

ON PRIMITIVE 3-SMOOTH PARTITIONS OF n

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This paper is dedicated to Paul Erdős, on the sad occasion of his recent death.

ABSTRACT. A primitive 3-smooth partition of n is a representation of n as the sum of numbers of the form $2^a 3^b$, where no summand divides another. Partial results are obtained in the problem of determining the maximal and average order of the number of such representations. Results are also obtained regarding the size of the terms in such a representation, resolving questions of Erdős and Selfridge.

0. INTRODUCTION

Recently Erdős proposed the following problem: let $r(n)$ be the number of representations of n as the sum of 3-smooth numbers (integers of the form $2^a 3^b$ with $a, b \geq 0$), which are primitive (no summand divides another). Determine

- i. the maximal order of $r(n)$,
- ii. the average order of $r(n)$.

It is easy to show that $r(n) \geq 1$ for all n (see[1],[2]). In this paper partial results are obtained for both of these problems. Specifically, we prove:

Theorem 1. For $n \geq 5$,

$$r(n) \leq \frac{1}{2} \cdot n^\alpha,$$

where $\alpha = \log 2 / \log 3$ (≈ 0.631).

Theorem 2. Let $R(x) = \sum_{n \leq x} r(n)$. Then

$$\frac{x^\beta}{(\log x)^{\beta+3/2}} \ll R(x) \ll \frac{x^\beta}{(\log x)^{3/2}}$$

where

$$\beta = \frac{1}{\log 3} \cdot \log\left(\frac{\log 6}{\log 2}\right) + \frac{1}{\log 2} \cdot \log\left(\frac{\log 6}{\log 3}\right) (\approx 1.570).$$

Define $g(n)$ as the maximum, over all representations, of the minimum term (e.g. $11=8+3=9+2$, so $g(11) = 3$). Erdős has asked if $\lim_{n \rightarrow \infty} g(n) = \infty$. We answer this in the affirmative by proving:

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Theorem 3. $g(n) \geq \frac{3}{5} \cdot n^\gamma$, where $\gamma = \frac{\log 2}{\log 5} (\approx 0.431)$.

Concerning $g(n)$, Selfridge has asked if, for each n , $g(n)$ appears as the least term in exactly one representation of n . We answer this in the negative by providing an infinite number of counter-examples where it appears in two representations.

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1. PROOF OF THEOREM 1

Lemma 1. $r(2^a 3^b m) = r(m)$

Proof. It is enough to show that $r(2m) = r(m)$ and $r(3m) = r(m)$. There cannot exist a term 3^β in any representation of $2m$, for that would make the sum odd. Thus all terms are even. Dividing through by 2 gives a 1-1 correspondence between representations of $2m$ and m . Likewise, for $3m$, a term 2^α would mean the sum is not divisible by 3. Thus all terms are divisible by 3. Dividing by 3 gives a similar correspondence between $3m$ and m .

The above lemma tells us that we need only deal with n coprime to 6 to determine an upper bound. Each representation of a number coprime to 6 must have a term 3^j . Define

$$k = \lfloor \log_3 n \rfloor$$

and for these n , define $r_j(n)$ = the number of representations with summand 3^j , for $j \geq 0$. Note that

$$r(n) = \sum_{j=0}^k r_j(n).$$

Also define

$$s_j(n) = \sum_{i=0}^j r_i(n),$$

(so that $r(n) = s_j(n)$ for $j \geq k$).

Lemma 2. For $n \geq 5$, and $i = 1, 2, \dots, k$

$$s_i(n) \leq 2^{i-1}.$$

[Note that $s_0(n) = r_0(n) = 0$ if $n \neq 1$, and is 1 if $n = 1$.]

Proof. We use complete induction on n . For $n=5$ or 7 , we have $k = 1$. Since $5=3+2$, and $7=3+4$ are their only representations,

$$s_1(n) = r_1(n) = 1 = 2^0$$

so the inequalities hold. Note that for $n > 1$, $s_1(n) = r_1(n)$ = the number of representations of n in the form $3 + 2^a$. Obviously, then

$$s_1(n) = 0 \text{ or } 1 \leq 2^0.$$

Now assume for all $m < n$ and $i = 1, 2, \dots, k = k(m)$ that

$$s_i(m) \leq 2^{i-1}.$$

It is easy to see that, for $j \geq 2$,

$$r_j(n) = s_{j-1} \left(\frac{n - 3^j}{2^{a_j}} \right)$$

where a_j is defined by $2^{a_j} \parallel n - 3^j$. Therefore, from the inductive hypothesis, we deduce that for $j \geq 2$

$$r_j(n) \leq 2^{j-2}.$$

Thus, for $n \geq 5$ and $i \geq 2$

$$s_i(n) = \sum_{j=1}^i r_j(n) \leq 1 + \sum_{j=2}^i 2^{j-2} = 2^{i-1}.$$

This completes the proof of the lemma.

The above lemma yields

$$r(n) = s_k(n) \leq 2^{k-1} \leq \frac{1}{2} \cdot 2^{\log_3(n)} = \frac{1}{2} \cdot n^\alpha$$

which completes the proof of Theorem 1.

2. PROOF OF THEOREM 2

We introduce the following notation:

- $\sum^*$ indicates a sum over numbers prime to 6,
- $R^*(x) = \sum_{n \leq x}^* r(n)$,
- $S(x)$ = the number of primitive sums $\sum 2^a 3^b = n$ with each term $2^a 3^b \leq x$ and $(n, 6) = 1$,
- $S_j(x)$ = the number of sums described above that have the summand 3^j ,
- $\alpha = \frac{\log 2}{\log 3}$.

For any $x \geq 3$, there exists $L \in \mathbb{N}$ such that $3^L \leq x < 3^{L+1}$. Obviously, if we prove the theorem for $x = 3^L$, the general result will follow from $R(3^L) \leq R(x) \leq R(3^{L+1})$. Therefore, we henceforth assume that $x = 3^L$.

Note that every $n \leq x$ can be uniquely written as $n = 2^a 3^b m$ with $(m, 6) = 1$, and $m \leq x/2^a 3^b$. Since $r(n) = r(m)$, by lemma 1, it follows that

$$R(x) = \sum_{\substack{a, b \geq 0 \\ 2^a 3^b \leq x}} R^* \left(\frac{x}{2^a 3^b} \right).$$

Using this fact, the theorem will follow once we prove, for $x \geq 2$,

$$(2.1) \quad \frac{x^\beta}{(\log x)^{\beta+3/2}} \ll R^*(x) \ll \frac{x^\beta}{(\log x)^{3/2}}.$$

Specifically, the theorem's lower bound follows trivially from this, and on the other hand we would have

$$\begin{aligned} R(x) &\ll \sum_{2^a 3^b \leq \sqrt{x}} \left(\frac{x}{2^a 3^b} \right)^\beta \cdot (\log x)^{-3/2} + \sum_{\sqrt{x} < 2^a \cdot 3^b \leq x} x^{\beta/2} \\ &\ll x^\beta \cdot (\log x)^{-3/2} + \log^2 x \cdot x^{\beta/2} \\ &\ll x^\beta \cdot (\log x)^{-3/2}. \end{aligned}$$

Consider $S(x/L)$. This counts primitive sums $\sum 2^a 3^b$ with $2^a 3^b \leq x/L$. Then $b \leq L-1$, and we have at most L terms. Thus, it follows from the definitions that

$$S\left(\frac{x}{L}\right) \leq R^*(x) \leq S(x).$$

This will imply (2.1) once we prove:

$$S(x) \asymp (\log x)^{-3/2} \cdot x^\beta,$$

and this would follow from

$$(2.2) \quad S(x) \asymp \frac{1}{L} \cdot \binom{[L/\alpha] + L}{L},$$

since, by Stirling's formula, the binomial coefficient above equals

$$\exp\{([L/\alpha] + L + 1/2) \log([L/\alpha] + L) - (L + 1/2) \log L - ([L/\alpha] + 1/2) \log([L/\alpha]) + O(1)\}$$

$$= \exp\{[L/\alpha] \cdot \log(1 + \alpha) + L \cdot \log(1 + 1/\alpha) - 1/2 \cdot \log L + O(1)\}$$

$$\asymp x^\beta \cdot (\log x)^{-1/2}.$$

It is easy to see that

$$S\left(\frac{x}{3}\right) \leq S_L(x) < S(x).$$

The second inequality trivially follows from the definitions. Since $x = 3^L$, no sum counted by $S(x/3)$ has a term divisible by 3^L . Thus, if we take any such sum, multiply through by 2 and add 3^L , we obtain a sum counted by $S_L(x)$. Hence, the first inequality holds, and (2.2) would follow from

$$(2.3) \quad S_L(x) \asymp \frac{1}{L} \cdot \binom{[L/\alpha] + L}{L}.$$

We now establish a recursive formula for the S_j and use it to prove (2.3).

The sums counted by $S_j(x)$ can all be written in the form

$$3^j + 2^l \cdot (3^i + \dots)$$

where $0 \leq i \leq j-1$, and $1 \leq l \leq [(L-i)/\alpha]$. If $0 \leq i \leq j-2$, replacing 3^j by 3^{j-1} yields a 1-1 correspondence with precisely those sums counted by $S_{j-1}(x)$. If $i = j-1$, the number of different sums in the parentheses, for a given l , is $S_{j-1}(x/2^l)$. Therefore

$$(2.4) \quad S_j(x) = \sum_{l=0}^{[(L-j+1)/\alpha]} S_{j-1}\left(\frac{x}{2^l}\right).$$

In particular

$$S_L(x) = S_{L-1}(x) + S_{L-1}\left(\frac{x}{2}\right).$$

We can generalize this to a relationship of the form

$$(2.5) \quad S_L(x) = \sum_{l=0}^{[j/\alpha]} A_j(l) \cdot S_{L-j}\left(\frac{x}{2^l}\right)$$

via (2.4), where $A_1(0) = A_1(1) = 1$ by the above. Note that the sums counted by $S_{L-j-1}(x/2^{l+\lambda})$ have $3^{L-j-1} \leq 3^L/2^{l+\lambda}$, which implies $\lambda \leq [(j+1)/\alpha] - l$. Now if we combine the above with (2.4) we obtain

$$\begin{aligned} S_L(x) &= \sum_{l=0}^{[j/\alpha]} A_j(l) \sum_{\lambda=0}^{[\frac{j+1}{\alpha}] - l} S_{L-j-1}\left(\frac{x}{2^{l+\lambda}}\right) \\ &= \sum_{m=0}^{[\frac{j+1}{\alpha}]} S_{L-j-1}\left(\frac{x}{2^m}\right) \sum_{l=0}^{\min(m, [\frac{j}{\alpha}])} A_j(l), \end{aligned}$$

so we conclude

$$(2.6) \quad A_{j+1}(m) = \sum_{l=0}^{\min(m, [j/\alpha])} A_j(l).$$

It is trivially true that $S_0(x) = 1$ for $x \geq 1$, so with $j = L$, (2.5) yields

$$(2.7) \quad S_L(x) = \sum_{l=0}^{[L/\alpha]} A_L(l).$$

Now let $B_j(l) = \binom{l+j}{l} - \alpha \cdot \binom{l+j}{l-1}$ (with the conventions that $\binom{0}{0} = 1$ and $\binom{n}{-1} = 0$).

Claim: $B_{j-1}(l) \leq A_j(l) \leq B_{j+1}(l)$.

Once this claim is proved, (2.3) will follow, and the proof will be complete. Specifically, since $\sum_{l=0}^m \binom{l+j}{l} = \binom{m+j+1}{m}$,

$$(2.8) \quad \sum_{l=0}^m B_j(l) = B_{j+1}(m).$$

Therefore, combining this with (2.7) and the claim,

$$B_L([L/\alpha]) \leq S_L(x) \leq B_{L+2}([L/\alpha]).$$

This is the same as

$$\binom{[L/\alpha] + L}{[L/\alpha]} \cdot \left(1 - \alpha \cdot \frac{[L/\alpha]}{L+1}\right) \leq S_L(x) \leq \binom{[L/\alpha] + L + 2}{[L/\alpha]} \cdot \left(1 - \alpha \cdot \frac{[L/\alpha]}{L+3}\right)$$

and (2.3) follows at once.

Proof of claim: We use induction on j . For $j = 1$ the inequalities follow immediately from the fact that $A_1(0) = A_1(1) = 1$.

Now assume the inequalities hold for a given $j \geq 1$ and $0 \leq l \leq [j/\alpha]$. Suppose first that $0 \leq m \leq [j/\alpha]$. Summing the inequalities over $l \in [0, m]$, applying (2.6) and (2.8) yields the corresponding inequalities for $A_{j+1}(m)$.

We are left to handle the $[j/\alpha] < m \leq [(j+1)/\alpha]$, so that $m = [j/\alpha] + i$, where $i = 1$ or possibly 2. For these m , (2.6) tells us that $A_{j+1}(m) = A_{j+1}([j/\alpha])$. The already established inequalities from the preceding paragraph yield

$$B_j([j/\alpha]) \leq A_{j+1}(m) \leq B_{j+2}([j/\alpha]).$$

The proof of the claim will be complete if we show

(i) $B_{j+2}([j/\alpha]) < B_{j+2}([j/\alpha] + 1) < B_{j+2}([j/\alpha] + 2)$
and

(ii) $B_j([j/\alpha]) > B_j([j/\alpha] + 1) > B_j([j/\alpha] + 2)$.

Towards this, note that for $i = 1$ or 2 , and $j \in \mathbb{N}$,

$$\alpha \cdot \frac{1}{j+2} < \frac{1}{[j/\alpha] + i} < \alpha \cdot \frac{1}{j}.$$

These two inequalities respectively imply

$$B_{j+1}([j/\alpha] + i) > 0$$

and

$$B_{j-1}([j/\alpha] + i) < 0.$$

Noting that $\binom{n}{k} = \binom{n+1}{k+1} - \binom{n}{k+1}$, the definition of $B_j(l)$ yields

$$B_j(l) = B_j(l+1) - B_{j-1}(l+1).$$

Combined with the above inequalities, this establishes the validity of (i) and (ii), hence the claim, and hence we have theorem 2.

3. PROOF OF THEOREM 3

Since $g(2^a 3^b n) = 2^a 3^b \cdot g(n)$ (from the proof of lemma 1), we may assume $(n, 6) = 1$. It is easy to check that the theorem holds for $n=1,5,7$. We proceed inductively. We may now assume

$$(3.1) \quad 3^L < n < 3^{L+1}, \quad L \geq 2$$

so that

$$(3.2) \quad n = 3^L + 2^a k, \quad a \geq 1, k \equiv 1(2), k < 3^L.$$

Evidently

$$g(n) \geq \min(3^L, 2^a \cdot g(k)).$$

If $g(n) \geq 3^L$, then $g(n) > n/3$ and the theorem holds. Otherwise

$$g(n) \geq 2^a \cdot \frac{3}{5} \cdot k^\gamma$$

by the inductive hypothesis.

Case I: $k \geq 3^{L-a}$

We wish to show that

$$2^a \cdot k^\gamma \geq n^\gamma$$

or

$$k \geq \frac{n}{2^{a/\gamma}} = \frac{3^L + 2^a \cdot k}{2^{a/\gamma}}$$

or

$$k \geq \frac{3^L}{2^{a/\gamma} - 2^a}.$$

Thus it is enough to have

$$3^{L-a} \geq \frac{3^L}{2^{a/\gamma} - 2^a}$$

or

$$\gamma \leq \min_{a \geq 1} \frac{a \cdot \log 2}{\log(3^a + 2^a)} = \frac{\log 2}{\log 5},$$

so the theorem holds. Note that if $a \geq L$, we certainly have $k \geq 3^{L-a}$, so that in particular the theorem holds in that situation.

Case II: $k < 3^{L-a}$ and $1 \leq a \leq L-1$

From (3.2):

$$\begin{aligned} n &= (3^L - 3^L \cdot (2/3)^a) + (2^a \cdot 3^{L-a} + 2^a \cdot k) \\ &= (3^{L-1} + 2 \cdot 3^{L-2} + 2^2 \cdot 3^{L-3} + \dots + 2^{a-1} \cdot 3^{L-a}) + 2^{a+1} \left(\frac{3^{L-a} + k}{2} \right). \end{aligned}$$

Now either

$$g \left(2^{a+1} \left(\frac{3^{L-a} + k}{2} \right) \right) \geq 2^{a-1} \cdot 3^{L-a}$$

or

$$g \left(2^{a+1} \left(\frac{3^{L-a} + k}{2} \right) \right) < 2^{a-1} \cdot 3^{L-a}.$$

In the former circumstance

$$\begin{aligned} g(n) &\geq 2^{a-1} \cdot 3^{L-a} \\ &\geq \frac{3}{4} \cdot 2^L \\ &= \frac{3}{4} \cdot (3^L)^{\log 2 / \log 3} \\ &> \frac{3}{8} \cdot n^{\log 2 / \log 3} \\ &\geq \frac{3}{5} \cdot n^{\log 2 / \log 5} \end{aligned}$$

(using (3.1) and $n \geq 11$ in the last two lines respectively). In the latter circumstance

$$g(n) \geq 2^{a+1} \cdot \frac{3}{5} \left(\frac{3^{L-a} + k}{2} \right)^\gamma$$

by the inductive hypothesis. To complete the proof, it is enough to show that

$$2^{a+1} \geq \left(\frac{2 \cdot (3^L + 2^a \cdot k)}{3^{L-a} + k} \right)^\gamma$$

(which combined with the above and (3.2) yields the result).

The right side is maximized at $k = 1$, so it is

$$\begin{aligned} &\leq \left(\frac{2(3^L + 2^a)}{3^{L-a} + 1} \right)^\gamma \\ &< (2 \cdot 3^a)^\gamma \\ &< 2^{a+1}. \end{aligned}$$

4. NON-UNIQUENESS OF REPRESENTATION WITH $g(n)$

Lastly, we show the counter-examples to Selfridge's question. Let

$$n = 3^b + 11 \cdot 2^a$$

with $a \geq 5$, and

$$(4.1) \quad 11 \cdot 2^{a-3} < 3^b < 99 \cdot 2^{a-6}.$$

For any given a , there are 0 or 1 values of b satisfying (4.1). However, for an infinite number of values of a , there exists a value of b satisfying (4.1). This is because it is equivalent to

$$\frac{\log 11}{\log 3} + (a-3) \cdot \frac{\log 2}{\log 3} < b < \frac{\log 99}{\log 3} + (a-6) \cdot \frac{\log 2}{\log 3},$$

or

$$\frac{\log(64/33)}{\log 3} < a \cdot \frac{\log 2}{\log 3} - (b-1) < \frac{\log(24/11)}{\log 3}.$$

Since the number $\log 2 / \log 3$ is irrational, the fractional part of $a \cdot \log 2 / \log 3$ is thus dense in $[0,1]$ and so lies between $\log(64/33)/\log 3$ and $\log(24/11)/\log 3$ for infinitely many values of n .

We will show that $g(n) = 3^b$. This will suffice since

$$\begin{aligned} n &= 3^b + 3 \cdot 2^a + 2^{a+3} \\ &= 3^b + 9 \cdot 2^a + 2^{a+1} \end{aligned}$$

and (4.1) implies that 3^b is the smallest term in both representations.

We must now demonstrate that there does not exist a representation of n with all its terms $> 3^b$. Suppose that such a representation exists. Since n is odd, any representation

must have a term 3^j , and since (4.1) implies that $n < 3^{b+2}$, this representation must have the term 3^{b+1} . Thus it is in the form

$$n = 3^{b+1} + 2 \cdot n_1$$

where

$$n_1 = 11 \cdot 2^{a-1} - 3^b.$$

Since n_1 is odd, its corresponding representation has a term 3^j . Since $n < 3^{b+2}$, we have $n_1 < 3^{b+1}$, and thus $j \leq b$. If j were less than b , the corresponding term $2 \cdot 3^j$ in the representation of n would contradict our assumption that all the terms in that representation are greater than 3^b . Following this same line of reasoning, the alleged representation unfolds as follows:

$$\begin{aligned} n &= 3^{b+1} + 2 \cdot 3^b + 4(11 \cdot 2^{a-2} - 3^b) \\ &= 3^{b+1} + 2 \cdot 3^b + 4 \cdot 3^{b-1} + 16(11 \cdot 2^{a-4} - 3^{b-1}) \\ &= 3^{b+1} + 2 \cdot 3^b + 4 \cdot 3^{b-1} + 16 \cdot 3^{b-2} + 16(11 \cdot 2^{a-4} - 4 \cdot 3^{b-2}). \end{aligned}$$

However

$$\begin{aligned} m &=: 16(11 \cdot 2^{a-4} - 4 \cdot 3^{b-2}) \\ &= 11 \cdot 2^a - \frac{64}{9} \cdot 3^b \\ &< 8 \cdot 3^b - \frac{64}{9} \cdot 3^b \\ &= \frac{8}{9} \cdot 3^b, \end{aligned}$$

and clearly $m \neq 0$. This contradicts the assumption that all terms are $> 3^b$, and completes the proof.

REFERENCES

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