

A note on the minimum size of Turán systems

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Abstract

For positive integers $n \geq s > r$, a *Turán (n, s, r) -system* is an n -vertex r -graph in which every set of s vertices contains at least one edge. Let $T(n, s, r)$ denote the the minimum size of a Turán (n, s, r) -system.

Upper bounds on $T(n, s, r)$ were established by Sidorenko [20] for the case $s - r = \Omega(r/\ln r)$, based on a construction of Frankl–Rödl [7], and by a number of authors in the case $s - r = O(1)$. Motivated by these results and recent work [16] of the second author, we investigate in this note the intermediate regime where $s = s(r)$ satisfies both $s - r = \Omega(1)$ and $s - r = O(r/\ln r)$, and establish upper bounds for $T(n, s, r)$ in this range as $r \rightarrow \infty$.

Mathematics Subject Classifications: 05D05

1 Introduction

Given an integer $r \geq 2$, an **r -uniform hypergraph** (henceforth an **r -graph**) \mathcal{H} is a collection of r -subsets of some set V . We call V the **vertex set** of \mathcal{H} and denote it by $V(\mathcal{H})$. When V is understood, we usually identify a hypergraph \mathcal{H} with its set of edges.

For positive integers $n \geq s > r$, a **Turán (n, s, r) -system** is an r -graph \mathcal{H} on an n -set V such that every s -subset $S \subseteq V$ contains at least one edge from \mathcal{H} . Denote by $T(n, s, r)$ the smallest **size** (i.e. the number of edges) of a Turán (n, s, r) -system. Observe that $T(n, s, r) = \binom{n}{r} - \text{ex}(n, K_s^r)$, where $\text{ex}(n, K_s^r)$ denotes the Turán number of the complete r -graph on s vertices K_s^r . A simple averaging argument shows that $\text{ex}(n, K_s^r)/\binom{n}{r}$ is non-increasing (see e.g. [10]), and hence the following limit exists:

$$t(s, r) := \lim_{n \rightarrow \infty} \frac{T(n, s, r)}{\binom{n}{r}}. \quad (1)$$

Determining the value of $t(s, r)$ is a central topic in Extremal Combinatorics. The seminal paper of Turán [22] established that $t(s, 2) = \frac{1}{s-1}$ for all $s \geq 3$ (with the case $s = 3$

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solved earlier by Mantel [14]). Erdős [6] offered \$500 for the determination of $t(s, r)$ for a single pair s, t with $s > r \geq 3$. This prize is still unclaimed, despite decades of active attempts. Turán and other researchers conjectured that $t(s, 3) = \frac{4}{(s-1)^2}$ for $s \geq 4$. Various constructions achieving this bound are known (see e.g. [19]). For $r \geq 4$, there is no general conjectured value for $t(s, r)$, except for the case $(s, r) = (5, 4)$ (see [9, 15]). For further related results, we refer the reader to surveys such as [8, 5, 19, 11].

In this note, we focus on the case where $r \rightarrow \infty$, and all asymptotics are taken with respect to r unless otherwise specified. The trivial lower bound is $t(s, r) \geq 1/\binom{s}{r}$, which follows from the monotonicity of the ratio in (1). For convenience, let us define

$$\mu(s, r) := t(s, r) \cdot \binom{s}{r}.$$

Note that $\mu(s, r)$ is always at least 1.

The best-known general lower bound, $t(s, r) \geq 1/\binom{s-1}{r-1}$ (i.e. $\mu(s, r) \geq s/r$) is due to de Caen [3]. In particular, $t(r+1, r) \geq 1/r$, a result that was independently proved by de Caen [4], Sidorenko [18], and Tazawa–Shirakura [21]. Further improvements on $t(r+1, r)$ in lower order terms were made by Giraud (unpublished, see [5, Page 189]), Chung–Lu [2], and Lu–Zhao [13].

Improving the previous upper bounds established by Sidorenko [17], Kim–Roush [12], Frankl–Rödl [7], and Sidorenko [20], the second author established the following upper bound, which disproved the \$500 conjecture of de Caen [5] that $r \cdot t(r+1, r) \rightarrow \infty$.

Theorem 1 ([16]). *For every integer $R \geq 1$, it holds that*

$$\mu(r+R, r) \leq \alpha + o(1), \quad \text{as } r \rightarrow \infty,$$

where $\alpha := (c_0 + 1)^{R+1}/c_0^R$ with $c_0 = c_0(R)$ being the largest real root of the equation $e^x = (x+1)^{R+1}$. In particular, $\mu(r+1, r) \leq 4.911$ for all sufficiently large r .

An immediate corollary of Theorem 1 (for derivation see [16, Corollary 1.3]) is that, for every sufficiently large R ,

$$\limsup_{r \rightarrow \infty} \mu(r+R, r) \leq R \ln R + 3R \ln \ln R = (1 + o(1))R \ln R. \quad (2)$$

This improves asymptotically the previous bound by Frankl–Rödl [7] which states that, for any fixed $R \geq 1$, we have

$$\mu(r+R, r) \leq (1 + o(1))R(R+4) \ln r, \quad \text{as } r \rightarrow \infty.$$

For the case $R \geq \frac{r}{\log_2 r}$, Sidorenko [20] (by analyzing the construction of Frankl–Rödl [7] in this regime) proved that

$$\mu(r+R, r) \leq (1 + o(1))R \ln \binom{r+R}{R}. \quad (3)$$

Although Turán systems were actively studied, it seems that no general upper bounds on $t(r + R, r)$ have been published in the intermediate regime $1 \ll R \leq r/\log_2 r$. This is the case we address in this note. In brief, we show that Sidorenko's bound in (3) applies in the whole range as long as R is sufficiently large, while the asymptotic bound in (2) can be extended from constant R to any $R = o(\sqrt{r})$.

Theorem 2. *For every $\varepsilon > 0$, there exists r_0 such that the following statements hold for all r, R satisfying $R \geq r_0$.*

(i) *It holds that $\mu(r + R, r) \leq (1 + \varepsilon)R \ln \binom{r+R}{R}$.*

(ii) *Suppose that $R \leq \sqrt{18r \ln r}$. Then*

$$\mu(r + R, r) \leq e^{18R^2/r} \cdot (1 + \varepsilon)R \ln R.$$

Our construction for Theorem 2 (i) will be a straightforward modification of the random construction used by Frankl and Rödl [7, Theorem 3], while the construction for Theorem 2 (ii) will be an extension of the recursive construction of the second author [16].

2 Proofs

In this section, we prove Theorem 2. We will use the following notation. For an integer $n \geq 1$ and a set X , we denote $[n] := \{1, \dots, n\}$ and $\binom{X}{n} := \{Y \subseteq X : |Y| = n\}$.

Let us begin with the proof of Theorem 2 (i).

Proof of Theorem 2 (i). Given $\varepsilon > 0$, let r_0 be sufficiently large. Take any r, R with $R \geq r_0$. By (3), we can assume that $R \leq \frac{r}{\ln r}$. Let $s := r + R$. Define

$$N := \left\lfloor r(r-1) \binom{s}{R} / (2R) \right\rfloor \quad \text{and} \quad \ell := \left\lfloor \binom{s}{R} / \ln \left(\binom{s}{R}^2 \binom{N-s}{R} \right) \right\rfloor.$$

In the rest of the proof, we repeat the construction of Frankl and Rödl [7, Theorem 3], arguing that our choices of N and ℓ give the stated bound.

Consider a random colouring $c : \binom{[N]}{r} \rightarrow [\ell]$. For each s -set $A \in \binom{[N]}{s}$ we have a bad event that some colour is not present in $\binom{A}{r}$. Its probability p is at most $\ell(1 - 1/\ell)^{\binom{s}{r}} \leq \ell e^{-(\binom{s}{r})/\ell}$ (note that $\binom{s}{r} = \binom{s}{s-r} = \binom{s}{R}$). Define the obvious dependency graph D (with loops) on $\binom{[N]}{s}$ where $A \sim B$ if $|A \cap B| \geq r$. It is regular of degree $\Delta := \sum_{i=r}^s \binom{s}{i} \binom{N-s}{s-i}$.

Our goal is to use the Lovász Local Lemma to deduce the existence of a colouring in which none of the bad events occur. In order to apply the variant of the Lovász Local Lemma given in [1, Corollary 5.1.2], it suffices to check that $ep\Delta < 1$, for which it suffices to prove

$$e^{\binom{s}{R}/\ell} > e\ell\Delta. \tag{4}$$

First, we estimate Δ . Consider the ratio of two consecutive terms:

$$\frac{\binom{s}{i+1} \binom{N-s}{s-i-1}}{\binom{s}{i} \binom{N-s}{s-i}} = \frac{(s-i)^2}{(i+1)(N-2s+i+1)}.$$

Since $3 \leq R \leq s/2$ and $N \geq \binom{s}{R} \geq \binom{s}{3}$, this ratio is at most $1/2$ for every $i \in [r, s]$, provided r_0 is sufficiently large. Thus, we can bound

$$\Delta \leq 2 \binom{s}{r} \binom{N-s}{s-r} = 2 \binom{s}{R} \binom{N-s}{R}. \quad (5)$$

Therefore, by the choice of ℓ , we have

$$e^{\binom{s}{R}/\ell} \geq \binom{s}{R}^2 \binom{N-s}{R} \geq e \cdot \frac{\binom{s}{R}}{2 \ln \binom{s}{R}} \cdot 2 \binom{s}{R} \binom{N-s}{R} \geq e\ell\Delta,$$

as desired. (In fact, our definition of ℓ comes from (4), given N .)

Thus by the Lovász Local Lemma, there exists a colouring c with no bad events, meaning that each colour gives a Turán (N, s, r) -system on $[N]$. Let $\mathcal{A} \subseteq \binom{[N]}{r}$ be the r -graph formed by the least frequent colour. We have

$$|\mathcal{A}| \leq \frac{1}{\ell} \binom{N}{r}. \quad (6)$$

Now let $n := mN$ with integer $m \rightarrow \infty$. Our Turán (n, s, r) -system \mathcal{B} on $[n]$ is made of a blowup of \mathcal{A} plus all r -sets that intersect at least one part in more than one vertex. We have

$$|\mathcal{B}| \leq m^r |\mathcal{A}| + N \binom{m}{2} \binom{mN-2}{r-2} \leq m^r \cdot \frac{1}{\ell} \binom{N}{r} + \frac{r(r-1)}{2N} \binom{mN}{r} \leq \binom{mN}{r} f,$$

where $f := \frac{1}{\ell} + \frac{r(r-1)}{2N}$. Since f does not depend on m , it gives an upper bound on $t(r+R, r)$ and thus $f \cdot \binom{s}{R} \geq \mu(s, r)$. Note that we have $N \leq \binom{s}{R}^2$ and $\ln N \leq 2R \ln s$. Therefore,

$$\ell \geq \frac{\binom{s}{R}}{2 \ln \binom{s}{R} + R \ln N} - 1$$

can be forced to be arbitrarily large if r_0 was sufficiently large. Thus the rounding in the definition of ℓ gives only a multiplicative term that is arbitrarily close to 1 and we have

$$\begin{aligned} \binom{s}{R} f &\leq (1 + \varepsilon/4) \left(2 \ln \binom{s}{R} + R \ln N + \frac{r(r-1)\binom{s}{R}}{2N} \right) \\ &\leq (1 + \varepsilon/2) \left(2 \ln \binom{s}{R} + R \left(2 \ln r + \ln \binom{s}{R} \right) + R \right) \leq (1 + \varepsilon) R \ln \binom{s}{R}, \end{aligned}$$

proving Theorem 2 (i). Note that our choice of N was determined by using the fact that the function $\ln x + c/x$, with fixed $c > 0$, is minimized on the interval $(0, \infty)$ at $x = c$. \square

Next, we prove Theorem 2 (ii). Before doing so, let us present some preliminary results. The following lemma is derived from the proof of [16, Lemma 2.3].

Lemma 3 ([16]). *For all integers $r, R \geq 1$, $k \in [R, r-1]$, and a real number $c \in [0, \binom{k}{R}]$, it holds that*

$$\mu(r+R, r) \leq \binom{r+R}{R} \left(\frac{c}{\binom{k}{R}} + \frac{\mu(r-k+R, r-k)}{e^c \cdot \binom{r-k+R}{R}} \right). \quad (7)$$

Sketch of Proof. Let S be a random subset of $\binom{[n]}{k-R}$ where each $(k-R)$ -set is included into S with probability $p := c/\binom{k}{R}$. Let $S^* \subseteq \binom{[n]}{r}$ consist of those r -sets $\{x_1 < \dots < x_r\}$ such that $\{x_1, \dots, x_{k-R}\} \in S$, that is, we include an r -set into S^* if its initial $(k-R)$ -segment is in S . Let $T \subseteq \binom{[n]}{k}$ consist of those k -sets X such that $\binom{X}{k-R} \cap S = \emptyset$, that is, X does not contain any element of S as a subset. For each $y \in [n]$, take a minimum Turán $(n-y, r-k+R, r-k)$ -system F_y on $\{y+1, \dots, n\}$. Let $T^* \subseteq \binom{[n]}{r}$ be the union over $Y \in T$ of the r -graphs $\{Y \cup Z : Z \in F_{\max Y}\}$. Informally speaking, we extend every $Y \in T$ by a minimum Turán system to the right of Y .

It is easy to check that $G := S^* \cup T^*$ is a Turán $(n, r+R, r)$ -system, regardless of the choice of S . By taking S such that $|G|$ is at most its expected value, it is routine to see that

$$\begin{aligned} T(n, r+R, r) &\leq \mathbb{E}|S^*| + \mathbb{E}|T^*| \\ &= p \binom{n}{r} + \sum_{y=k}^n (1-p) \binom{k}{R} \binom{y-1}{k-1} \cdot T(n-y, r-k+R, r-k) \\ &\leq \left(\frac{c}{\binom{k}{R}} + \frac{e^{-c}}{\binom{r-k+R}{R}} \mu(r-k+R, r-k) \right) \binom{n}{r}, \end{aligned}$$

giving the required bound. \square

Fact 4. *For any integers $r_1 \geq r_2 > R$, we have*

$$\binom{r_1}{R} / \binom{r_2}{R} = \prod_{i=0}^{R-1} \frac{r_1 - i}{r_2 - i} \leq \left(\frac{r_1 - R}{r_2 - R} \right)^R.$$

Fact 5. *Let $r \geq 1, R \geq 1$ be integers and δ be a real number satisfying $18R^2/r \leq \delta \leq R$. Let $k := \lceil \frac{Rr}{R+\delta} \rceil + R$. Then*

$$k \leq r-1, \quad \frac{r}{k-R} \leq 1 + \frac{\delta}{R}, \quad \text{and} \quad \frac{r}{r-k} \leq \frac{3R}{\delta}.$$

Proof of Fact 5. Since $\delta \geq \frac{18R^2}{r}$, straightforward calculations show that

$$\begin{aligned} \frac{\delta r}{R+\delta} - R - 2 &= \frac{\delta r - (R+\delta)(R+2)}{R+\delta} \geq \frac{18R^2 - (R+R)(R+2)}{R+\delta} \geq 0 \quad \text{and} \\ \frac{Rr}{R+\delta/2} - \frac{Rr}{R+\delta} &= \frac{\delta r R}{(R+\delta)(2R+\delta)} \geq \frac{18R^3}{(R+R)(2R+R)} \geq R+1. \end{aligned}$$

It follows that

$$\begin{aligned} k &\leq \frac{Rr}{R+\delta} + 1 + R = r - 1 - \left(\frac{\delta r}{R+\delta} - R - 2 \right) \leq r - 1, \\ \frac{r}{k-R} &\leq \frac{r}{\frac{Rr}{R+\delta} + R - R} = \frac{R+\delta}{R} = 1 + \frac{\delta}{R}, \quad \text{and} \\ \frac{r}{r-k} &\leq \frac{r}{r - \frac{Rr}{R+\delta} - 1 - R} \leq \frac{r}{r - \frac{Rr}{R+\delta/2}} = \frac{2R+\delta}{\delta} \leq \frac{3R}{\delta}, \end{aligned}$$

which proves Fact 5. \square

We are now ready to present the proof of Theorem 2 (ii).

Proof of Theorem 2(ii). Given $\varepsilon > 0$, choose a sufficiently small real number $\varepsilon_1 > 0$ and then a sufficiently large integer r_0 . Take any integers r, R such that $R \geq r_0$ and $R \leq \sqrt{18r \ln r}$.

Case 1. Suppose that $R \geq \ln r$.

We define

$$\delta := \max \left\{ \varepsilon_1, \frac{18R^2}{r} \right\}, \quad k := \left\lceil \frac{Rr}{R+\delta} \right\rceil + R, \quad \text{and} \quad c := R \ln \left(\frac{3R}{\delta} \right) + \ln(2R^3).$$

Clearly, $c \leq \binom{k}{R}$. Since R is large (which can be ensured by choosing r_0 sufficiently large) and $\ln r \leq R$, it follows from Theorem 2 (i) that

$$\mu(r - k + R, r - k) \leq 2R \ln \left(\frac{r - k + R}{R} \right) \leq 2R^2 \ln r \leq 2R^3.$$

Combining this with Lemma 3, Facts 4 and 5, we obtain

$$\begin{aligned} \mu(r + R, r) &\leq \binom{r + R}{R} \left(\frac{c}{\binom{k}{R}} + \frac{2R^3}{e^c \cdot \binom{r-k+R}{R}} \right) \\ &\leq \left(\frac{r}{k - R} \right)^R \cdot c + \left(\frac{r}{r - k} \right)^R \cdot \frac{2R^3}{e^c} \\ &\leq \left(1 + \frac{\delta}{R} \right)^R \cdot c + \left(\frac{3R}{\delta} \right)^R \cdot \frac{2R^3}{e^c} \leq e^\delta \cdot c + e^{R \ln \left(\frac{3R}{\delta} \right) - c} \cdot 2R^3 \leq e^\delta \cdot c + 1. \end{aligned}$$

If $\frac{18R^2}{r} \leq \varepsilon_1$, that is, $R \leq \sqrt{\varepsilon_1 r / 18}$, then

$$\begin{aligned} \mu(r + R, r) &\leq e^{\varepsilon_1} \cdot c + 1 \leq (1 + 2\varepsilon_1) \left(R \ln \left(\frac{3R}{\varepsilon_1} \right) + \ln(2R^3) \right) + 1 \\ &= (1 + 2\varepsilon_1) \left(R \left(\ln R + \ln \left(\frac{3}{\varepsilon_1} \right) \right) + 3 \ln R + \ln 2 \right) + 1 \\ &\leq (1 + \varepsilon) R \ln R, \end{aligned}$$

as desired. Note that the last inequality holds since $R \geq r_0$ (and we choose r_0 sufficiently large depending on ε_1) and $\varepsilon_1 \ll \varepsilon$.

If $\frac{18R^2}{r} > \varepsilon_1$, that is, $r < \frac{18R^2}{\varepsilon_1}$, then

$$\begin{aligned}\mu(r + R, r) &\leq e^{18R^2/r} \cdot c + 1 \leq e^{18R^2/r} \left(R \ln \left(\frac{r}{6R} \right) + \ln(2R^3) \right) + 1 \\ &\leq e^{18R^2/r} \left(R \ln \left(\frac{3R}{\varepsilon_1} \right) + \ln(2R^3) \right) + 1 \\ &\leq e^{18R^2/r} (1 + \varepsilon) R \ln R,\end{aligned}$$

also as desired. Note that, as in the previous calculation, the last inequality holds since $R \geq r_0 \gg 1/\varepsilon_1$ and $\varepsilon_1 \ll \varepsilon$.

Case 2. Suppose that $R < \ln r$.

Let $r_1 := r$ and, inductively for $i = 1, 2, \dots$, define

$$k_i := \left\lceil \frac{Rr_i}{R + \varepsilon_1} \right\rceil + R, \quad \text{and} \quad r_{i+1} := r_i - k_i;$$

if $r_{i+1} < 18R^2/\varepsilon_1$ then let $t := i$ and stop. Since r_i decreases each time, this process terminates.

We prove by backward induction on $i \in [t]$ that

$$\mu(r_i + R, r_i) \leq (1 + \varepsilon) R \ln R. \quad (8)$$

First, consider the base case $i = t$. Note that, for $i \leq t$, we have by $r_i \geq 18R^2/\varepsilon_1$ that

$$r_{i+1} = r_i - k_i \geq r_i - \frac{Rr_i}{R + \varepsilon_1} - 1 - R = \frac{\varepsilon_1 r_i}{R + \varepsilon_1} - (R + 1) \geq \frac{\varepsilon_1 r_i}{2(R + \varepsilon_1)}.$$

In particular this holds for $i = t$, giving that $r_t \leq \frac{18R^2}{\varepsilon_1} \cdot \frac{2(R + \varepsilon_1)}{\varepsilon_1}$, which is clearly at most e^R . Thus $R \geq \ln r_t$ and the desired conclusion follows by Case 1.

Now consider the inductive step for some $i \in [t-1]$. Let

$$c := R \ln(3R/\varepsilon_1) + \ln(2R \ln R) \leq \binom{k_i}{R}.$$

It follows from Lemma 3, Facts 4 and 5 (note that $\varepsilon_1 \geq 18R^2/r_t \geq 18R^2/r_i$), and the inductive hypothesis that

$$\begin{aligned}\mu(r_i + R, r_i) &\leq \binom{r_i + R}{R} \left(\frac{c}{\binom{k_i}{R}} + \frac{\mu(r_{i+1} + R, r_{i+1})}{e^c \cdot \binom{r_{i+1} + R}{R}} \right) \\ &\leq \left(\frac{r_i}{k_i - R} \right)^R \cdot c + \left(\frac{r_i}{r_{i+1}} \right)^R \cdot \frac{(1 + \varepsilon) R \ln R}{e^c} \\ &\leq \left(1 + \frac{\varepsilon_1}{R} \right)^R \cdot c + \left(\frac{3R}{\varepsilon_1} \right)^R \cdot \frac{(1 + \varepsilon) R \ln R}{e^c}\end{aligned}$$

$$\begin{aligned}
&\leq e^{\varepsilon_1} \cdot c + e^{R \ln \left(\frac{3R}{\varepsilon_1} \right) - c} \cdot 2R \ln R \\
&\leq (1 + 2\varepsilon_1) \left(R \ln \left(\frac{3R}{\varepsilon_1} \right) + \ln(2R \ln R) \right) + 1 \leq (1 + \varepsilon) R \ln R,
\end{aligned}$$

as desired. \square

This completes the proof of Theorem 2. \square

Remark 6. We did not optimise the bound in Theorem 2(ii) when $R = \Omega(\sqrt{r})$, since our main aim was to extend the inequality $\mu(r + R, r) \leq (1 + o(1))R \ln R$ for constant R from [16] to as large as possible range of functions $R(r)$.

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