 & 252 & 462 \\ -1 & 0 & 1 & 7 & 28 & 84 & 462 & 924 \end{pmatrix} \begin{array}{l} \text{extra} \\ \text{rows} \\ \\ \\ \text{ith row} \\ \text{deleted} \\ \\ \end{array} \begin{array}{l} \\ \\ \text{extra} \\ \text{columns} \\ \\ \text{jth column} \\ \text{deleted} \end{array}$$

With this optimization, it is not difficult to compute the first 100 terms, and using these, the technique of [7] is able to guess a convincing recurrence equation of order 9 and degree 36. It is not reproduced here.

For  $r = 4$ , we explicitly delete two rows and columns and add two rows and columns with alternating signs, as shown in Figure 4 on the left. This allows us to reduce the original 8-fold sum to a 4-fold sum. A 4-fold sum is also sufficient for  $r = 5$ , where we can even eliminate six summations by adding extra rows and columns, as shown in Figure 4 on the right. By computing the sums over all these determinants, we were able to determine the first 70 terms of the sequences  $H_4(n, n)$  and  $H_5(n, n)$ . Unfortunately, these terms were not sufficient to find a recurrence by guessing.

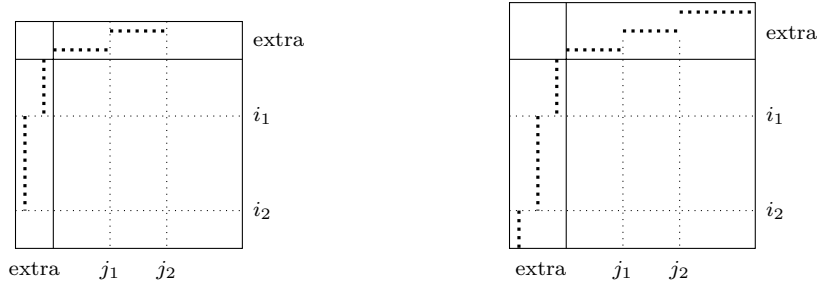


Figure 4: Left: the computation of  $\sum_{i_1 < i_2 < i_3 < i_4} \sum_{j_1 < j_2 < j_3 < j_4} \Delta(n-1)_{i_1, i_2, i_3, i_4}^{j_1, j_2, j_3, j_4}$  is equivalent to the computation of the sum over  $i_1, i_2$  and  $j_1, j_2$  of the determinants constructed as shown in the figure. Light dots indicated omitted rows and columns; strong dots indicate regions filled with alternating signs.

Right: the computation of  $\sum_{i_1 < i_2 < i_3 < i_4 < i_5} \sum_{j_1 < j_2 < j_3 < j_4 < j_5} \Delta(n-1)_{i_1, i_2, i_3, i_4, i_5}^{j_1, j_2, j_3, j_4, j_5}$  is equivalent to the computation of the sum over  $i_1, i_2$  and  $j_1, j_2$  of the determinants constructed as shown in the figure.

However, the terms are enough to obtain convincing conjectured expressions for their asymptotics. We obtained the following conjectures:

$r$	asymptotics	remark
0	1	trivial
1	$\frac{1}{2^{23}} 4^n$	by Theorem 1
2	$\frac{1}{2^{23} 4 \pi} 16^n n^{-1}$	from the proven recurrence
3	$\frac{1}{2^{23} 9 \pi} 64^n n^{-3}$	from the guessed recurrence
4	$\frac{2^2}{3^{16} \pi^2} 256^n n^{-6}$	from the first 70 terms
5	$\frac{2^4}{3^{23} \pi^2} 1024^n n^{-10}$	from the first 70 terms

Altogether, it seems that for every  $r \geq 0$ , we have

$$H_r(n, n) \sim c 2^{2rn} n^{-\binom{r}{2}} \quad (n \rightarrow \infty)$$

for some constant  $c$  that can be expressed as a power product of 2, 3, and  $\pi$ .

At least for specific values of  $r$ , it might be possible to prove these conjectured asymptotic formulas using the powerful techniques of analytic combinatorics in several variables [13, 12]. However, in order to invoke these techniques, we would need to know more about the bivariate sequences  $H_r(n, k)$ . Unfortunately, while we found an explicit expression for  $H_1(n, k)$ , we were not able to show that  $H_r(n, k)$  is D-finite as a bivariate sequence in  $n$  and  $k$  for any  $r \geq 2$ , although we suspect it to be.

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