

Note on a Conjecture of Talagrand: Expectation Thresholds vs. Fractional Expectation Thresholds

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Abstract

We show that a restricted version of a conjecture of M. Talagrand on the relation between “expectation thresholds” and “fractional expectation thresholds” follows easily from a strong version of a second conjecture of Talagrand, on “selector processes”. The selector process conjecture was proved by Park and Pham, and the quantitative strengthening used here is due to Bednorz, Martyněk, and Meller.

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1 Introduction

The purpose of this note is to point out that a restricted version of a conjecture of M. Talagrand (Conjecture 1 below) is an easy consequence of [1], a quantitative strengthening of [5, Theorem 1.5]. We first briefly recall definitions. (For a less hurried introduction to the problem, see e.g. [2, 4].)

For a finite set V , 2^V is the power set of V . An $\mathcal{F} \subseteq 2^V$ is *increasing* if $B \supseteq A \in \mathcal{F} \Rightarrow B \in \mathcal{F}$. For $\mathcal{G} \subseteq 2^V$ we use $\langle \mathcal{G} \rangle$ for the increasing family generated by $\mathcal{G}$, namely $\{B \subseteq V : \exists A \in \mathcal{G}, B \supseteq A\}$. We assume throughout that $|V| = n$ and $\mathcal{F} \subseteq 2^V$ is increasing.

Say $\mathcal{F}$ is *p-small* if there is a $\mathcal{G} \subseteq 2^V$ such that

$$\langle \mathcal{G} \rangle \supseteq \mathcal{F} \text{ and } \sum_{S \in \mathcal{G}} p^{|S|} \leq 1/2,$$

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and set $q(\mathcal{F}) = \max\{p : \mathcal{F} \text{ is } p\text{-small}\}$. Say $\mathcal{F}$ is *weakly p -small* if there is a $\lambda : 2^V \rightarrow \mathbb{R}^+ (:= [0, \infty))$ such that

$$\sum_{S \subseteq I} \lambda_S \geq 1 \quad \forall I \in \mathcal{F} \quad \text{and} \quad \sum_S \lambda_S p^{|S|} \leq 1/2, \quad (1)$$

and set $q_f(\mathcal{F}) = \max\{p : \mathcal{F} \text{ is weakly } p\text{-small}\}$. If there is λ as in (1) supported on sets of size at most r , we say $\mathcal{F}$ is *weakly (p, r) -small*.

We are interested in the following conjecture of Talagrand [6, Conjecture 6.3].

Conjecture 1. There is a universal $L > 0$ such that for every finite V and increasing $\mathcal{F} \subseteq 2^V$,

$$q(\mathcal{F}) \geq q_f(\mathcal{F})/L.$$

That is, weakly p -small implies (p/L) -small.

Our goal here is the following statement, which says that for any r , weakly (p, r) -small does imply (p/L) -small, but with L depending on r . The case $r = 2$, which was proved in [4], was one of a pair of test cases for Conjecture 1 suggested in [6]. (The other was resolved in [2].) The simplicity of the present argument versus the rather difficult [4] is a nice illustration of the power of the selector process results we are using.

Theorem 2. For any r , there is a J such that any weakly (Jp, r) -small $\mathcal{F}$ is p -small.

This turns out to be an easy consequence of Lemma 3 below, which is essentially [1, Lemma 3.4]. For the statement of the lemma, we need a little more preparation. Let $\mathcal{F} \subseteq 2^V$ be given, and for each $I \in \mathcal{F}$, suppose we are given weights $(\mu_I(v))_{v \in I}$ with $\mu_I(v) \geq 0$ and

$$\mu_I(I) := \sum_{v \in I} \mu_I(v) = 1.$$

Say $X \subseteq V$ is c -bad ($c > 0$) if

$$\sup_{I \in \mathcal{F}} \mu_I(X \cap I) < c.$$

Lemma 3. For each positive integer r there is a C such that if $\mathcal{F}$ is not p -small and W is chosen uniformly from the m -element subsets of V , then

$$\mathbb{P} \left(W \text{ is } \left(1 - \frac{1}{2r} \right)\text{-bad} \right) < 2 \sum_{t=1}^n \left(C \frac{np}{m} \right)^t.$$

When $r = 1$ this follows from [1, Lemma 3.4], and direct extension of the argument of [1] gives Lemma 3 with $C = (2er)^2$.

2 Proof of Theorem 2

We will prove Theorem 2 with $J = 10rC$, where $C = C(r)$ is as in Lemma 3.¹ For the rest of our discussion, X is a uniformly random m -element subset of V , with $m = Jpn/(2r) = 5Cpn$.

Suppose $\lambda : \binom{V}{\leq r} \rightarrow \mathbb{R}^+$ satisfies (1) with Jp in place of p . We may assume $\lambda_\emptyset = 0$, since otherwise λ' given by $\lambda'_\emptyset = 0$ and $\lambda'_S = \lambda_S/(1 - \lambda_\emptyset)$ if $S \neq \emptyset$ also satisfies (1) (with $1/2$ improved to $(1 - 2\lambda_\emptyset)/(2(1 - \lambda_\emptyset))$). We use Lemma 3 in combination with the following easy point.

Observation 4. *With probability at least $1/2$, $\sum_{S \subseteq X} \lambda_S < 1/(2r)$.*

Proof. This follows from Markov's Inequality, once we notice that our assumptions on X and λ imply

$$\mathbb{E} \left[\sum_{S \subseteq X} \lambda_S \right] \leq \sum_S (m/n)^{|S|} \lambda_S = \sum_S (Jp/(2r))^{|S|} \lambda_S \leq (2r)^{-1} \sum_S (Jp)^{|S|} \lambda_S \leq 1/(4r). \quad \square$$

For $I \in \mathcal{F}$, let $\nu_I = \sum_{S \subseteq I} |S| \lambda_S \geq 1$ (see (1), recalling that $\lambda_\emptyset = 0$). Define $\mu_I : I \rightarrow \mathbb{R}^+$ by

$$\mu_I(v) = \nu_I^{-1} \sum_{v \in S \subseteq I} \lambda_S \quad \forall v \in I$$

and notice that $\mu_I(I) = 1$. Suppose for a contradiction that $\mathcal{F}$ is not p -small. Then by Lemma 3 and our choice of m , the probability that X is $(1 - 1/(2r))$ -bad (with respect to the μ_I 's) is less than $1/2$. So the following claim gives the desired contradiction.

Claim 5. *With probability at least $1/2$, X is $(1 - 1/(2r))$ -bad with respect to the μ_I 's.*

Proof. For any $Y \subseteq V$ and $I \in \mathcal{F}$,

$$\begin{aligned} \mu_I(Y \cap I) &= \nu_I^{-1} \sum_{v \in Y \cap I} \sum_{v \in S \subseteq I} \lambda_S \leq \nu_I^{-1} \sum_{\substack{S \subseteq I \\ S \not\subseteq Y}} (|S| - 1) \lambda_S + \nu_I^{-1} \sum_{S \subseteq Y \cap I} |S| \lambda_S \\ &= \nu_I^{-1} \sum_{S \subseteq I} (|S| - 1) \lambda_S + \nu_I^{-1} \sum_{S \subseteq Y \cap I} \lambda_S \\ &\leq \nu_I^{-1} (1 - 1/r) \sum_{S \subseteq I} |S| \lambda_S + \nu_I^{-1} \sum_{S \subseteq Y \cap I} \lambda_S \\ &= 1 - 1/r + \nu_I^{-1} \sum_{S \subseteq Y \cap I} \lambda_S \leq 1 - 1/r + \sum_{S \subseteq Y \cap I} \lambda_S. \end{aligned}$$

When $Y = X$ we may bound the last sum (for any I) by $\sum_{S \subseteq X} \lambda_S$, which by Observation 4 is less than $(2r)^{-1}$ with probability at least $1/2$; and it follows that with probability at least $1/2$,

$$\sup_{I \in \mathcal{F}} \mu_I(X \cap I) < 1 - 1/r + (2r)^{-1} = 1 - 1/(2r). \quad \square$$

¹A more involved argument, using ideas from the proof of [3, Theorem 8], would allow us to take $J = O(C)$, but for simplicity we settle for the present version.

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