On the Dimer Problem and the Ising Problem in Finite 3-dimensional Lattices

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Abstract

We present a new expression for the partition function of the dimer arrangements and the Ising partition function of the 3-dimensional cubic lattice. We use the Pfaffian method. The partition functions are expressed by means of expectations of determinants and Pfaffians of matrices associated with the cubic lattice.

1 Introduction

The close-packed dimer model of statistical mechanics can be stated as follows. One considers a set of sites and a set of bonds connecting certain pairs of sites. Each bond b can absorb a 'dimer' (which represents a diatomic molecule) with corresponding energy E_b . It is required that each site is occupied exactly once by one of the atoms of a dimer. A state s is an arrangement of dimers which meets this requirement, and its energy E(s) is $\sum E_b$ where the sum is taken over all bonds b which absorb a dimer. Then the partition function of the dimer model may be viewed as a density function of energy levels.

The dimer model was first considered by Roberts [18] in 1935, and by Fowler and Rushbrook [3]. The dimer model for 2-dimensional lattices appears in calculations of the thermodynamic properties of a system of diatomic molecules-dimers. It has been solved

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by Kasteleyn [13] and by Temperley and Fisher [11]. The same problem for 3-dimensional lattices remains an important open problem of statistical physics (see [14] for references).

Many fundamental observations about the dimer and monomer-dimer model in general lattice graphs have been given by Heilmann and Lieb [9], [10].

Another model we consider here is the Ising version of the Edwards-Anderson model. It can be described as follows. A coupling constant J_{ij} is assigned to each bond $\{i, j\}$ of a given lattice graph G; the coupling constant characterizes the interaction between the particles represented by sites *i* and *j*. A physical state of the system is an assignment of spin $\sigma_i \in \{+1, -1\}$ to each site *i*. The Hamiltonian (or energy function) is defined as $H(\sigma) = -\sum_{\{i,j\}\in E} J_{ij}\sigma_i\sigma_j$. The distribution of physical states over all possible energy levels is encapsulated in the partition function $Z(\beta) = \sum_{\sigma} e^{-\beta H(\sigma)}$ from which all fundamental physical quantities may be derived.

The literature on the 3-dimensional dimer problem and the 3-dimensional Ising problem is vast but there is a general feeling and some evidence (see e.g. [12]) that no closed solution similar to the solutions of the 2-dimensional case nor a deterministic efficient algorithm may be found for the cubic lattices.

This however does not rule out a statistical treatment. We believe that our new expressions are natural enough to allow such further analysis.

Recent papers [16], [17] also study the problems using a Pfaffian method. They obtain new expressions by means of a series of Pfaffians with a topological signature. Our approach is more combinatorial in nature. We express the partition functions by means of expectations of the determinants of matrices naturally associated with the cubic lattice. Determinants and spectral properties of random matrices are extensively studied (see e.g. [8]) and a goal of this paper is to draw attention to possible applications of related machinery to the 3-dimensional statistical mechanics problems.

We may reformulate the dimer problem and the Ising problem in graph theoretic terms as follows. A graph is a pair G = (V, E) where V is a set of *vertices* and E is now the set of edges (not the energy). A graph with some regularity properties may be called a *lattice graph*. We associate with each edge e of G a weight w(e) and for a subset of edges $A \subset E$, w(A) will denote the sum of the weights w(e) associated with the edges in A.

A subset of edges $P \subset E$ is called a *perfect matching* or *dimer arrangement* if each vertex belongs to exactly one element of P. The dimer partition function may be viewed as a polynomial $\mathcal{P}(G, x)$ which equals the sum of $x^{w(P)}$ over all perfect matchings P of G. This polynomial is also called the *generating function of perfect matchings*.

The Ising partition function is very close to the generating function of cuts which is a standard concept in graph theory. A cut of a graph G = (V, E) is a partition of its vertices into two disjoint subsets $V_1, V_2 \subset V$, and the implied set of edges between the two parts:

$$C(V_1, V_2) = \{\{u, v\} \in E : u \in V_1, v \in V_2\}$$

The generating function of cuts $\mathcal{C}(G, x)$ equals the sum of $x^{w(C)}$ over all cuts C of G.

If we set the coupling constant J_{ij} as the weight $w(\{i, j\})$ of the edge $\{i, j\}$, the

generating function of cuts becomes very similar to the partition function:

$$Z(\beta) = 2\sum_{cutC} e^{-\beta(2w(C)-W)} = 2e^{\beta W} \mathcal{C}(G, e^{-2\beta})$$

where W is the sum of all the edge weights.

The generating functions of perfect matchings and cuts may be defined in a more general way as follows: associate a variable x_e with each edge e of graph G, let $x(A) = \prod_{e \in A} x_e$ and let e.g. the generating function of perfect matchings be the sum of x(P), P perfect matching of G. All results introduced in this paper also hold in this more general setting; however the presentation using weights rather than variables is perhaps more natural.

This paper studies properties of finite cubic lattices. Let us now fix some notation for them. Let m be an odd positive integer and k an even positive integer. The cubic lattice Q_{mmk} is the following graph:

 Q_{mmk} has vertices V_{xyz} , x, y = 1, ..., m, z = 1, ..., k, and the following edges:

- 1. The vertical edges $v_{xyz} = \{V_{xyz}, V_{xy(z+1)}\}, z = 1, ..., k 1,$
- 2. The width edges $w_{xyz} = \{V_{xyz}, V_{x(y+1)z}\},\ y = 1, ..., m-1,$
- 3. The horizontal edges $h_{xyz} = \{V_{xyz}, V_{(x+1)yz}\},\$

$$x = 1, ...m - 1$$

Let us denote the ordered set $(V_{xy1}, ..., V_{xyk})$ by V_{xy} . V_{xy} will also stand for the vertical path of Q_{mmk} from V_{xy1} to V_{xyk} . Let \overline{V}_{xy} denote the reversal of V_{xy} .

 Q_{mmk} is a bipartite graph, which means that its vertices may be partitioned into two sets Z_1, Z_2 such that if e is an edge of Q_{mmk} then $|e \cap Z_1| = |e \cap Z_2| = 1$. Moreover, we have also that $|Z_1| = |Z_2| = mmk/2$. Let \mathcal{Z} be the square $(Z_1 \times Z_2)$ matrix defined by $\mathcal{Z}_{ij} = x^{w(ij)}$ if $e = \{ij\}$ is an edge of Q_{mmk} with weight w(e) = w(ij), and $\mathcal{Z}_{ij} = 0$ otherwise.

We will consider matrix \mathcal{Z} with its rows and columns ordered in agreement with the *natural order* $(V_{11}, \bar{V}_{12}, ..., V_{1m}, \bar{V}_{21}, ..., V_{mm})$ and we will assume that $V_{111} \in Z_1$.

Note that $\mathcal{P}(Q_{mmk}, x)$ equals the permanent of \mathcal{Z} . In this paper we show that $\mathcal{P}(Q_{mmk}, x)$ may be computed from the average of determinants of CERTAIN signings of \mathcal{Z} , where a *signing* of a matrix is obtained by multiplying some of the entries of the matrix by -1.

The signings of \mathcal{Z} correspond to orientations of Q_{mmk} .

An orientation of a graph G = (V, E) is a digraph D = (V, A) obtained from G by assigning an orientation to each edge of G, i.e., by ordering the elements of each edge of G. The elements of A are called *arcs*. We say that signing Z of \mathcal{Z} corresponds to orientation D of Q_{mmk} if $Z_{ij} = x^{w(ij)}$ if $(ij) \in A(D)$, $Z_{ij} = -x^{w(ij)}$ if $(ji) \in A(D)$, and $Z_{ij} = 0$ otherwise.

An expression of similar flavor as our result exists already: a seminal observation of Heilmann and Lieb [9], [10] asserts that $\mathcal{P}(Q_{mmk}, x^2)$ equals the average of $(det(Z))^2$ over ALL signings Z of \mathcal{Z} .

The following short proof of this observation is taken from the monograph [15]. If D is an orientation of Q_{mmk} then let A(D) denote the skew-symmetric adjacency matrix of D, i.e. matrix consisting of 4 blocks where both blocks on the main diagonal are 0-matrices and the remaining two blocks equal Z and -Z, where Z is the signing of \mathcal{Z} corresponding to D. Clearly $det(A(D)) = (det(Z))^2$, hence we need to show that $\mathcal{P}(Q_{mmk}, x^2)$ equals the expectation of det(A(D)) over all orientations D of Q_{mmk} . For the expectation we have

$$E(det(A(D))) = \sum sgn(\pi)E(a_{1\pi(1)}...a_{n\pi(n)})$$

where n = mmk and $A(D) = (a_{ij})$ by the linearity of expectation. If π is a permutation having a fix point or such that i and $\pi(i)$ are non-adjacent for some $i \leq n$ then the term corresponding to π equals 0. If there is i such that $\pi(\pi(i)) \neq i$ then the random variable $a_{i\pi(i)}$ occurs in the term corresponding to π but the random variable $a_{\pi(i)i}$ does not. Hence

$$E(a_{1\pi(1)}...a_{n\pi(n)}) = E(a_{i\pi(i)})E(a_{1\pi(1)}...a_{(i-1)\pi(i-1)}a_{(i+1)\pi(i+1)}...a_{n\pi(n)}) = 0.$$

So we are left with the terms corresponding to those permutations which have no fix point, for which i and $\pi(i)$ are adjacent and $(\pi)^2$ is the identity. Such permutations uniquely correspond to perfect matchings of Q_{mmk} and the signs turn out correct. \Box

A difference between our expression and the result of Heilmann and Lieb is that we replace the average of a multi quadratic function by the average of a multi linear function, with a restricted range.

1.1 Statement of the main result.

An orientation D of Q_{mmk} is called *stable* if all vertical edges are oriented in D from the 'smaller' to the 'bigger' vertex in the natural order. For edge e we let $s_D(e) = 0$ if the orientation of e agrees with the natural order, and $s_D(e) = 1$ otherwise.

Theorem 1.1

$$\mathcal{P}(Q_{mmk}, x) = -2^{C_r} x^{w(M)} + \alpha (2^{C_r} + 1)$$

where $C_r = km(m-1)$, M is the unique perfect matching of Q_{mmk} consisting of vertical edges only and α equals the average of det(Z(D)), D stable orientation of Q_{mmk} satisfying

$$\sum_{A} s_D(w_{x,2y,z}) s_D(w_{x,2y'-1,z'}) = \sum_{B} s_D(h_{2x,y,z}) s_D(h_{2x'-1,y',z'})$$

modulo 2, where

$$A = \{(x, y, z, y', z'); 1 \le x \le m, 1 \le y \le (m-1)/2, 1 \le z \le k, 1 \le y' \le y, z \le z' \le z+1\}$$

and

$$B = \{(x, y, z, x', y', z'); 1 \le x \le (m - 1)/2, 1 \le z \le k, 1 \le y \le m, 1 \le x' \le x, (y, z) \le (y', z') \le (y'', z'')\}.$$

In the definition of B the order on pairs of integers is lexicographic order and (y'', z'') is the immediate successor of (y, z); if the immediate successor does not exist than we let (y'', z'') = (y, z).

Theorem 1.1 holds also for $\mathcal{P}(Q_{m_1m_2k}, x)$ with $m_1 \neq m_2$ odd. In this more general setting $C_r = 1/2km_1(m_2-1) + 1/2km_2(m_1-1)$,

$$A = \{(x, y, z, y', z'); 1 \le x \le m_1, 1 \le y \le (m_2 - 1)/2, 1 \le z \le k, 1 \le y' \le y, z \le z' \le z + 1\}$$

and

$$B = \{(x, y, z, x', y', z'); 1 \le x \le (m_1 - 1)/2, 1 \le z \le k, 1 \le y \le m_2, 1 \le x' \le x, (y, z) \le (y', z') \le (y'', z'')\}.$$

Example. Let us illustrate the statement of Theorem 1.1 by calculation of $\mathcal{P}(Q_{3,1,2}, x)$ with w(e) = 0 for each edge e. $Q_{3,1,2}$ has no width edges: it is simply a square (3×2) grid and thus it has 6 vertices and 3 dimer arrangements. Hence $\mathcal{P}(Q_{3,1,2}, x) = 3$.

On the other hand there are 2^4 stable orientations of $Q_{3,1,2}$ and those relevant for α are characterized by the equation

$$s_D(h_{2,1,1})s_D(h_{1,1,1}) + s_D(h_{2,1,2})s_D(h_{1,1,2}) + s_D(h_{2,1,1})s_D(h_{1,1,2}) = 0$$

modulo 2.

A simple calculation reveals that there are 10 such stable orientations and 6 stable orientations that are irrelevant. Hence α equals average of 10 determinants of signings of (3×3) matrix \mathcal{Z} . We can check by hand that $\alpha = 7/5$. Since $C_r = 2$, we have

$$-2^{C_r} x^{w(M)} + \alpha (2^{C_r} + 1) = -4 + 5(7/5) = 3.$$

If we want to use Theorem 1.1 to calculate $\mathcal{P}(Q_{3,3,2}, x)$ we would have $C_r = 12$ and α equal the average of $2^{23} + 2^{11}$ determinants of signings of a (9×9) matrix.

These huge numbers which appear even for very small lattices demonstrate the character of Theorem 1.1: it certainly does not aim to a computational efficiency.

Having the expression for the partition function of the dimer problem given by Theorem 1.1, let me briefly indicate how to transform the 3-dimensional Ising problem to the dimer problem of locally modified cubic lattice. This transformation goes back to Kasteleyn [13] and Fisher [2] and it is well described e.g. in [6]. An *eulerian subgraph* of a graph G = (V, E) is a set of edges $U \subset E$ such that each vertex of V is incident with an even number of edges from U. The generating function of eulerian subgraphs $\mathcal{E}(G, x)$ equals the sum of $x^{w(U)}$ over all eulerian subgraphs U of G. The partition function of the Ising problem of a graph (with zero magnetic field) can be expressed as the generating function of eulerian subgraphs of the same graph, with modified edge weights. This classic relation between the Ising partition function and the generating function of eulerian subgraphs was discovered by van der Waerden:

$$Z(\beta) = 2^n \prod_{\{i,j\}\in E} \cosh(\beta J_{ij}) \quad \mathcal{E}(G, \tanh(\beta J_{ij})),$$

see [20]. Hence it remains to transform the generating function of eulerian subgraphs of the cubic lattice Q_{mmk} into the generating function of perfect matchings of a locally modified graph Q^*_{mmk} . We use Fisher's construction [2] since it is local in the sense that it only modifies each vertex in a way dependent on its degree and it may be performed so that the embedding of Q_{mmk} is preserved. Fisher's construction may be described as follows:

Let G = (V, E) be a graph embedded in an orientable surface of genus g, and $v \in V$ a vertex. Let $e_1, e_2, ..., e_d \in E$ denote the edges incident with v, ordered clockwise as they spread out from v in the embedding. Then the even splitting of v is a graph G' = (V', E') where

- $V' = V \setminus \{v\} \cup \{v_1, ..., v_d, v'_1, ..., v'_d\}$
- $E' = E \setminus \{e_1, e_2, ..., e_d\} \cup \{e'_1, e'_2, ..., e'_d\} \cup E^A$

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$$E^A = \{\{v_i, v'_i\}; i = 1, ..., d\} \cup \{\{v_i, v'_{i-1}\}; i = 2, ..., d\} \cup \{\{v'_i, v'_{i+1}\}; i = 1, ..., d-1\}$$

The edges $e'_i \in E'$ (*image edges*) are obtained from $e_i \in E$ by replacing the vertex v by v_i . The edges E^A will be called *auxiliary*.

Note that the graph obtained by even splitting can be again embedded in the same surface since the transformation replaces a vertex $v \in V$ by a cluster of 2d vertices and 3d - 2 edges. The cluster itself is a planar graph which can be embedded in a small neighborhood of the original location of the vertex v. The images of the edges incident with v can be embedded in the same way as they were in the original graph.

Let G = (V, E) be a graph and $G^* = (V^*, E^*)$ the graph obtained by successive even splitting of all vertices in V. If there are weights w(e) assigned to edges $e \in E$, we assign the same weights to their images in E^* : w(e') = w(e). The auxiliary edges $f \in E^*$ get assigned w(f) = 0. With this assignment of weights, the generating function of perfect matchings of G^* is equal to the generating function of eulerian subgraphs of G,

$$\mathcal{P}(G^*, x) = \mathcal{E}(G, x).$$

This may be observed as follows: if M is a perfect matching in G^* , it must cover each of its vertices exactly once. Because the cluster replacing every vertex has an even number of vertices, and any of the auxiliary edges which is in M covers a pair of vertices of the cluster, there remain an even number of vertices to be covered by the image edges incident with the cluster. Therefore, every cluster coincides with an even number of image edges which are in M; in other words, these edges form the image of an eulerian subgraph of G.

Vice versa, the image of any eulerian subgraph of G can be extended (uniquely) by adding some of the auxiliary edges in G^* to make a perfect matching in G^* . Thus, there is a one-to-one correspondence between the perfect matchings of G^* and the eulerian subgraphs of G. As all the auxiliary edges have weights equal to 0, the corresponding terms contributing to either of the generating functions are equal. Consequently, the two generating functions are equal.

Further in sections 2,3 we show how to calculate $\mathcal{P}(Q_{mmk}, x)$ by embedding Q_{mmk} into a generalised surface S_g so that Q_{mmk} becomes a generalised g-graph. This embedding of Q_{mmk} has a 'planar part' consisting of all the vertical edges, and this part does not play a role in the derivation of the formula, where the 'non-planar' edges are vital. The advantage of the Fisher's construction is that the even splitting of the vertices may be performed in the planar part of Q_{mmk} , hence the paths of vertical edges are replaced by the 'paths of triangles', and the non-planar part of Q_{mmk} remains untouched. Hence Q_{mmk} is turned into Q_{mmk}^* without changing the embedding and an expression analogous to the one described in Corollory 3.9 for the dimer problem holds for the Ising problem as well.

In fact, one should find an analogous expression for the 3-dimensional variants of the problems which may be treated by the Pfaffian method in 2 dimensions, like a variant of the ice problem.

The proof of our result is involved: this paper may be viewed as a continuation of the papers [4], [5], [6], [7]. A theorem of Galluccio and Loebl [4] expresses $\mathcal{P}(G, x)$, where

G is an arbitrary graph, as a linear combination of Pfaffians of matrices associated with *relevant* orientations of G. When G is a bipartite graph like the cubic lattice, the Pfaffians may be turned into determinants. The relevant orientations may be naturally described when the graph is embedded in a certain way on an orientable surface.

This 'Pfaffian approach' to the dimer problem has been started by Kasteleyn [13]. Kasteleyn [13] and Fisher [2] also described methods how to find the Ising partition function for a graph G as the dimer partition function of a locally modified G. In [5] and [6], the Pfaffian method leads to an efficient algorithmic treatment of the Ising problem for finite lattices which may be embedded on a fixed surface, e.g. on a torus. This approach has been recently extended in [19] to non-orientable surfaces.

We use the Pfaffian method to prove Theorem 1.1 as follows: we embed the threedimensional cubic lattice to a 2-dimensional orientable surface, use the theory developed in [4] and finally characterize the coefficients of the resulting linear combination and turn it into a probabilistic expression.

Applying elementary probabilistic analysis to the statement of Theorem 1.1 I have obtained a curious corollary which may be of independent interest. Once discovered, the corollary may be proved directly without using Theorem 1.1.

Let Q' be a cubic lattice with added boundary edges, i.e. the degree of each vertex of Q' is six. A subset C of vertices of Q' is called a *cover* if each edge of Q' is incident with exactly one vertex of C. Note that Q' has exactly 2 covers. We fix one of them and denote it by C. A subgraph of Q' is called a *plane* if it is obtained from Q' by deleting both horizontal and/or vertical edges incident with each vertex of the cover C. Hence each vertex of C has degree 2 or 4 in any plane. A plane P is called *even* if the number of vertices of C of degree 2 in P is even, and P is called *odd* otherwise.

Theorem 1.2

$$\mathcal{P}(Q', x) = \sum_{W \in A} \mathcal{P}(W, x) - \sum_{W \in B} \mathcal{P}(W, x)$$

where A consists of the even planes and B consists of the odd planes.

Proof. Let M be a perfect matching of Q'. We will compute how M contributes to the RHS. Let Z be the subset of vertices of C incident to the width edges of M, and let z = |Z|. M contributes to a term of the RHS corresponding to a plane P if and only if M is a perfect matching of P. Which planes contain M? Assume M is a perfect matching of a plane P and let $x \in C$. First let x be incident with a horisontal or a vertical edge e of M. Then the degree of x in P is 4 and e determines which edges of P are incident with x. Secondly let x be incident with a width edge e of M. Then all three possibilities may occur in P: x may be incident with the width edges only, or with the width and the horisontal edges, or with the width and the vertical edges. Let P have $i \leq z$ vertices of Z incident with the width edges only. Then M contributes $(-1)^i$ to the term of the RHS corresponding to P. Hence, the total contribution of M equals $\sum_{i=0}^{z} (-1)^i 2^{z-i} {z \choose i}$, which equals $(2-1)^z = 1$ by binomial theorem. \Box

The next section will describe a theorem of Galluccio and Loebl ([4]) which forms a basis of the proof of Theorem 1.1. The basic notation, definitions and some relevant simple facts may be found in the appendix.

2 Generalized g-graphs.

It is recommended to read the appendix first before starting with this section.

Definition 2.1 A surface S_g of genus g consists of a base B_0 and 2g bridges B_j^i , i = 1, ..., g and j = 1, 2, where

- i) B_0 is a convex 4g-gon with vertices $a_1, ..., a_{4g}$ numbered clockwise;
- ii) B_1^i , i = 1, ..., g, is a 4-gon with vertices $x_1^i, x_2^i, x_3^i, x_4^i$ numbered clockwise. It is glued with B_0 so that the edge $[x_1^i, x_2^i]$ of B_1^i is identified with the edge $[a_{4(i-1)+1}, a_{4(i-1)+2}]$ of B_0 and the edge $[x_3^i, x_4^i]$ of B_1^i is identified with the edge $[a_{4(i-1)+3}, a_{4(i-1)+4}]$ of B_0 ;
- iii) B_2^i , i = 1, ..., g, is a 4-gon with vertices $y_1^i, y_2^i, y_3^i, y_4^i$ numbered clockwise. It is glued with B_0 so that the edge $[y_1^i, y_2^i]$ of B_2^i is identified with the edge $[a_{4(i-1)+2}, a_{4(i-1)+3}]$ of B_0 and the edge $[y_3^i, y_4^i]$ of B_2^i is identified with the edge $[a_{4(i-1)+4}, a_{4(i-1)+5(mod4g)}]$ of B_0 .

Observe that in Definition 2.1 we denote by [a, b] edges of polygons and not edges of graphs. The usual representation in the space of an orientable surface S of genus g may be then obtained from its polygonal representation S_g by the following operation: for each bridge B, glue together the two segments which B shares with the boundary of B_0 , and delete B.

Definition 2.2 A graph G is called a g-graph if it may be embedded on S_g so that all the vertices belong to the base B_0 , and the embedding of each edge uses at most one bridge. The set of the edges embedded entirely on the base will be denoted by E_0 and the set of the edges embedded on the bridge B_j^i will be denoted by E_j^i , $i = 1, \ldots, g$, j = 1, 2. We also let $G_0 = (V, E_0)$ and $G_j^i = (V, E_0 \cup E_j^i)$. Moreover the following conditions need to be satisfied too.

- 1. the outer face of $G_0 = (V, E_0)$ is a cycle, and it is embedded on the boundary of B_0 ,
- 2. if $e \in E_1^i$ then e is embedded entirely on B_1^i and one end vertex of e belongs to $[x_1^i, x_2^i]$ and the other one belongs to $[x_3^i, x_4^i]$. Similarly, if $e \in E_2^i$ then e is embedded entirely on B_2^i and one end vertex of e belongs to $[y_1^i, y_2^i]$ and the other one belongs to $[y_3^i, y_4^i]$.

From now on, we shall consider g-graphs together with a fixed embedding on S_g . Given a g-graph G, we denote by C_0 the cycle which forms the outer face of G_0 . **Definition 2.3** Let G be a g-graph and let $G_j^i = (V, E_0 \cup E_j^i)$. If we draw $B_0 \cup B_j^i$ on the plane as follows: B_0 along with the edges of the polygons belonging to its boundary is unchanged, and the edge $[x_1^i, x_4^i]$ ($[y_1^i, y_4^i]$ respectively) of B_j^i is drawn so that it belongs to the external boundary of $B_0 \cup B_j^i$, we obtain a planar embedding of G_j^i . This embedding will be called planar projection of E_j^i outside B_0 .

Definition 2.4 Let G = (V, E) be a g-graph. An orientation D_0 of G_0 such that each inner face of each 2-connected component of G_0 is clockwise odd in D_0 is called a basic orientation of G_0 .

Note that a basic orientation always exists for a planar graph. Kasteleyn [13] proved that if D is a basic orientation of a planar graph G then the contributions of all perfect matchings of G have the same sign in Pf(A(D)).

From now on we shall fix a basic orientation D_0 for each g-graph.

Definition 2.5 Let G = (V, E) be a g-graph and D_0 a basic orientation of G_0 . We define the orientation D_j^i of each G_j^i as follows: We consider G_j^i embedded on the plane by the planar projection of E_j^i outside B_0 (see Definition 2.3), and complete the basic orientation D_0 of G_0 to an orientation of G_j^i so that each inner face of each 2-connected component of G_j^i is clockwise odd.

The orientation $-D_i^i$ is defined by reversing the orientation D_i^i of G_i^i .

Observe that after fixing a basic orientation D_0 , the orientation D_j^i is uniquely determined for each i, j.

Definition 2.6 Let G be a g-graph, $g \ge 1$. An orientation D of G which equals the basic orientation D_0 on G_0 and which equals D_j^i or $-D_j^i$ on E_j^i is called relevant. We define its type $r(D) \in \{+1, -1\}^{2g}$ as follows: For $i = 0, \ldots, g-1$ and $j = 1, 2, r(D)_{2i+j}$ equals +1 or -1 according to the sign of D_j^{i+1} in D.

Definition 2.7 Let G be a g-graph and D a relevant orientation of G. Let $r(D) = (r_1, ..., r_{2g})$. We let c(r(D)) equal the product of c_i , i = 0, ..., g-1, where $c_i = c(r_{2i+1}, r_{2i+2})$ and c(1, 1) = c(1, -1) = c(-1, 1) = 1/2 and c(-1, -1) = -1/2.

Observe that $c(r(D)) = (-1)^n 2^{-g}$, where $n = |\{i; r_{2i+1} = r_{2i+2} = -1\}|$.

The following theorem is proved in Galluccio, Loebl [4]. See appendix for the definition of s(D, M).

Theorem 2.8 Let G be a g-graph with a perfect matching $M_0 \subset E_0$. If we order the vertices of G so that $s(D_0, M_0)$ is positive then

$$\mathcal{P}(G, x) = \sum_{i=1}^{4^g} c(r(D_i)) Pf(A(D_i))$$

where D_i , $i = 1, ..., 4^g$, are the relevant orientations of G.

We need a generalization of the notion of a g-graph.

Definition 2.9 Any graph G obtained by the following construction will be called generalized g-graph.

- 1. Let $g = g_1 + ... + g_n$ be a partition of g into positive integers.
- 2. Let S_{g_i} be a surface of genus g_i , i = 1, ..., n. Let us denote the basis and the bridges of S_{g_i} by B_0^i and $B_{i,k}^i$, i = 1, ..., n, $j = 1, ..., g_i$ and k = 1, 2.
- 3. For i = 1, ..., n let H_i be a g_i -graph with the property that the subgraph of H_i embedded on B_0^i is a cycle, embedded on the boundary of B_0^i . Let us denote it by C^i .
- 4. Let G_0 be a 2-connected plane graph and let $F_1, ..., F_n$ be a subset of faces of G_0 . Let K^i be the cycle bounding F_i , i = 1, ..., n. Let each K^i be isomorphic to C^i .
- 5. Then G is obtained by glueing the H_i 's into G_0 so that each K^i is identified with C_i .

For each generalized g-graph G we can define 4^g relevant orientations $D_1, ..., D_{4^g}$ with respect to a fixed basic orientation of G_0 , and coefficients $c(r(D_i))$, i = 1, ..., n in the same way as for a g-graph. The following theorem can be proved in the same way as Theorem 2.8 since each H_i may be treated independently. In fact, G. Tessler chose this more general setting in his paper [19].

Theorem 2.10 Let G be a generalized g-graph with a perfect matching M_0 of G_0 . Let D_0 be a basic orientation of G_0 . If we order the vertices of G so that $s(D_0, M_0)$ is positive then

$$\mathcal{P}(G, x) = \sum_{i=1}^{4^g} c(r(D_i)) Pf(A(D_i))$$

where D_i , $i = 1, ..., 4^g$, are the relevant orientations of G.

3 Cubic lattices as generalized g-graphs.

In this section we will describe how to draw 3-dimensional cubic lattices as generalized g-graphs. Let m, n be odd positive integers such that k = (n - 1)/2 is even. Let us use Q to denote the cubic lattice $Q_{m,m,n}$. Let us denote vertical path $(V_{xy1}, ..., V_{xyn})$ of Q by $V_{xy}(Q) = V_{xy}$ and let \bar{V}_{xy} denote V_{xy} traversed in the opposite direction.

Let $H_x(Q) = H_x = \{h_{xyz}; z = 1, ..., n, y = 1, ..., m\}$ and $W_{xy}(Q) = W_{xy} = \{w_{xyz}; z = 1, ..., n\}$.

How to draw Q on the plane.

First draw the paths V_{xy} along a cycle in the following natural way:

 $V_{11}, \bar{V}_{12}, V_{13}, \dots V_{1m}, \bar{V}_{2m}, V_{2(m-1)}, \dots, \bar{V}_{21}, V_{31}, \dots, V_{mm}.$

Next, draw the horizontal edges inside this cycle, and the width edges outside of this cycle as depicted in Fig. 1 below where $Q = Q_{3,3,3}$ is properly drawn.



Figure 1

For each x = 1, ..., m - 1 the curves representing the edges of H_x are pairwise disjoint and for x = 2, ..., m - 2 the curves representing the edges of H_x intersect the curves representing the edges of H_{x-1} and H_{x+1} . We keep the following rule: the interiors of the curves representing h_{xyz} and $h_{(x+1)yz}$ intersect if and only if z is even.

For each x = 1, ..., m and y = 1, ..., (m - 1) the curves representing the edges of W_{xy} are pairwise disjoint and for y = 2, ..., m - 2 the curves representing the edges of W_{xy} intersect the curves representing the edges of $W_{x(y-1)}$ and $W_{x(y+1)}$. We again keep the rule that the interiors of the curves representing w_{xyz} and $w_{x(y+1)z}$ intersect if and only if z is even. The curve representing an edge e will be denoted by C(e).

Now we modify Q into a generalized g-graph Q'.

Width construction. First we describe the modification for W_x , x = 1, ..., m. The modification is described by Fig. 2 where the construction is illustrated on edges among $V_{x(y-1)}, V_{xy}$ and $V_{x(y+1)}$ for x odd and y < m - 1 even.



For each x = 1, ..., m perform the following construction:

- 1. For each y even let $Aux_1 = \{w_{xyz}; z \text{ odd }\}$. For each edge e of Aux_1 introduce a new vertex to each intersection of C(e) with the curves representing the edges of $W_{x(y-1)} \cup W$, where $W = W_{x(y+1)}$ in case y < m-1 and $W = \emptyset$ otherwise. By this operation, each $e \in Aux_1$ is replaced by a path. Call each edge of this path *auxiliary*.
- 2. For each y even let $Aux_2 = \{w_{x(y-1)1}, w_{x(y-1)n}\} \cup A$, where $A = \{w_{x,(y+1)1}, w_{x(y+1)n}\}$ in case y < m - 1 and $A = \emptyset$ otherwise. For each edge e of Aux_2 introduce a new vertex to each intersection of C(e) with the curves representing the edges of W_{xy} . Hence each $e \in Aux_2$ is replaced by a path. Call each edge of this path *auxiliary*.

For each y even the edges $v_{xy1}, v_{xy(n-1)}$ and also $v_{x(y+1)1}, v_{x(y+1)(n-1)}$ will also be called *auxiliary*.

In Figure 2, the auxiliary edges are represented by dashed lines.

- 3. We introduce a new variable a which we associate with each auxiliary edge e and we let w(e) = 1. Hence the term associated with each auxiliary edge e has form $a^{w(e)} = a$.
- 4. The edges w_{xyz} , y even and z even will be called *relevant* for Q. If y < m 1 then the relevant edges are subdivided by two vertices (added in 2.) into three edges of Q'. The middle one will be called *special* and the other two *long*.

If y = m - 1 then the relevant edge w_{xyz} is subdivided by one vertex into two edges of Q'. The one incident to V_{xm} will be called *special* and the other one *long*.

If e is a relevant edge of Q, then we choose a corresponding long edge f and we let w(e) = w(f). We let the weight of the special edge and of the remaining long edge be equal to 0.

- 5. The edges of $W_{x(y-1)} \cup W$ also got subdivided by new vertices introduced in step 2 and step 3.
- 6. We delete all edges of the paths obtained from $w_{x(y-1)z}$ and $w_{x(y+1)z}$, 1 < z < n odd. In Figure 2, the deleted edges are represented by dotted lines.
- 7. Each edge $e \in \{w_{x(y-1)z}, w_{x(y+1)z}; z \text{ even }\}$, is subdivided by new vertices introduced in 2. into a path. We let the weights assigned to the edges of the path equal 0 except of one initial edge whose weight is let equal w(e). The edge e of this path such that the interior of C(e) does not intersect interior of any curve representing a long edge will also be *special*.
- 8. All vertical edges which are not auxiliary (see 2.) will be called *special*. In Figure 2, the special edges are represented by normal lines.

This finishes the construction for the width edges. In Figure 2, the edges which are neither auxiliary nor special nor deleted are represented by fat lines.

Horizontal construction. Now we perform an analogous construction with the horizontal edges of Q.

- 1. For each x even let $Aux_3 = \{h_{xyz}; z \text{ odd }\}$. For each edge e of Aux_3 introduce a new vertex to each intersection of C(e) with the curves representing the edges of $H_{x-1} \cup K$, where $K = H_{x+1}$ in case x < m-1 and K = set otherwise. By this operation, each $e \in Aux_3$ is replaced by a path. Call each edge of this path *auxiliary*.
- 2. For each x even let $Aux_4 = \{h_{(x-1)11}, h_{(x-1)nn}\} \cup B$, where $B = \{h_{(x+1)11}, h_{(x+1)nn}\}$ in case x < m-1 and $B = \emptyset$ otherwise. For each edge e of Aux_4 introduce a new vertex to each intersection of C(e) with the curves representing the edges of H_x . Hence each $e \in Aux_4$ is replaced by a path. Call each edge of this path *auxiliary*.
- 3. We again associate variable a with each auxiliary edge e and we let w(e) = 1.
- 4. The edges h_{xyz} , x even and z even will be called *relevant* for Q. If x < m 1 then the relevant edges are subdivided by two vertices (added in 2.) into three edges of Q'. The middle one will be called *special* and the other two *long*.

If x = m - 1 then the relevant edge h_{xyz} is subdivided by one vertex into two edges of Q'. The one incident to V_m , will be called *special* and the other one *long*.

If e is a relevant edge of Q, then we choose a corresponding long edge f and we let w(e) = w(f). We let the weight of the special edge and of the remaining long edge equal 0.

5. The edges of $H_{x-1} \cup K$ also got subdivided by new vertices introduced in step 2 and step 3.

- 6. We delete all edges of the paths obtained from $h_{(x-1)yz}$ and $h_{(x+1)yz}$, 1 < z < n odd.
- 7. Each edge $h \in \{h_{(x-1)yz}, h_{(x+1)yz}; z \text{ even }\}$, is subdivided by new vertices introduced in 2. into a path. We let the weights assigned to the edges of the path equal 0 except of one initial edge whose weight is let equal w(h). Each edge e of this path such that the interior of C(e) does not intersect interior of any curve representing a long edge will also be *special*.

Final steps. Let Aux denote the set of all auxiliary edges. Then Q' - Aux is a subdivision of Q_{mmk} . We subdivide some special edges so that the graph Q = Q' - Aux is an even subdivision of Q_{mmk} . All these new edges will be special, and we set their weights equal 0.

The subgraph of Q' formed by the special edges consists of vertical paths (of odd length) and some other mutually disjoint paths which may have at most one vertex in common with the vertical paths. Hence there is matching M'_0 of special edges covering all the vertices of the vertical paths, and all but possibly one vertex of each of the additional paths of special edges. We conclude the construction of Q' by subdividing some auxiliary edges in such a way that Q' has a perfect matching M' consisting of special and auxiliary edges only, which extends M'_0 (i.e. $M'_0 \subset M'$). All these new edges will be auxiliary; we again associate variable a with them and we set their weights equal 1.

This finishes the construction of Q'.

Some properties of Q'.

- 1. Each edge e of Q' such that C(e) does not intersect any curve representing other edge in its interior is auxiliary or special. Let us denote the plane subgraph of Q'formed by the auxiliary and special edges by Q^p .
- 2. Any other edge of Q' is drawn on a face of Q^p . Moreover, the edges drawn on a face of Q^p may be drawn onto a pair of bridges above this face, where one of the bridges contains one long edge, and the other bridge contains the remaining edges. Hence, we may view Q' as a generalized g-graph with the planar part equaled to Q^p .
- 3 The special edges form an acyclic subgraph of Q^p (see Fig. 2). Hence any orientation of the special edges may be extended into a basic orientation of Q^p . We will choose basic orientation D^p of Q^p with the following properties:
 - 1. D^p on special edges is in agreement with the natural ordering ,
 - 2. Perfect matching M' has positive sign in $Pf(A(D^p))$,
 - 3. The orientation of edges on a bridge has positive sign if and only if it is in agreement with the natural ordering $(V_{11}\bar{V}_{12}....V_{1m}\bar{V}_{21}....V_{mm})$.
- 4. We constructed Q' so that Q = Q' Aux is an even subdivision of Q_{mmk} . If w is a vector of weights associated with Q_{mmk} then let w' be the vector of weights associated with Q and induced by w and let w'' be the vector of weights associated with Q' which equals w' on the edges of Q and w''(e) = 1 for each auxiliary edge e of Q'. If we let a = 0 we have $\mathcal{P}(Q', x, a) = \mathcal{P}(Q_{mmk}, x)$.

We have described how to view Q_{mmk} , m odd and k even, as a generalized g-graph Q'. Now we can use Theorem 2.10 for Q' to compute $\mathcal{P}(Q_{mmk}, x)$.

The relevant orientations of Q'.

Each relevant edge of Q corresponds to unique edge of Q_{mmk} ; this unique edge will also be called *relevant* in Q_{mmk} . Hence the relevant edges of Q_{mmk} are: $w_{xyz}(Q_{mmk})$, x = 1, ..., m, y even and and $h_{xyz}(Q_{mmk})$, x even. Hence there are 1/2km(m-1) relevant width edges and 1/2(m-1)mk relevant horizontal edges in Q_{mmk} .

We let \mathcal{R} be the set of relevant edges of Q_{mmk} and $C_r = |\mathcal{R}| = km(m-1)$ denote the number of relevant edges of Q_{mmk} .

The set \mathcal{S} of the edges of $V_{ij}(Q_{mmk})$, i, j = 1, ...m, corresponds to a subset of special edges of Q'. The orientation D^p induces orientation \mathcal{S}^d of \mathcal{S} which is in agreement with the natural ordering $(V_{11}(Q_{mmk})\bar{V}_{12}(Q_{mmk})....V_{1m}(Q_{mmk})\bar{V}_{21}(Q_{mmk})....V_{mm}(Q_{mmk}))$.

Each relevant orientation D' of Q' is determined by the fixed basic orientation D^p of Q^p , and by a pair of signs for each pair of bridges. Each pair of bridges is associated with a long edge of Q'. Hence these signs may be given by specifying $(d_{D'}^1(e), d_{D'}^2(e)) \in \{+-\}^2$, for each long edge e, where $d_{D'}^1(e)$ denotes the sign of the bridge containing e, and $d_{D'}^2(e)$ denotes the sign of the other bridge.

The long edges of Q' are associated with relevant edges of Q, and hence also with relevant edges of Q_{mmk} .

The relevant edges $w_{x(m-1)z}(Q_{mmk})$ and $h_{(m-1)yz}(Q_{mmk})$ are associated with only one long edge of Q'. If e is such relevant edge of Q_{mmk} , we will call e border edge and we denote by e_1 the corresponding long edge. We let $d_{D'}(e) = (d_{D'}^1(e_1), d_{D'}^2(e_1), +, +)$.

Let $C_b = 2mk$ denote the number of border edges.

Each relevant non-border edge e of Q_{mmk} has two long edges e_1, e_2 associated with it. We let $d_{D'}(e) = (d_{D'}^1(e_1), d_{D'}^2(e_1), d_{D'}^1(e_2), d_{D'}^2(e_2))).$

A relevant vector is any element r of $[\{+,-\}^4]^{\mathcal{R}}$ such that $r(e)_3 = r(e)_4 = +$ for each relevant border edge e of Q_{mmk} . Hence there are $4^{2C_r-C_b}$ relevant vectors.

There is a natural bijection between relevant orientations of Q' and relevant vectors. If s is a relevant vector, then let D'(s) denote the corresponding relevant orientation of Q'and let sgn(s) of a relevant vector s be calculated according to Theorem 2.10 as follows: $sgn(s) = (-1)^{|\{(e,i);i=0,1,s(e)_{2i+1}=s(e)_{2i+2}=-1\}|}$.

If D'(s) is a relevant orientation of Q' then let D(s) denote the orientation of Q_{mmk} induced by D'(s) (see 5.4). An orientation of Q_{mmk} will be called *relevant* if it equals D(s) for some relevant vector s.

Let M be the (uniquely determined) perfect matching of Q_{mmk} consisting of the edges of \mathcal{S} only. Recall that perfect matching M' has positive sign in $Pf(A(D^p))$ and similarly, perfect matching M has positive sign in $Pf(A(\mathcal{S}^d))$. Using Theorem 2.10 and Theorem 5.5 we have the following.

Theorem 3.1

$$\mathcal{P}(Q_{mmk}, x) = 2^{-2C_r + C_b} \sum sgn(s) Pf(A(D(s)))$$

where the sum is over all relevant vectors.

Note that possibly D(r) = D(r') for relevant vectors $r \neq r'$. Next we clarify this.

Definition 3.2 We define an equivalence * on the relevant vectors as follows. r * s if the following holds: there is exactly one relevant non-border edge e such that $r(e) \neq s(e)$ and r(f) = s(f) for each $f \neq e$. Moreover, $r(e)_1 \neq s(e)_1$, $r(e)_3 \neq s(e)_3$, $r(e)_2 = s(e)_2 \neq$ $r(e)_4 = s(e)_4$.

Proposition 3.3 If r * s then D(r) = D(s) and $sgn(r) \neq sgn(s)$.

Proof. If r * s then D(r) = D(s) by the definition of '*'. Moreover $sgn(r) \neq sgn(s)$ since $r(e)_2 = s(e)_2 \neq r(e)_4 = s(e)_4$ where e is the only relevant edge for which $r(e) \neq s(e)$. \Box

Definition 3.4 A relevant vector r is called useful if it forms a one-element class w.r.t. equivalence *, i.e. if $r(e)_2 = r(e)_4$ for each relevant non-border edge e.

Corollary 3.5

$$\mathcal{P}(Q_{mmk}, x) = 2^{-2C_r + C_b} \sum sgn(r) Pf(A(D(r)))$$

where the sum is over all useful vectors r.

Definition 3.6 If r, s are useful vectors we write r * *s if D(r) = D(s).

Proposition 3.7 1. Each equivalence class of '**' has $2^{C_r-C_b}$ elements.

2. If r * *s then sgn(r) = sgn(s).

Proof. Let r be a useful vector. Then D(r) determines uniquely $r(e)_2$ and $r(e)_4$ for each relevant edge e and also $r(f)_1$ for each relevant border edge f. Hence D(r) determines uniquely r(f) for each relevant border edge f. Moreover D(r) determines uniquely the product $r(e)_1 \times r(e)_3$ for each relevant non-border edge e. Since there are $C_r - C_b$ relevant non-border edges, each equivalence class of '**' has $2^{C_r-C_b}$ elements.

Let r, s be useful and let r * *s. Then $r(e)_2 = s(e)_2 = s(e)_4 = r(e)_4$ for each relevant non-border edge e and $r(e)_2 = s(e)_2$ and $s(e)_1 = r(e)_1$ for each relevant border edge. This implies that sgn(r) = sgn(s). \Box

Proposition 3.8 If D is an orientation of Q_{mmk} that extends S^d then there is uniquely determined class C of equivalence ** such that D = D(r) for each $r \in C$.

Hence, given an orientation D of Q_{mmk} that extends S^d , let us call it *stable orientation* and let us define its sign sgn(D) to be equal to sgn(r) for any useful vector r such that D = D(r). This is well defined by Proposition 3.7.

Corollary 3.9

$$\mathcal{P}(Q_{mmk}, x) = 2^{-C_r} \sum sgn(D) Pf(A(D))$$

over all stable orientations D.

We continue by characterizing sgn(D).

As we noticed before, $sgn(r) = (-1)^{|\{(e,i);i=0,1,r(e)_{2i+1}=r(e)_{2i+2}=-1\}|}$. If r is a stable vector then $r(e)_2 = r(e)_4$ for each relevant non-border edge e and we get the following observation.

Proposition 3.10 Let r be a stable vector. Then $sgn(r) = (-1)^{|\{e; r(e)_1 r(e)_3 = -1, r(e)_2 = -1\}|}$.

Definition 3.11 Let D be a stable orientation. We define orientation \overline{D} as follows:

- 1. For each x, y, z such that y < m odd do the following: let n(xyz) be the number of arcs w_{xab} , $a \leq y$ odd and $z \leq b \leq (z+1)$, oriented in D against the natural ordering. If n(xyz) odd then we orient w_{xyz} in \overline{D} against the natural ordering, else according to the natural ordering.
- 2. For each x, y, z such that x < m odd do the following: let n(xyz) be the number of arcs h_{abc} oriented in D against the natural ordering. Here (abc) are the triples of indices satisfying $a \leq x$ odd and $(y, z) \leq (b, c) \leq (y', z')$ where the order is lexicographic and (y', z') is immediate successor of (y, z).

If n(xyz) odd then we orient h_{xyz} in \overline{D} against the natural order, else according to the natural order.

3. All the remaining arcs orient in \overline{D} in the same way as in D.

Note that relevant edges are oriented in the same way in both D and \overline{D} .

Proposition 3.12 Let D be a stable orientation and let r be a stable vector such that D = D(r). Let e be a relevant edge of Q_{mmk} . Then

- 1. $r(e)_1 \times r(e)_3 = -1$ if and only if e is oriented in D (and hence also in D) against the natural ordering.
- 2. If $e = w_{xyz}$, y even then $r(e)_2 = -1$ if and only if $w_{x(y-1)z}$ is oriented in D against the natural ordering. If $e = h_{xyz}$, x even then $r(e)_2 = -1$ if and only if $h_{(x-1)yz}$ is oriented in \overline{D} against the natural ordering.

Proof. $r(e)_1 \times r(e)_3 = -1$ if and only if exactly one long edge of Q' corresponding to e is oriented in D' (we remind that D is induced by orientation D' of Q') against the natural ordering. Since D is induced by D', this happens if and only if e is oriented in both D and \overline{D} against the natural ordering.

It remains to prove 2. We will show the case $e = w_{xyz}$ since the other case is completely analogous. If f is an edge of Q_{mmk} then we let f(D) = 1 if f is oriented in D according to the natural order, and we let f(D) = -1 otherwise. We proceed by induction on (y, k-z).

Firstly assume y = 2 and z = k. In this simplest case $r(e)_2 = w_{x1k}(D)$ (see Fig. 2). Moreover the orientation of w_{x1k} is the same in both D and \overline{D} .

Secondly let y = 2 and z < k. Then $w_{x1z}(D) = w_{x1(z+1)}(D) \times r(e)_2$ (see Fig. 2). It follows from Definition 3.11 that $r(e)_2 = -1$ if and only if w_{x1z} is oriented in \overline{D} against the natural order.

Thirdly, let y = 4 and z = k. Then $w_{x3k}(D) = w_{x1k}(D) \times r(e)_2$ (see Fig. 2). It follows from Definition 3.11 that $r(e)_2 = -1$ if and only if w_{x3k} is oriented in \overline{D} against the natural order.

Fourthly, let y = 4 and z < k. Then $w_{x3z}(D) = w_{x3(z+1)}(D) \times r(e)_2 \times r(w_{x2z})_2 = w_{x3(z+1)}(D) \times r(e)_2 \times w_{x1z}(D) \times w_{x1(z+1)}(D)$ (see Fig. 2). It follows from Definition 3.11 that $r(e)_2 = -1$ if and only if w_{x3z} is oriented in \overline{D} against the natural ordering.

In general if $e = w_{xyz}$, y even and z = k then $w_{x(y-1)k}(D) = r(w_{x(y-2)k})_2 \times r(e)_2 = r(e)_2 \times \prod(w_{xy'k}(D); y' < y - 1 \text{ odd })$. It follows from Definition 3.11 that $r(e)_2 = -1$ if and only if $w_{x(y-1)k}$ is oriented in \overline{D} against the natural order.

Finally if $e = w_{xyz}$, y even and z < k then $w_{x(y-1)z}(D) = w_{x(y-1)(z+1)}(D) \times r(e)_2 \times r(w_{x(y-2)z})_2 = w_{x(y-1)(z+1)}(D) \times r(e)_2 \times \prod(w_{xy'z}(D); y' < y-1 \text{ odd }) \times \prod(w_{xy'z}(z+1)(D); y' < y-1 \text{ odd })$. It follows from Definition 3.11 that $r(e)_2 = -1$ if and only if w_{x3z} is oriented in \overline{D} against the natural order. \Box

Corollary 3.13 Let D be a stable orientation. Then $sgn(D) = (-1)^{h+\sum_{x=1}^{m} w(x)}$, where $w(x) = |\{(yz); y \text{ even and both } w_{xyz}, w_{x,(y-1),z} \text{ are oriented against the natural ordering in } \bar{D}\}|; h = |\{(xyz); x \text{ even and both } w_{xyz}, w_{(x-1),y,z} \text{ are oriented against the natural ordering in } \bar{D}\}|.$

We can also write the sign in the following form: for edge e we let $s_D(e) = 0$ if the orientation of e agrees with the natural order, and $s_D(e) = 1$ otherwise.

Corollary 3.14 Let D be a stable orientation. Then

$$sgn(D) = (-1)^{\sum_A s_D(w_{x,2y,z})s_D(w_{x,2y'-1,z'}) + \sum_B s_D(h_{2x,y,z})s_D(h_{2x'-1,y',z'})}$$

where

$$A = \{(x, y, z, y', z'); 1 \le x \le m, 1 \le y \le (m-1)/2, 1 \le z \le k, 1 \le y' \le y, z \le z' \le z+1\}$$

and

$$B = \{(x, y, z, x', y', z'); 1 \le x \le (m - 1)/2, 1 \le z \le k, 1 \le y \le m, 1 \le x' \le x, (y, z) \le (y', z') \le (y'', z'')\}.$$

In the definition of B the order on pairs of integers is lexicographic order and (y'', z'') is the immediate successor of (y, z).

Proposition 3.15 There are 2^{2C_r} stable orientations. There are $2^{C_r-1}(2^{C_r}+1)$ stable orientations with positive sign.

Proof. The first statement follows directly from the definition of a stable orientation. For Q_{132} there are $16 = 2^4 = 2^{2C_r}$ stable orientations (see the remark after Theorem 1.1 for the definition of C_r), from which 10 have positive sign. Hence the difference between the number of stable orientations of positive sign and stable orientations of negative sign is $4 = 2^{C_r}$. For Q_{152} there are $4^4 = 2^{2C_r}$ stable orientations from which $10 \times 10 + 6 \times 6$ have positive sign and $2(6 \times 10) = 120$ have negative sign. Hence the difference between the number of stable orientations of positive sign and stable orientations of negative sign is $10 \times 10 + 6 \times 6 - 2(6 \times 10) = (10 - 6)(10 - 6) = 2^{C_r}$. Similarly by induction on a we get that for $Q_{1(2a+1)2}$ there are $4^{2a} = 2^{2C_r}$ stable orientations, and the difference between the number of stable orientations of positive sign and stable orientations of negative sign is 2^{2a} . Similarly we calculate the difference for Q_{13k} , k even, and by induction also for Q_{1mk} . That takes care of one layer of width edges, and the layers are independent so the differences multiply as indicated above. In the same way we calculate the contribution of the horisontal edges. Summarising for Q_{mmk} the difference between the number of stable orientations of positive sign and those of negative sign equals $2^{C_r} = 2^{(m-1)km}$. From this Proposition follows. \Box

4 From Pfaffians to determinants.

In the introduction we let \mathcal{Z} be square $(Z_1 \times Z_2)$ matrix defined by $\mathcal{Z}_{ij} = x_{ij}$ if ij is an edge of Q_{mmk} and $\mathcal{Z}_{ij} = 0$ otherwise.

Let D be an orientation of Q_{mmk} . In the introduction we associate a signing Z(D) of \mathcal{Z} with it such that $Z(D)_{ij} = x_{ij}$ if $(ij) \in E(D)$, $Z(D)_{ij} = -x_{ij}$ if $(ji) \in E(D)$, and $Z(D)_{ij} = 0$ otherwise.

Note that Pf(A(D)) = det(Z(D)). Hence we can reformulate Corollary 3.9:

Corollary 4.1 $\mathcal{P}(Q_{mmk}, x) = 2^{-C_r} \sum sgn(D)det(Z(D))$ where the sum is over all stable orientations D of Q_{mmk} .

Recall that M is unique perfect matching consisting only of vertical edges.

Proposition 4.2 The average of det(Z(D)), D stable, equals $x^{w(M)}$.

Proof. By the linearity of expectation and the definition of stable orientations, the contribution of other than vertical edges cancel out when we calculate the average of det(Z(D)), D stable. Since Q_{mmk} has exactly one perfect matching consisting of vertical edges only, Proposition follows. \Box

Proof of Theorem 1.1. By Corollary 4.1 and Proposition 3.15 we have that

$$\mathcal{P}(Q_{mmk}, x) = 2^{-C_r} [-2^{2C_r} x^{w(M)} + 2\alpha (2^{C_r-1} (2^{C_r}+1)] = -2^{-C_r+2C_r} x^{w(M)} + \alpha (2^{C_r}+1). \qquad \Box$$

Conclusion. We have expressed the partition functions of the dimer problem and the Ising problem in 3-dimensional finite cubic lattices by means of expectations of the determinants of matrices associated with the cubic lattices. This may open a possibility to apply fundamentally different statistical methods and Monte Carlo simulations to study these problems.

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5 Appendix: Basic notation, definitions and facts.

Let G = (V, E) be a graph. We will assume that a weight w(e) is associated with each edge e of G. If $A \subset E$ then we let $w(A) = \sum_{e \in A} w(e)$. A graph G' = (V', E') is called a *subgraph* of G if $V' \subset V$ and $E' \subset E$. Let $\{v_1, e_1, v_2, e_2, ..., v_i, e_i, v_{i+1}, ..., e_n, v_{n+1}\}$ be a sequence such that each v_j is a vertex of a graph G, each e_j is an edge of G and $e_j = v_j v_{j+1}$, and $v_i \neq v_j$ for i < j except if i = 1 and j = n + 1. If also $v_1 \neq v_{n+1}$ then P is called a path of G. If $v_1 = v_{n+1}$ then P is called a cycle of G. In both cases the length of P equals n. When no confusion arises we shall also denote paths by simply listing their vertices or edges, namely $P = (v_1, v_2, \ldots, v_{n+1})$ or $P = (e_1, e_2, \ldots, e_n)$.

A graph G = (V, E) is connected if it has a path between any pair of vertices, and it is 2-connected if the graph $G_v = (V - \{v\}, \{e \in E; v \notin e\})$ is connected for each vertex vof G. Each maximal 2-connected subgraph of G is called a 2-connected component of G.

A graph G' is a subdivision of a graph G if some edges of G are replaced in G' by paths so that the inner vertices of each such new path have all degree 2 in G', and both terminal vertices coincide with the vertices of the corresponding deleted edge of G. G' is an *even subdivision* of G if the new paths all have odd lengths. Let w be the vector of the weights associated with the edges of G. We define *induced weights* w' for G' as follows: if e is an edge of G which is replaced by path (e_1, \ldots, e_n) in G' consisting of n edges then $w'(e_1) = w(e), w'(e_j) = 0$ for each j > 1 and w'(f) = w(f) for the remaining edges f of G.

Let G = (V, E) be a graph with 2n vertices and D an orientation of G. Denote by A(D) the skew-symmetric matrix with the rows and the columns indexed by V, where $a_{uv} = x^{w(u,v)}$ in case (u, v) is an arc of D, $a_{u,v} = -x^{w(u,v)}$ in case (v, u) is an arc of D, and $a_{u,v} = 0$ otherwise.

The *Pfaffian* of A(D) is defined as

$$Pf(A(D)) = \sum_{P} s^*(P)a_{i_1j_1}\cdots a_{i_nj_n}$$

where $P = \{\{i_1 j_1\}, \dots, \{i_n j_n\}\}$ is a partition of the set $\{1, \dots, 2n\}$ into pairs, $i_k < j_k$ for $k = 1, \dots, n$, and $s^*(P)$ equals the sign of the permutation $i_1 j_1 \dots i_n j_n$ of $12 \dots (2n)$.

Each nonzero term of the expansion of the Pfaffian of A(D) equals $x^{w(P)}$ or $-x^{w(P)}$ where P is a perfect matching of G. If s(D, P) denote the sign of the term $x^{w(P)}$ in the expansion, we may write

$$Pf(A(D)) = \sum_{P} s(D, P) x^{w(P)}.$$

The Pfaffian is a determinant-type expression. Note the following classic result of Cayley (see [1]).

Theorem 5.1 Let G be a graph and let D be an orientation of G. Then

$$Pf^{2}(A(D)) = det(A(D)).$$

Let $A\Delta B$ denote the symmetric difference of the sets A and B and let $a \stackrel{2}{=} b$ denote a = b modulo 2.

Let M, N be two perfect matchings of a graph G. Then $M\Delta N$ consists of vertex disjoint cycles of even length. These cycles are called *alternating cycles* of M and N.

Let C be a cycle of G of an even length and let D be an orientation of G. C is said to be *clockwise even* in D if it has an even number of edges directed in D in agreement with the clockwise traversal. Otherwise C is called *clockwise odd*. **Definition 5.2** Let G be a graph and let D be an orientation of G. Let M be a perfect matching of G. For each perfect matching P of G let $sgn(D, M\Delta P) = (-1)^n$ where n is the number of clockwise even alternating cycles of M and P, and let $\mathcal{P}(D, M)$ be the sum of $sgn(D, M\Delta P)x^{w(P)}$ over all perfect matchings P of G.

The following theorem was proved by Kasteleyn [13].

Theorem 5.3 Let G be a graph and D an orientation of G. Let P, M be two perfect matchings of G. Then

$$s(D, P) = s(D, M)sgn(D, M\Delta P).$$

Hence,

$$Pf(A(D)) = \sum_{P} s(D, P) x^{w(P)} = s(D, M) \sum_{P} sgn(D, M\Delta P) x^{w(P)} = s(D, M) \mathcal{P}(D, M).$$

In the construction of section 3 we rely on the following definition and theorem.

Definition 5.4 Let G be a graph and let G' be an even subdivision of G. Let D' be an orientation of G'. An orientation D of G induced by D' is constructed as follows: for each edge e of G which was changed into a path P_e in the construction of G', orient e in the direction in which an odd number of edges of P_e is directed in D': this is uniquely determined since P_e has an odd length.

Let G be a graph and let w be a vector of weights associated with the edges of G. Let G' be an even subdivision of G and let w' be the vector of weights associated with the edges of G' induced by w. Each perfect matching P of G gives naturally rise to perfect matching P' of G' such that $x^{w(P)} = x^{w'(P')}$.

Observe that $sgn(D, P\Delta Q) = sgn(D', P'\Delta Q')$ for each pair of perfect matchings P, Q of G. Hence the following theorem follows from Theorem 5.3.

Theorem 5.5 Let G be a graph and let w be a vector of weights associated with the edges of G. Let G' be an even subdivision of G and let w' be a vector of weights associated with the edges of G' induced by w. Let D' be an orientation of G' and let D be the orientation of G induced by D'. Let M be an arbitrary perfect matching of G. Then s(D, M)s(D', M')Pf(A(D')) = Pf(A(D)).

An *embedding* of a graph on a surface is defined in a natural way: the vertices are embedded as points, and each edge is embedded as a continuous non-self-intersecting curve connecting the embedding of its end vertices. The interiors of the embedding of the edges are pairwise disjoint and the interiors of the curves embedding edges do not contain points embedding vertices.

A graph is called *planar* if it may be embedded on the plane. A *plane graph* is a planar graph together with its planar embedding. The embedding of a plane graph partitions

the plane into connected regions called *faces*. The (unique) unbounded face is called *outer face* and the bounded faces are called *inner faces*.

Plane graphs with some regularities are sometimes called 2-dimensional lattices.

Let G be a plane graph. A subgraph of G consisting of the vertices and the edges embedded on the boundary of a face will also be called *a face*. If a plane graph is 2connected then each face is a cycle.

The genus g of a graph G is that of the orientable surface $\mathcal{S} \subset \mathbb{R}^3$ of minimal genus on which G may be embedded.