

Weakly distance-regular digraphs of valency three, I

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Abstract

Suzuki (2004) classified thin weakly distance-regular digraphs and proposed the project to classify weakly distance-regular digraphs of valency 3. The case of girth 2 was classified by the third author (2004) under the assumption of the commutativity. In this paper, we continue this project and classify these digraphs with girth more than 2 and two types of arcs.

Keywords: Weakly distance-regular digraph; Cayley digraph

1 Introduction

A *digraph* Γ is a pair (X, A) where X is a finite set of *vertices* and $A \subseteq X^2$ is a set of *arcs*. Throughout this paper we use the term ‘digraph’ to mean a finite directed graph with no loops. We write $V\Gamma$ for X and $A\Gamma$ for A . A *path* of length r from u to v is a finite sequence of vertices $(u = w_0, w_1, \dots, w_r = v)$ such that $(w_{t-1}, w_t) \in A\Gamma$ for $t = 1, 2, \dots, r$. A digraph is said to be *strongly connected* if, for any two distinct vertices x and y , there is a path from x to y . The length of a shortest path from x to y is called the *distance* from x to y in Γ , denoted by $\partial_\Gamma(x, y)$. The *diameter* of Γ is the maximum value of the distance function in Γ . Let $\tilde{\partial}_\Gamma(x, y) = (\partial_\Gamma(x, y), \partial_\Gamma(y, x))$ and $\tilde{\partial}(\Gamma) = \{\tilde{\partial}_\Gamma(x, y) \mid x, y \in V\Gamma\}$. If no confusion occurs, we write $\partial(x, y)$ (resp. $\tilde{\partial}(x, y)$) instead of $\partial_\Gamma(x, y)$ (resp. $\tilde{\partial}_\Gamma(x, y)$). An arc (u, v) of Γ is of *type* $(1, r)$ if $\partial(v, u) = r$. A path $(w_0, w_1, \dots, w_{r-1})$ is said to be a

circuit of length r if $\partial(w_{r-1}, w_0) = 1$. A circuit is *undirected* if each of its arcs is of type $(1, 1)$. The *girth* of Γ is the length of a shortest circuit.

Let $\Gamma = (X, A)$ and $\Gamma' = (X', A')$ be two digraphs. Γ and Γ' are *isomorphic* if there is a bijection σ from X to X' such that $(x, y) \in A$ if and only if $(\sigma(x), \sigma(y)) \in A'$. In this case, σ is called an *isomorphism* from Γ to Γ' . An isomorphism from Γ to itself is called an *automorphism* of Γ . The set of all automorphisms of Γ forms a group which is called the *automorphism group* of Γ and denoted by $\text{Aut}(\Gamma)$. A digraph Γ is *vertex transitive* if $\text{Aut}(\Gamma)$ is transitive on $V\Gamma$.

Lam [5] introduced a concept of distance-transitive digraphs. A strongly connected digraph Γ is said to be *distance-transitive* if, for any vertices x, y, x' and y' of Γ satisfying $\partial(x, y) = \partial(x', y')$, there exists an automorphism σ of Γ such that $x' = \sigma(x)$ and $y' = \sigma(y)$. Damerell [4] generalized this concept to that of distance-regular digraphs. He showed that the girth g of a distance-regular digraph of diameter d is either 2, d or $d + 1$, and the one with $d = g$ is a coclique extension of a distance-regular digraph with $d = g - 1$. Bannai, Cameron and Kahn [2] proved that a distance-transitive digraph of odd girth is a Paley tournament or a directed cycle. Leonard and Nomura [6] proved that except directed cycles all distance-regular digraphs with $d = g - 1$ have girth $g \leq 8$. In order to find ‘better’ classes of digraphs with unbounded diameter, Damerell [4] also proposed a more natural definition of distance-transitivity, i.e., weakly distance-transitivity. In [8], Wang and Suzuki introduced weakly distance-regular digraphs as a generalization of distance-regular digraphs and weakly distance-transitive digraphs.

A strongly connected digraph Γ is said to be *weakly distance-transitive* if, for any vertices x, y, x' and y' satisfying $\tilde{\partial}(x, y) = \tilde{\partial}(x', y')$, there exists an automorphism σ of Γ such that $x' = \sigma(x)$ and $y' = \sigma(y)$. A strongly connected digraph Γ is said to be *weakly distance-regular* if, for all $\tilde{h}, \tilde{i}, \tilde{j} \in \tilde{\partial}(\Gamma)$ and $\tilde{\partial}(x, y) = \tilde{h}$, the number $p_{\tilde{i}, \tilde{j}}^{\tilde{h}} := |P_{\tilde{i}, \tilde{j}}^{\tilde{h}}(x, y)|$ depends only on $\tilde{h}, \tilde{i}, \tilde{j}$, where

$$P_{\tilde{i}, \tilde{j}}^{\tilde{h}}(x, y) = \{z \in V\Gamma \mid \tilde{\partial}(x, z) = \tilde{i} \text{ and } \tilde{\partial}(z, y) = \tilde{j}\}.$$

The nonnegative integers $p_{\tilde{i}, \tilde{j}}^{\tilde{h}}$ are called the *intersection numbers*. We say that Γ is *commutative* (resp. *thin*) if $p_{\tilde{i}, \tilde{j}}^{\tilde{h}} = p_{\tilde{j}, \tilde{i}}^{\tilde{h}}$ (resp. $p_{\tilde{i}, \tilde{j}}^{\tilde{h}} \leq 1$) for all $\tilde{i}, \tilde{j}, \tilde{h} \in \tilde{\partial}(\Gamma)$. Note that a weakly distance-transitive digraph is weakly distance-regular.

Let G be a finite group and S a subset of G not containing the identity. The *Cayley digraph* $\Gamma = \text{Cay}(G, S)$ is a digraph with the vertex set G and the arc set $\{(x, sx) \mid x \in G, s \in S\}$.

In [8], Wang and Suzuki determined all commutative 2-valent weakly distance-regular digraphs. In [7], Suzuki determined all thin weakly distance-regular digraphs and proved the nonexistence of noncommutative weakly distance-regular digraphs of valency 2. Moreover, he proposed the project to classify weakly distance-regular digraphs of valency 3. In [9], Wang classified all commutative weakly distance-regular digraphs of valency 3 and girth 2. In this paper, we continue this project, and obtain the following result.

Theorem 1. *Let Γ be a weakly distance-regular digraph of valency 3 and girth more than 2. If Γ has two types of arcs, then Γ is isomorphic to one of the following digraphs:*

- (i) $\text{Cay}(\mathbb{Z}_4 \times \mathbb{Z}_g, \{(0, 1), (2, 1), (1, 0)\})$, where $g = 3$ or $g \geq 5$.
- (ii) $\Gamma_{q,2mq,1}$, $\Gamma_{q,mq+2,q}$ or $\Gamma_{q,2mq-2q+2t,q+1-t}$ in Construction 3, where $q \geq 3$, $m \geq 1$ and $2 \leq t \leq q - 1$.

This paper is organized as follows. In Section 2, we construct two families of weakly distance-regular digraphs of valency 3. In Section 3, we discuss some properties for circuits of weakly distance-regular digraphs. In Section 4, we prove our main theorem.

2 Constructions

In this section, we construct two families of weakly distance-regular digraphs of valency 3. For any element x in a residue class ring, we assume that \hat{x} denotes the minimum nonnegative integer in x . Denote $\beta(w) = (1 + (-1)^{w+1})/2$ for any integer w .

Proposition 2. *Let $g \geq 3$. Then $\Gamma_g := \text{Cay}(\mathbb{Z}_4 \times \mathbb{Z}_g, \{(1, 0), (0, 1), (2, 1)\})$ is a weakly distance-regular digraph if and only if $g \neq 4$.*

Proof. For any vertex (a, b) distinct with $(0, 0)$, we have

$$\tilde{\partial}((0, 0), (a, b)) = \begin{cases} (\hat{a}, 4 - \hat{a}), & \text{if } b = 0, \\ (\hat{b} + \beta(\hat{a}), g - \hat{b} + \beta(\hat{a})), & \text{if } b \neq 0. \end{cases}$$

Suppose $g \neq 4$. We will show that Γ_g is weakly distance-transitive. Let (a, b) and (x, y) be any two vertices satisfying $\tilde{\partial}((0, 0), (a, b)) = \tilde{\partial}((0, 0), (x, y))$. It suffices to verify that there exists an automorphism σ of Γ_g such that $\sigma(0, 0) = (0, 0)$ and $\sigma(a, b) = (x, y)$. If $(a, b) = (x, y)$, then the identity permutation is a desired automorphism. Now suppose $(a, b) \neq (x, y)$. Then $b \neq 0$, $y \neq 0$ and $(\hat{b} + \beta(\hat{a}), g - \hat{b} + \beta(\hat{a})) = (\hat{y} + \beta(\hat{x}), g - \hat{y} + \beta(\hat{x}))$. It follows that $b = y$ and $a - x = 2$. Let σ be the permutation on $V\Gamma_g$ such that

$$\sigma(u, v) = \begin{cases} (u, v), & \text{if } v \neq b, \\ (u + 2, v), & \text{if } v = b. \end{cases}$$

Routinely, σ is a desired automorphism.

In Γ_4 , $\tilde{\partial}((0, 0), (0, 2)) = \tilde{\partial}((0, 0), (2, 0)) = (2, 2)$. But $P_{(1,3),(3,3)}((0, 0), (0, 2)) = \{(1, 0)\}$ and $P_{(1,3),(3,3)}((0, 0), (2, 0)) = \emptyset$. Hence, Γ_4 is not a weakly distance-regular digraph. \square

Construction 3. Let q, s, k be integers with $q > 2$, $s > 2$ and $\max\{1, q - s + 2\} \leq k \leq q$. Write $s = 2mq + p$ with $m \geq 0$ and $0 \leq p < 2q$. Let $\Gamma_{q,s,k}$ be the digraph with the vertex set $\mathbb{Z}_q \times \mathbb{Z}_s$ whose arc set consists of $((a, b), (a+1, b))$, $((a, c), (a, c+1))$, $((a, d), (a+1, d-1))$, $((a, -1), (a - k + 1, 0))$ and $((a, 0), (a + k, -1))$, where $a \in \mathbb{Z}_q$, $b, c, d \in \mathbb{Z}_s$, $\hat{c} \neq s - 1$ and $d \neq 0$. See Figure 1.

In the following, we will prove that $\Gamma_{q,s,k}$ is a weakly distance-regular digraph if and only if one of the following holds:

- C1: $p = 0$ and $k = 1$.
- C2: $p = q + 2$ or $p = 2$, and $k = q$.
- C3: $4 \leq p \leq 2q - 2$, p is even and $k = q + 1 - p/2$.

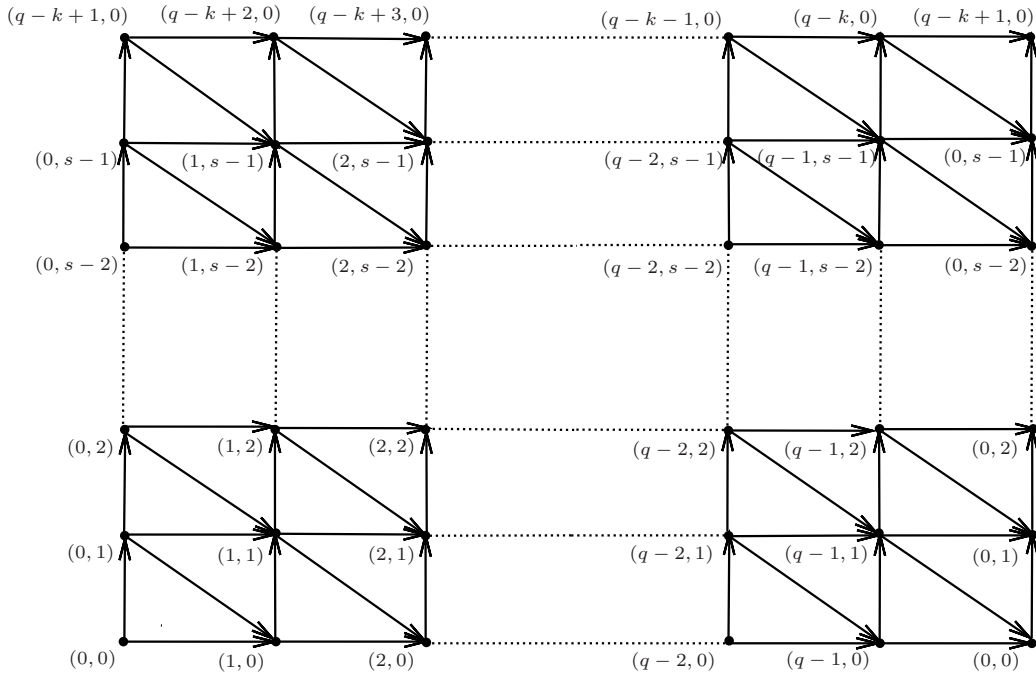


Figure 1: The digraph $\Gamma_{q,s,k}$.

Lemma 4. $\Gamma_{q,s,k}$ is a vertex transitive digraph.

Proof. Pick any vertex (a, b) . It suffices to show that there exists an automorphism σ of $\Gamma_{q,s,k}$ such that $\sigma(0, 0) = (a, b)$. Let σ be the permutation on $V\Gamma_{q,s,k}$ such that

$$\sigma(x, y) = \begin{cases} (x + a, y + b), & \text{if } y \in \{0, 1, 2, \dots, s - 1 - \hat{b}\}, \\ (x + a - k + 1, y + b), & \text{otherwise.} \end{cases}$$

Routinely, σ is a desired automorphism. \square

For any two integers i and j , we write $i \equiv j$ instead of $i \equiv j \pmod{q}$. For any vertex (a, b) of $\Gamma_{q,s,k}$, let $f(a, b)$, $g(a, b)$ and $h(a)$ be nonnegative integers less than q such that

$$f(a, b) \equiv \hat{a} + \hat{b} - k - p + 1, \quad g(a, b) \equiv q - \hat{a} - \hat{b} \text{ and } h(a) \equiv k - \hat{a} - 1. \quad (1)$$

By the structure of $\Gamma_{q,s,k}$, we have

$$\tilde{\partial}((0, 0), (a, b)) = (\min\{\hat{a} + \hat{b}, s - \hat{b} + f(a, b)\}, \min\{\hat{b} + g(a, b), s - \hat{b} + h(a)\}). \quad (2)$$

Lemma 5. Let C1, C2 or C3 hold. In $\Gamma_{q,s,k}$, $\partial((0, 0), (a, b)) = \hat{a} + \hat{b}$ if and only if $\partial((a, b), (0, 0)) = \hat{b} + g(a, b)$.

Proof. Let $M = s - 2\hat{b} - \hat{a} + f(a, b)$, $N = s - 2\hat{b} + h(a) - g(a, b)$ and $\hat{b} = n'q + r'$ with $0 \leq r' < q$. By (2), we only need to prove $M \geq 0$ if and only if $N \geq 0$. From (1), note that $f(a, b) + g(a, b)$ equals to $k - 1$ or $q + k - 1$, and $h(a)$ equals to $k - \hat{a} - 1$ or $q + k - \hat{a} - 1$.

Case 1. $f(a, b) + g(a, b) = k - 1$ and $h(a) = k - \hat{a} - 1$, or $f(a, b) + g(a, b) = q + k - 1$ and $h(a) = q + k - \hat{a} - 1$.

In this case, it is routine to check $M = N$, as desired.

Case 2. $f(a, b) + g(a, b) = k - 1$ and $h(a) = q + k - \hat{a} - 1$.

In this case, only C1 or C3 holds by $k < \hat{a} + 1$.

Assume that C1 holds. Then $f(a, b) = g(a, b) = 0$ and $h(a) = q - \hat{a}$, which imply that $\hat{a} + r' = 0$ or q . Since $h(a) < q$, one gets $\hat{a} \neq 0$. Hence, $\hat{a} + r' = q$. Then $M = 2(m - n')q - q - r' \geq 0$ if and only if $N = 2(m - n')q - r' \geq 0$.

Assume that C3 holds. Then $f(a, b) + g(a, b) = q - p/2$ and $h(a) = 2q - p/2 - \hat{a}$. By (1), note that $g(a, b) = q - \hat{a} - r'$ or $2q - \hat{a} - r'$, and $p/2 + \hat{a} > q$. Suppose $g(a, b) = q - \hat{a} - r'$. Then $f(a, b) = \hat{a} + r' - p/2$. From $\hat{a} + r' \leq q$, we have $0 < p/2 - r' < q$, which implies $M = 2(m - n')q + p/2 - r' \geq 0$ if and only if $N = 2(m - n')q + q + p/2 - r' \geq 0$. Suppose $g(a, b) = 2q - \hat{a} - r'$. Since $f(a, b) = \hat{a} - q + r' - p/2$, we have $0 \leq r' - p/2 < q$, which implies $M = 2(m - n')q - q + p/2 - r' \geq 0$ if and only if $N = 2(m - n')q + p/2 - r' \geq 0$.

Case 3. $f(a, b) + g(a, b) = q + k - 1$ and $h(a) = k - \hat{a} - 1$.

In this case, only C1 or C3 holds by $f(a, b) + g(a, b) \leq 2(q - 1)$.

Assume that C1 holds. Then $h(a) = \hat{a} = 0$ and $r' \neq 0$, which imply that $f(a, b) = r'$ and $g(a, b) = q - r'$. Then $M = 2(m - n')q - r' \geq 0$ if and only if $N = 2(m - n')q - q - r' \geq 0$.

Assume that C3 holds. Then $f(a, b) + g(a, b) = 2q - p/2$ and $h(a) = q - p/2 - \hat{a}$. By (1), note that $g(a, b) = q - \hat{a} - r'$ or $2q - \hat{a} - r'$, and $p/2 + \hat{a} \leq q$. Suppose $g(a, b) = q - \hat{a} - r'$. Then $f(a, b) = q + \hat{a} + r' - p/2$ and $0 \leq q + r' - p/2 < q$, which imply $M = 2(m - n' + 1)q - q - r' + p/2 \geq 0$ if and only if $N = 2(m - n')q + p/2 - r' \geq 0$. Suppose $g(a, b) = 2q - \hat{a} - r'$. Since $f(a, b) = \hat{a} + r' - p/2$, we have $p/2 - r' \leq \hat{a} < q$, which implies $M = 2(m - n')q + p/2 - r' \geq 0$ if and only if $N = 2(m - n')q - q + p/2 - r' \geq 0$.

Thus, desired result follows. \square

Lemma 6. *If C1, C2 or C3 holds, then $\Gamma_{q,s,k}$ is a weakly distance-regular digraph.*

Proof. We will prove that $\Gamma_{q,s,k}$ is weakly distance-transitive. Let (a, b) and (x, y) be two vertices satisfying $\tilde{\partial}((0, 0), (a, b)) = \tilde{\partial}((0, 0), (x, y))$. It suffices to find $\sigma \in \text{Aut}(\Gamma_{q,s,k})$ such that $\sigma(0, 0) = (0, 0)$ and $\sigma(a, b) = (x, y)$. By (2), we divide the proof into two cases.

Case 1. $\partial((0, 0), (a, b)) = \hat{a} + \hat{b}$.

Suppose $\partial((0, 0), (x, y)) = \hat{x} + \hat{y}$. Then $g(a, b) = g(x, y)$. By Lemma 5, we have $\hat{b} + g(a, b) = \hat{y} + g(x, y)$. This implies that $a = x$ and $b = y$. Hence, the identity permutation is a desired automorphism.

Suppose $\partial((0, 0), (x, y)) = s - \hat{y} + f(x, y)$. Then $\hat{a} + \hat{b} = s - \hat{y} + f(x, y)$, which implies $\hat{x} \equiv \hat{a} + \hat{b} + k - 1$ by $s \equiv p$. Hence, $\hat{x} = f(a, b)$ and $g(a, b) = h(x)$. From Lemma 5, we have $\hat{b} + g(a, b) = s - \hat{y} + h(x)$. This implies $\hat{y} = s - \hat{b}$. Let σ be the permutation on $V\Gamma_{q,s,k}$ such that

$$\sigma(u, v) = \begin{cases} (u, v), & \text{if } v = 0, \\ (f(u, v), -v), & \text{if } v \neq 0. \end{cases}$$

Routinely, σ is a desired automorphism.

Case 2. $\partial((0, 0), (a, b)) = s - \hat{b} + f(a, b)$.

Suppose $\partial((0, 0), (x, y)) = s - \hat{y} + f(x, y)$. Then $\hat{y} - \hat{b} = f(x, y) - f(a, b)$. We have $\hat{y} - \hat{b} \equiv \hat{x} + \hat{y} - \hat{a} - \hat{b}$. This implies $x = a$. By Lemma 5, one gets $s - \hat{b} + h(a) = s - \hat{y} + h(x)$, which implies that $y = b$. Hence, the identity permutation is a desired automorphism.

Suppose $\partial((0, 0), (x, y)) = \hat{x} + \hat{y}$. It is similar to Case 1 and the desired result holds. \square

By Lemma 4, for vertices (a, b) and (x, y) of $\Gamma_{q,s,k}$, we have

$$\tilde{\partial}((a, b), (x, y)) = \begin{cases} \tilde{\partial}((0, 0), (x - a, y - b)), & \text{if } \hat{y} \in \{\hat{b}, \hat{b} + 1, \dots, s - 1\}, \\ \tilde{\partial}((0, 0), (x - a + k - 1, y - b)), & \text{otherwise.} \end{cases}$$

Lemma 7. *If $\Gamma_{q,s,k}$ is weakly distance-regular, then C1, C2 or C3 holds.*

Proof. Suppose for the contrary that C1, C2 and C3 do not hold. Let $e = (0, 0)$, $z = (0, 1)$, $w = (k, s - 1)$ be the vertices of $\Gamma_{q,s,k}$, and $\alpha(v) = (3 + (-1)^v)/4$ for $v \in \mathbb{Z}$. By (2), note that $\tilde{\partial}(e, z) = \tilde{\partial}(e, w) = (1, q)$. To prove this lemma, we would pick proper $x, y \in V\Gamma_{q,s,k}$ such that $\tilde{\partial}(e, x) = \tilde{\partial}(e, y)$, and

$$P_{(1,q),\tilde{\partial}(z,x)}(e, y) = \emptyset \quad \text{or} \quad P_{(1,q),\tilde{\partial}(w,x)}(e, y) = \emptyset, \quad (3)$$

which contrary to $z \in P_{(1,q),\tilde{\partial}(z,x)}(e, x)$ or $w \in P_{(1,q),\tilde{\partial}(w,x)}(e, x)$.

Case 1. $k \neq q$ and $2k + p \geq 2(q - \alpha(p) + 2)$.

Let $x = (k, \alpha(s) - 1 + s/2)$ and $y = (k, \alpha(s) + s/2)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = \tilde{\partial}(z, y) = ((m + 1)q + 1, (m + 2)q - k)$ and $\tilde{\partial}(w, x) = ((m + 1)q, -\alpha(s) + s/2)$. Since $\partial(y, w) = -\alpha(s) + s/2 - 1$, we have (3) holds.

Case 2. $k \neq q$, $p \geq q$ and $2k + p \leq 2(q - \alpha(p) + 1)$.

Let $x = (0, \alpha(s) + s/2)$ and $y = (q + \alpha(p) + 1 - k - p/2, (m - 1)q + p + k - 1)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = (\alpha(s) + s/2, k - \alpha(s) - 1 + s/2)$ and $\tilde{\partial}(z, x) = (\alpha(s) - 1 + s/2, k - \alpha(s) + s/2)$. Since $\partial(z, y) = (\alpha(s) - 1 + s/2, k - \alpha(s) - 1 + s/2)$ and $\partial(w, y) = \alpha(s) + s/2$, we have (3) holds.

Case 3. $k \neq q$, $p < q$ and $2(\alpha(p) + 1) < 2k + p \leq 2(q - \alpha(p) + 1)$.

Let $x = (0, \alpha(s) + s/2)$ and $y = (q + \alpha(p) + 1 - k - p/2, (m - 1)q + p + k - 1)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = (\alpha(s) + s/2, k - \alpha(s) - 1 + s/2)$ and $\partial(w, y) = \alpha(s) + s/2$. If $p = 2(1 - \alpha(p))$, then $\tilde{\partial}(z, x) = (\alpha(s) - 1 + s/2, \alpha(s) - 1 + s/2)$ and $\tilde{\partial}(z, y) = (\alpha(s) - 1 + s/2, (m - 1)q + p + k - 2)$; if $p \neq 2(1 - \alpha(p))$, then $\tilde{\partial}(z, x) = (\alpha(s) - 1 + s/2, k - \alpha(s) + s/2)$ and $\tilde{\partial}(z, y) = (\alpha(s) - 1 + s/2, k - \alpha(s) - 1 + s/2)$. Hence, (3) holds.

Case 4. $k \neq q$, $p < q$ and $2k + p \leq 2(\alpha(p) + 1)$.

Let $x = (k, \alpha(s) - 1 + s/2)$ and $y = (k, \alpha(s) + s/2)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = \tilde{\partial}(z, y) = (mq + 1, q - \alpha(s) - 1 + s/2)$ and $\tilde{\partial}(w, x) = (mq, -\alpha(s) + s/2)$. Since $\partial(y, w) = -\alpha(s) - 1 + s/2$, we have (3) holds.

Case 5. $k = q$ and $p \geq q + 3$.

Let $x = (q - 1, (m - 1)q + p)$ and $y = (0, (m + 1)q)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = ((m + 1)q, (m + 1)q)$ and $\partial(z, x) = mq + p - 2$. Since $\partial(z, y) = (m + 1)q - 1$ and $\partial(w, y) = (m + 1)q$, we have (3) holds.

Case 6. $k = q$ and $3 \leq p \leq q + 1$.

Let $x = (q - 2, mq + 2)$ and $y = (q - 2, mq + 1)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = \tilde{\partial}(z, x) = ((m+1)q - 1, mq + 2)$. Since $\tilde{\partial}(z, y) = ((m+1)q - 2, mq + 2)$ and $\partial(y, w) = mq + 3$, we have (3) holds.

Case 7. $k = q$ and $p \leq 1$.

Let $x = (q - 1, mq + p)$ and $y = (0, mq)$. In this case, $\tilde{\partial}(e, x) = \tilde{\partial}(e, y) = (mq, mq)$ and $\partial(z, x) = (m + 1)q + p - 2$. Since $\partial(z, y) = mq - 1$ and $\partial(w, y) = mq$, we have (3) holds.

Therefore, the desired result holds. \square

Combining Lemmas 6 and 7, we obtain the following theorem.

Theorem 8. $\Gamma_{q,s,k}$ is weakly distance-regular, if and only if C1, C2 or C3 holds.

Finally, we shall show that every weakly distance-regular digraph $\Gamma_{q,s,k}$ is a Cayley digraph.

Proposition 9. Let $d = \frac{p}{2\gcd(q,p)}$, $l = \max\{w \mid 2^w \text{ divides } \gcd(q,p)\}$, $h = \frac{s}{2^l}$, $i = 2\{d\}$ and u be an integer such that $2^i q$ divides $(up - \gcd(q,p))$, where $\{d\}$ denotes the fractional part of d and $\gcd(q,p)$ denotes the greatest common divisor of q and p .

- (i) If C1 holds, then $\Gamma_{q,s,k}$ is isomorphic to $\text{Cay}(\mathbb{Z}_q \times \mathbb{Z}_{2mq}, \{(1, 0), (0, 1), (1, 2mq - 1)\})$, $m \geq 1$ and $q \geq 3$.
- (ii) If C2 holds, then $\Gamma_{q,s,k}$ is isomorphic to $\text{Cay}(\mathbb{Z}_{(mq+2)q}, \{1, mq + 2, mq + 1\})$, $m \geq 1$ and $q \geq 3$.
- (iii) If C3 holds, then $\Gamma_{q,s,k}$ is isomorphic to $\text{Cay}(\mathbb{Z}_{2^i q} \times \mathbb{Z}_{2^{-i}(2mq+p)}, \{(2^i u d, 1), (2^i - 2^i u d, i h - 1), (2^i, i h)\})$, where $q \geq 3$, $m \geq 0$, $4 \leq p \leq 2q - 2$ and p is even.

Proof. If C1 holds, then (i) is obvious. If C2 holds, then the mapping σ from $\Gamma_{q,s,k}$ to the digraph in (ii) satisfying $\sigma(a, b) = \hat{a}(mq + 2) + \hat{b}$ is an isomorphism.

Now suppose C3 holds. Let σ be the mapping from $\Gamma_{q,s,k}$ to the digraph in (iii) such that $\sigma(a, b) = (2^i \hat{a} + 2^i u d \hat{b}, i h \hat{a} + \hat{b})$. Note that σ is well defined.

We will show that σ is injective. It is clear for $i = 0$. If $i = 1$, by $2 \mid p$, then $l \geq 1$. Assume that $\sigma(x_1, y_1) = \sigma(x_2, y_2)$ for $(x_1, y_1), (x_2, y_2) \in V\Gamma_{q,s,k}$. Let $x = 2ud(\hat{y}_2 - \hat{y}_1) - 2(\hat{x}_1 - \hat{x}_2)$ and $y = (\hat{y}_2 - \hat{y}_1) - h(\hat{x}_1 - \hat{x}_2)$. Since $\sigma(x_1, y_1) = \sigma(x_2, y_2)$, we have $2q \mid x$ and $(mq + p/2) \mid y$, which imply $2^{l-1}h \mid y$ by $s = 2mq + p$. Hence, $h \mid (\hat{y}_2 - \hat{y}_1)$.

We claim $2^j \mid (\hat{y}_2 - \hat{y}_1)$ for $1 \leq j \leq l$. By $2q \mid (up - \gcd(q,p))$, we get $(2q/\gcd(q,p)) \mid (2ud - 1)$, which implies $2ud$ is odd. Since $2q \mid x$, one obtains $2 \mid (\hat{y}_2 - \hat{y}_1)$. Suppose $2^j \mid (\hat{y}_2 - \hat{y}_1)$ for some $j < l$. From $2^j \mid (mq + p/2)$, we have $2^j \mid y$ and $2^j \mid (\hat{x}_1 - \hat{x}_2)$, which imply $2^{j+1} \mid (\hat{y}_2 - \hat{y}_1)$ by $2^{j+1} \mid x$. So our claim is valid.

By $\gcd(2^l, h) = 1$, we obtain $(2mq + p) \mid (\hat{y}_2 - \hat{y}_1)$. Thus, $y_1 = y_2$ and $x_1 = x_2$. Therefore σ is a bijection. One can verify that $((x_1, y_1), (x_2, y_2))$ is an arc if and only if $(\sigma(x_1, y_1), \sigma(x_2, y_2))$ is an arc. Hence, σ is an isomorphism. \square

3 Circuits

In this section, we will discuss some properties for circuits of weakly distance-regular digraphs.

Let Γ be a digraph. Let $R = \{\Gamma_{\tilde{i}} \mid \tilde{i} \in \tilde{\partial}(\Gamma)\}$, where $\Gamma_{\tilde{i}} = \{(x, y) \in V\Gamma \times V\Gamma \mid \tilde{\partial}(x, y) = \tilde{i}\}$. If Γ is weakly distance-regular, then $(V\Gamma, R)$ is an association scheme. For more information about association schemes, see [3, 11]. For two nonempty subsets $E, F \subseteq R$, define

$$EF := \{\Gamma_{\tilde{h}} \mid \sum_{\Gamma_{\tilde{i}} \in E} \sum_{\Gamma_{\tilde{j}} \in F} p_{\tilde{i}, \tilde{j}}^{\tilde{h}} \neq 0\},$$

and write $\Gamma_{\tilde{i}}\Gamma_{\tilde{j}}$ instead of $\{\Gamma_{\tilde{i}}\}\{\Gamma_{\tilde{j}}\}$. For each nonempty subset F of R , define $\langle F \rangle$ to be the minimal equivalence relation containing F . Let

$$V\Gamma/F := \{F(x) \mid x \in V\Gamma\} \quad \text{and} \quad \Gamma_{\tilde{i}}^F := \{(F(x), F(y)) \mid y \in F\Gamma_{\tilde{i}}F(x)\},$$

where $F(x) := \{y \in V\Gamma \mid (x, y) \in \cup_{f \in F} f\}$. The digraph $(V\Gamma/F, \cup_{(1,s) \in \tilde{\partial}(\Gamma)} \Gamma_{1,s}^F)$ is said to be the *quotient digraph* of Γ over F , denoted by Γ/F . The size of $\Gamma_{\tilde{i}}^F(x) := \{y \in V\Gamma \mid \tilde{\partial}(x, y) = \tilde{i}\}$ depends only on \tilde{i} , denoted by $k_{\tilde{i}}$. For any $(a, b) \in \tilde{\partial}(\Gamma)$, we usually write $k_{a,b}$ (resp. $\Gamma_{a,b}$) instead of $k_{(a,b)}$ (resp. $\Gamma_{(a,b)}$).

Now we shall introduce some basic results which are used frequently in this paper.

Lemma 10. *Let Γ be a weakly distance-regular digraph. For each $\tilde{i} := (a, b) \in \tilde{\partial}(\Gamma)$, define $\tilde{i}^* = (b, a)$.*

- (i) $k_{\tilde{h}} p_{\tilde{i}, \tilde{j}}^{\tilde{h}} = k_{\tilde{i}} p_{\tilde{h}, \tilde{j}^*}^{\tilde{i}} = k_{\tilde{j}} p_{\tilde{i}^*, \tilde{h}}^{\tilde{j}}$.
- (ii) $k_{\tilde{i}} k_{\tilde{j}} = \sum_{\tilde{h} \in \tilde{\partial}(\Gamma)} k_{\tilde{h}} p_{\tilde{i}, \tilde{j}}^{\tilde{h}}$.
- (iii) $|\Gamma_{\tilde{i}} \Gamma_{\tilde{j}}| \leq \gcd(k_{\tilde{i}}, k_{\tilde{j}})$.

Proof. See Proposition 2.2 in [3, pp. 55-56] and Proposition 5.1 in [1]. □

In the remaining of this paper, we assume that Γ is a weakly distance-regular digraph of valency 3 satisfying $k_{1,q-1} = 1$ and $k_{1,g-1} = 2$, where $q, g \geq 3$ and $q \neq g$. Let $A_{i,j}$ denote a binary matrix with rows and columns indexed by $V\Gamma$ such that $(A_{i,j})_{x,y} = 1$ if and only if $\tilde{\partial}(x, y) = (i, j)$.

Lemma 11. *The following hold:*

$$A_{1,q-1}A_{1,g-1} = A_{1,g-1}A_{1,q-1}, \tag{4}$$

$$A_{1,g-1}A_{g-1,1} = A_{g-1,1}A_{1,g-1}. \tag{5}$$

Proof. By Lemma 10 (iii), we may assume that

$$A_{1,g-1}A_{1,q-1} = A_{i,j} \quad \text{and} \quad A_{1,q-1}A_{1,g-1} = A_{s,t}, \quad i, s \in \{1, 2\}.$$

We claim that $i = s = 2$. Suppose $i = 1$. Then $j = g - 1$ because of $k_{1,q-1} = 1$. By Lemma 10 (i), we get $p_{(g-1,1),(1,g-1)}^{(1,q-1)} = 2p_{(1,g-1),(1,q-1)}^{(1,g-1)} = 2$. By Lemma 10 (iii), $A_{g-1,1}A_{1,g-1} = 2I + 2A_{1,q-1}$, contrary to the fact that $A_{g-1,1}A_{1,g-1}$ is a symmetric matrix. Hence, $i = 2$. Similarly, $s = 2$ and our claim is valid.

Pick a path (x_0, x_1, x_2) with $\tilde{\partial}(x_0, x_1) = (1, g - 1)$ and $\tilde{\partial}(x_1, x_2) = (1, q - 1)$. Then $\partial(x_2, x_0) = j$. We may choose a path $(x_2, x_3, \dots, x_{j+1}, x_0)$. Since Γ has just two types of arcs, there exists an $i \in \{1, 2, \dots, j + 1\}$ such that $\tilde{\partial}(x_i, x_{i+1}) = (1, q - 1)$ and $\tilde{\partial}(x_{i+1}, x_{i+2}) = (1, g - 1)$, where $x_{j+2} = x_0$ and $x_{j+3} = x_1$. Since $\tilde{\partial}(x_i, x_{i+2}) = (2, t)$, one has $t \leq j$. Similarly, $j \leq t$. Hence, $j = t$ and (4) holds.

In view of Lemma 10 (iii), we have

$$A_{1,g-1}A_{g-1,1} = 2I + p_{(1,g-1),(g-1,1)}^{(s,s)} A_{s,s}, \quad s \geq 2, \quad (6)$$

$$A_{g-1,1}A_{1,g-1} = 2I + p_{(g-1,1),(1,g-1)}^{(t,t)} A_{t,t}, \quad t \geq 2. \quad (7)$$

By Lemma 10 (ii), we have $k_{s,s}p_{(1,g-1),(g-1,1)}^{(s,s)} = k_{t,t}p_{(g-1,1),(1,g-1)}^{(t,t)} = 2$, which implies that $p_{(1,g-1),(g-1,1)}^{(s,s)}, p_{(g-1,1),(1,g-1)}^{(t,t)} \in \{1, 2\}$. Let x_0 and x_s be two vertices satisfying $\tilde{\partial}(x_0, x_s) = (s, s)$. Suppose $p_{(1,g-1),(g-1,1)}^{(s,s)} = 2$. Pick two distinct vertices $x, y \in P_{(1,g-1),(g-1,1)}(x_0, x_s)$. By $x_0 \in P_{(g-1,1),(1,g-1)}(x, y)$ and (7), $\tilde{\partial}(x, y) = (t, t)$. It follows that $p_{(g-1,1),(1,g-1)}^{(t,t)} = 2$. Similarly, if $p_{(g-1,1),(1,g-1)}^{(t,t)} = 2$, then $p_{(1,g-1),(g-1,1)}^{(s,s)} = 2$ by (6). Hence, $p_{(1,g-1),(g-1,1)}^{(s,s)} = p_{(g-1,1),(1,g-1)}^{(t,t)}$. In order to show (5), we shall prove $s = t$. Pick $x \in P_{(1,g-1),(g-1,1)}(x_0, x_s)$ and a path $P := (x_0, x_1, \dots, x_s)$.

Case 1. P contains an arc of type $(1, g - 1)$.

By (4), without loss of generality, we may assume that $\tilde{\partial}(x_0, x_1) = (1, g - 1)$. Pick $y \in \Gamma_{1,g-1}(x_s) \setminus \{x\}$. In view of (7), if $x \neq x_1$, from $x_0 \in P_{(g-1,1),(1,g-1)}(x_1, x)$, then $\partial(x_1, x) = t \leq s$; if $x = x_1$, from $x_0 \in P_{(g-1,1),(1,g-1)}(x, y)$, then $\partial(x, y) = t \leq s$.

Case 2. P only contains arcs of type $(1, q - 1)$.

In this case, $A_{1,q-1}^s \neq I$. By (4), there exists a path $(x_0, y_1, y_2, \dots, y_s, x)$ containing the unique arc (x_0, y_1) of type $(1, g - 1)$. If $x = y_1$, by Lemma 10 (iii), we have $A_{1,q-1}^s = I$, a contradiction. Therefore, $x \neq y_1$. By $x_0 \in P_{(g-1,1),(1,g-1)}(y_1, x)$ and (7), one has $\partial(y_1, x) = t \leq s$.

Similarly, $t \geq s$, which implies $s = t$, as desired. \square

In the following, let $F = \langle \Gamma_{1,g-1} \rangle$ and fix $x \in V\Gamma$. Then Γ/F is isomorphic to a circuit C_m of length m . Let Δ be a digraph with the vertex set $F(x)$ such that (y, z) is an arc of Δ if (y, z) is an arc of type $(1, g - 1)$ in Γ .

Lemma 12. *Suppose that every circuit of length g contains arcs of the same type in Γ . Then $\Delta_{t,g-t}(x_0) = \Gamma_{t,g-t}(x_0)$ for each $x_0 \in F(x)$ and $t \in \{1, 2, \dots, g - 1\}$.*

Proof. Note that every arc of type $(1, g - 1)$ is contained in a circuit of length g with all arcs of type $(1, g - 1)$. It follows that, for any such circuit $(x_0, x_1, \dots, x_{g-1})$, we have

$\tilde{\partial}_\Gamma(x_0, x_i) = (i, g - i)$, where $1 \leq i \leq g - 1$. Then every arc of Δ is contained in a circuit of length g in Δ .

For any $x_t \in \Gamma_{t, g-t}(x_0)$, there exists a circuit $C_g := (x_0, x_1, \dots, x_t, \dots, x_{g-1})$ in Γ . Hence, C_g only contains the arcs of same type. Suppose that each arc of C_g is of type $(1, q - 1)$. Then, $q < g$ and every circuit of length q in Γ only contains arcs of type $(1, q - 1)$. It follows that $A_{1, q-1}^q = I$. Since $x_0 \neq x_l$ for $1 \leq l \leq g - 1$, $k_{1, q-1} = 1$ implies that g is the minimum positive integer such that $A_{1, q-1}^g = I$, a contradiction. Consequently, each arc of C_g is of type $(1, g - 1)$. Therefore, $(x_0, x_t) \in \Delta_{t, g-t}$; and so $\Gamma_{t, g-t}(x_0) \subseteq \Delta_{t, g-t}(x_0)$. Conversely, pick any $x_t \in \Delta_{t, g-t}(x_0)$. Then, in Γ , there exists a circuit $(x_0, x_1, \dots, x_t, \dots, x_{g-1})$ each of whose arcs is of type $(1, g - 1)$. Hence, $(x_0, x_t) \in \Gamma_{t, g-t}$; and so $\Delta_{t, g-t}(x_0) \subseteq \Gamma_{t, g-t}(x_0)$. Thus, the desired result holds. \square

Lemma 13. *If $F(x) = V\Gamma$, then there exists a circuit of length g containing different types of arcs.*

Proof. Suppose for the contrary that every circuit of length g contains the same type of arcs. By the Lemma 12, $\Gamma_{t, g-t} = \Delta_{t, g-t}$ for any $1 \leq t \leq g - 1$. By (5), the proof of Proposition 4.3 in [8] implies that Δ is isomorphic to $\Gamma_1 := \text{Cay}(\mathbb{Z}_{2g}, \{1, g + 1\})$ or $\Gamma_2 := \text{Cay}(\mathbb{Z}_g \times \mathbb{Z}_g, \{(0, 1), (1, 0)\})$.

Case 1. $\Delta \simeq \Gamma_1$.

Choose $y \in \mathbb{Z}_{2g} \setminus \{0, 1, g + 1\}$ and $t \in \mathbb{Z}_{2g}$ such that $\tilde{\partial}_\Gamma(0, y) = (1, q - 1)$, $\hat{t} \equiv \hat{y} \pmod{g}$ and $\hat{t} \in \{0, 2, 3, \dots, g - 1\}$. Since $(y + 1, y + 2, \dots, y - t + g - 1, 0, y)$ is a path of length $g - \hat{t}$, $\partial_\Gamma(y + 1, y) = g - 1 \leq g - \hat{t}$. It follows that $t = 0$, and so $\hat{y} = g$. Therefore, $\tilde{\partial}_\Gamma(0, g) = (1, q - 1)$. Similarly, $\tilde{\partial}_\Gamma(g, 0) = (1, q - 1)$. Hence, $q = 2$, a contradiction.

Case 2. $\Delta \simeq \Gamma_2$.

Pick $(i, j) \in \Gamma_{1, q-1}(0, 0)$. Since $\tilde{\partial}_\Delta((0, 0), (0, j)) = (\hat{j}, g - \hat{j})$, by Lemma 12, we have $\tilde{\partial}_\Gamma((0, 0), (0, j)) = (\hat{j}, g - \hat{j})$. It follows that $i \neq 0$. By Lemma 10 (i), one gets $p_{(\hat{i}, g-\hat{i}), (\hat{j}, g-\hat{j})}^{(1, q-1)} = k_{\hat{i}, g-\hat{i}} p_{(1, q-1), (g-\hat{j}, \hat{j})}^{(\hat{i}, g-\hat{i})}$. Since $(i, j) \in P_{(1, q-1), (g-\hat{j}, \hat{j})}((0, 0), (i, 0))$ in Γ , $p_{(1, q-1), (g-\hat{j}, \hat{j})}^{(\hat{i}, g-\hat{i})} = 1$, which implies that $p_{(\hat{i}, g-\hat{i}), (\hat{j}, g-\hat{j})}^{(1, q-1)} = k_{\hat{i}, g-\hat{i}}$.

Let $((a, b), (a', b'))$ be an arc of type $(1, q - 1)$. Then $P_{(\hat{i}, g-\hat{i}), (\hat{j}, g-\hat{j})}((a, b), (a', b')) = \Gamma_{\hat{i}, g-\hat{i}}(a, b)$. Since $(a + i, b), (a, b + i) \in \Delta_{\hat{i}, g-\hat{i}}(a, b)$, by Lemma 12, $(a', b') \in \Gamma_{\hat{j}, g-\hat{j}}(a + i, b) \cap \Gamma_{\hat{j}, g-\hat{j}}(a, b + i)$. By Lemma 12 again,

$$(a', b') \in \{(a + i + j, b), (a + i, b + j)\} \cap \{(a + j, b + i), (a, b + i + j)\}.$$

Since $i \neq 0$, we have $(a', b') = (a + i, b + j) = (a + j, b + i)$, which implies that $i = j$. Thus, $\Gamma \simeq \text{Cay}(\mathbb{Z}_g \times \mathbb{Z}_g, \{(1, 0), (0, 1), (i, i)\})$. Since $g \neq q$, one gets $i \neq 1$. Let $g = n\hat{i} + r$ with $0 \leq r \leq \hat{i} - 1$. If $r \neq 0$, then $\tilde{\partial}_\Gamma((0, 0), (1, 1)) = \tilde{\partial}_\Gamma((0, 0), (i, i + 1)) = (2, n + 2r - 2)$; if $r = 0$, then $\tilde{\partial}_\Gamma((0, 0), (1, 1)) = \tilde{\partial}_\Gamma((0, 0), (i, i + 1)) = (2, n + 2\hat{i} - 3)$. But we have $(1, 0) \in P_{(1, g-1), (1, g-1)}((0, 0), (1, 1))$ and $P_{(1, g-1), (1, g-1)}((0, 0), (i, i + 1)) = \emptyset$ in Γ , a contradiction. \square

Lemma 14. *Every circuit of length q in Γ only contains the arcs of the same type. In particular,*

$$A_{1, q-1}^2 = A_{2, q-2}. \quad (8)$$

Proof. If $F(x) = V\Gamma$, then $q < g$ by Lemma 13 and the desired result follows. Suppose $F(x) \neq V\Gamma$. Assume the contrary, namely, there exists a circuit $(x_0, x_1, \dots, x_{q-1})$ containing arcs of different types. Since $\Gamma/F \simeq C_m$ with $m \geq 2$, there exist at least two arcs of type $(1, q-1)$ in this circuit. By (4), we may assume that $\tilde{\partial}(x_0, x_1) = \tilde{\partial}(x_1, x_2) = (1, q-1)$ and $\tilde{\partial}(x_{q-1}, x_0) = (1, g-1)$. By the claim in Lemma 11, $\tilde{\partial}(x_{q-1}, x_1) = (2, q-2)$. Since $k_{1,q-1} = 1$, by Lemma 10 (ii), one has $k_{\tilde{\partial}(x_0, x_2)} = 1$. Therefore, $\tilde{\partial}(x_0, x_2) = (2, q-2)$. But $P_{(1,q-1),(1,q-1)}(x_0, x_2) = \{x_1\}$ and $P_{(1,q-1),(1,q-1)}(x_{q-1}, x_1) = \emptyset$, a contradiction. Lemma 10 (iii) implies (8). \square

Lemma 15. *For any circuit $(x_0, x_1, \dots, x_{l-1})$ with $\tilde{\partial}(x_{l-1}, x_0) = (1, g-1)$, there exists $i \in \{0, 1, \dots, l-2\}$ such that $\tilde{\partial}(x_i, x_{i+1}) = (1, g-1)$.*

Proof. Suppose for the contradiction that $\tilde{\partial}(x_i, x_{i+1}) = (1, q-1)$ for any $i = 0, 1, \dots, l-2$. By Lemma 10 (iii), we have $A_{g-1,1} = A_{1,q-1}^{l-1}$. Then $A_{g-1,1}$ is a permutation matrix, a contradiction. \square

Lemma 16. *$F(x) \neq V\Gamma$ if and only if every circuit of length g in Γ only contains the arcs of the same type.*

Proof. Suppose $F(x) \neq V\Gamma$. Assume the contrary, namely, $(x_0, x_1, \dots, x_{g-1})$ is a circuit containing arcs of different types such that $\tilde{\partial}(x_0, x_1) = (1, g-1)$. By (4) and Lemma 15, we may assume that $\tilde{\partial}(x_1, x_2) = (1, q-1)$ and $\tilde{\partial}(x_{g-1}, x_0) = (1, g-1)$. By the claim in Lemma 11, $\tilde{\partial}(x_0, x_2) = (2, g-2)$. If $\partial(x_{g-1}, x_1) = 1$, from $F(x) \neq V\Gamma$, then $\tilde{\partial}(x_{g-1}, x_1) = (1, g-1)$, which implies $(x_1, x_2, \dots, x_{g-1})$ is a circuit of length $g-1$ containing an arc of type $(1, g-1)$, a contradiction. Hence, $\tilde{\partial}(x_{g-1}, x_1) = (2, g-2)$. The fact that $x_2 \notin F(x_0)$ implies that $P_{(1,g-1),(1,g-1)}(x_0, x_2) = \emptyset$, contradicting to $x_0 \in P_{(1,g-1),(1,g-1)}(x_{g-1}, x_1)$.

The converse is true by Lemma 13. \square

4 The proof of Theorem 1

In this section, we assume that $F = \langle \Gamma_{1,g-1} \rangle$ and x is a fixed vertex of Γ .

Lemma 17. *If $F(x) \neq V\Gamma$, then $\Gamma/F \simeq C_2$.*

Proof. Suppose for the contradiction that $\Gamma/F \simeq C_m$ with $m \geq 3$. Choose a path (x_0, x_1, x_2, x_3) such that $\tilde{\partial}(x_0, x_1) = \tilde{\partial}(x_1, x_2) = (1, q-1)$ and $\tilde{\partial}(x_2, x_3) = (1, g-1)$. Since $\partial(F(x_0), F(x_2)) = 2$, $k_{1,q-1} = 1$ implies that $\tilde{\partial}(x_0, x_3) = (3, l)$ for some l . Then there exists a shortest path $(x_3, x_4, y_2, \dots, x_{l+2}, x_0)$. By Lemma 15 and (4), we may assume that $\tilde{\partial}(x_3, x_4) = (1, g-1)$. Since $\partial(F(x_1), F(x_4)) = 1$ and $k_{1,q-1} = 1$, we have $\partial(x_1, x_4) = 2$ or 3. If $\partial(x_1, x_4) = 2$, by $F(x) \neq V\Gamma$, then $\tilde{\partial}(x_2, x_4) = (1, g-1)$, which implies $g = 2$ by $x_2 \in P_{(g-1,1),(1,g-1)}(x_3, x_4)$ and (7), a contradiction. Hence, $\tilde{\partial}(x_1, x_4) = (3, t)$ for some $t \leq l$. From $m \geq 3$ and (4), there exists a path $(x_4, y_1, y_2, \dots, y_{t-2}, x_0, x_1)$. Then $(x_3, x_4, y_1, y_2, \dots, y_{t-2}, x_0)$ is a path of length t ; and so $l \leq t$. Hence, $l = t$. By (8), $x_2 \in P_{(2,q-2),(1,g-1)}(x_0, x_3)$. Then there exists $y \in P_{(2,q-2),(1,g-1)}(x_1, x_4)$. From $k_{1,q-1} = 1$, $\tilde{\partial}(x_2, y) = (1, q-1)$, which implies $\Gamma_{1,q-1} \in F$, a contradiction. \square

Proposition 18. *If $F(x) \neq V\Gamma$, then Γ is isomorphic to one of the digraphs in Theorem 1 (i).*

Proof. By Lemma 17, $V\Gamma$ has a partition $F(x) \dot{\cup} F(x')$. Let Δ and Δ' be the subdigraphs of Γ induced on $F(x)$ and $F(x')$, respectively. By (4) and $k_{1,q-1} = 1$, $\sigma : F(x) \rightarrow F(x')$, $y \mapsto y'$ is an isomorphism mapping from Δ to Δ' , where $y' \in \Gamma_{1,q-1}(y)$. By Lemmas 12 and 16, $\Gamma_{r,g-r}(y) = \Delta_{r,g-r}(y)$ for each $y \in F(x)$ and $r \in \{1, 2, \dots, g-1\}$. By (5), the proof of Proposition 4.3 in [8] implies that Δ is isomorphic to $\Gamma_1 := \text{Cay}(\mathbb{Z}_g \times \mathbb{Z}_g, \{(1, 0), (0, 1)\})$ or $\Gamma_2 := \text{Cay}(\mathbb{Z}_{2g}, \{1, g+1\})$. Suppose that τ_i is an isomorphism from Γ_i to Δ if Γ_i is isomorphic to Δ .

We claim that $\Delta \simeq \Gamma_2$. Suppose for the contrary that $\Delta \simeq \Gamma_1$. Write $\tau_1(a, b) = (a, b, 0)$ and $\sigma(a, b, 0) = (a, b, 1)$ for each $(a, b) \in \mathbb{Z}_g \times \mathbb{Z}_g$. Let $((0, 0, 1), (c, d, 0))$ be an arc of type $(1, q-1)$. By (8), $\tilde{\partial}_\Gamma((0, 0, 0), (c, d, 0)) = (2, q-2)$. Lemma 12 implies that $c \neq 0$ and $d \neq 0$. By Lemma 12 again, we have $(c, d, 0) \in P_{(2,q-2),(g-\hat{d},\hat{d})}((0, 0, 0), (c, 0, 0))$ and $\tilde{\partial}_\Gamma((0, 0, 0), (c, 0, 0)) = \tilde{\partial}_\Gamma((0, 0, 0), (0, c, 0))$. By $k_{2,q-2} = 1$, we have $(0, c, 0) \in \Gamma_{g-\hat{d},\hat{d}}(c, d, 0)$. Then $(0, c, 0) \in \{(c, 0, 0), (c-d, d, 0)\}$ by Lemma 12. Hence, $c = d$.

Suppose $\hat{c} = g-1$. Since $((0, 0, 1), (-1, -1, 0), (0, -1, 0), (0, 0, 0))$ is a shortest path, $q = 4$, contrary to Lemma 14. Suppose $\hat{c} \neq g-1$. Then $\tilde{\partial}_\Gamma((0, 0, 0), (c, c+1, 0)) = (3, l)$ for some l . Pick a path $((c, c+1, 0), x_1, x_2, \dots, x_{l-1}, (0, 0, 0))$. By Lemma 15 and (4), we may assume that $\tilde{\partial}_\Gamma((c, c+1, 0), x_1) = (1, g-1)$. By (7), we have $\tilde{\partial}_\Gamma((0, 0, 1), x_1) = (3, t)$ for some $t \leq l$. Since $F(x) \neq V\Gamma$, $k_{1,q-1} = 1$ implies that there exists a path $(x_1, y_1, y_2, \dots, y_{t-2}, (0, 0, 0), (0, 0, 1))$. Then $((c, c+1, 0), x_1, y_1, y_2, \dots, y_{t-2}, (0, 0, 0))$ is a path of length t ; and so $l \leq t$. Hence $l = t$. By (8) and $x_1 \in V\Delta$, one has $(c, c, 0) \in P_{(2,q-2),(1,g-1)}((0, 0, 0), (c, c+1, 0))$ and $P_{(2,q-2),(1,g-1)}((0, 0, 1), x_1) = \emptyset$ in Γ , a contradiction. Therefore, our claim is valid.

Write $\tau_2(a) = (a, 0)$ and $\sigma(a, 0) = (a, 1)$ for each $a \in \mathbb{Z}_{2g}$. Let $((a, 1), (a+k_a, 0))$ be an arc of type $(1, q-1)$. Then $k_a \neq 0$. By (8), $\tilde{\partial}_\Gamma((a, 0), (a+k_a, 0)) = (2, q-2)$. By Lemma 12, $\tilde{\partial}_\Delta((a, 0), (a+k_a, 0)) \neq (t, g-t)$ for any $t \in \{1, 2, \dots, g-1\}$. Since $\bigcup_{1 \leq t \leq g-1} \Delta_{t,g-t}(a, 0) = V\Delta \setminus \{(a, 0), (a+g, 0)\}$, one has $\hat{k}_a = g$. Then, $\Gamma \simeq \text{Cay}(\mathbb{Z}_4 \times \mathbb{Z}_g, \{(0, 1), (1, 0), (2, 1)\})$ and the result holds by Proposition 2. \square

Lemma 19. *If $F(x) = V\Gamma$, then $p_{(1,g-1),(1,g-1)}^{(1,q-1)} = 2$.*

Proof. By Lemma 13, there exists a circuit of length g with different types of arcs. Let $C := (x_0, x_1, \dots, x_{g-1})$ be such a circuit with the minimum number of arcs of type $(1, g-1)$. Suppose C contains t arcs of types $(1, g-1)$. Lemma 15 implies that $t \geq 2$. By (4), we may assume that $\tilde{\partial}(x_i, x_{i+1}) = (1, g-1)$ for $0 \leq i \leq t$. We claim that $\tilde{\partial}(x_0, x_2) = (1, q-1)$. Suppose not. By the claim in Lemma 11 and (7), we have $\tilde{\partial}(x_{g-1}, x_1) = \tilde{\partial}(x_0, x_2) = (2, g-2)$. Since $x_0 \in P_{(1,q-1),(1,g-1)}(x_{g-1}, x_1)$, there exists $x'_1 \in P_{(1,q-1),(1,g-1)}(x_0, x_2)$. The circuit $C' := (x_0, x'_1, x_2, \dots, x_{g-1})$ contains just $t-1$ arcs of type $(1, g-1)$, a contradiction. Thus, our claim is valid. It follows that $p_{(1,q-1),(g-1,1)}^{(1,g-1)} = 1$. By Lemma 10 (i), the desired result holds. \square

Let $H = \langle \Gamma_{1,q-1} \rangle$ and $H(x_{0,0}), H(x_{0,1}), \dots, H(x_{0,s-1})$ be all pairwise distinct vertices of Γ/H . Since $q < g$, the subdigraph induced on each $H(x_{0,j})$ is a circuit of length q with arcs of type $(1, q-1)$, say $(x_{0,j}, x_{1,j}, \dots, x_{q-1,j})$. It follows that $s \geq 2$.

Proposition 20. *If $F(x) = V\Gamma$, then Γ is isomorphic to one of the digraphs in Theorem 1 (ii).*

Proof. Suppose $\partial(H(x_{0,0}), H(x_{0,1})) = 1$. By (4), we may assume that $\tilde{\partial}(x_{0,0}, x_{0,1}) = (1, g-1)$. By Lemma 19, one has $\tilde{\partial}(x_{0,1}, x_{1,0}) = (1, g-1)$, which implies $\partial(H(x_{0,1}), H(x_{0,0})) = 1$. Since $F(x) = V\Gamma$, Γ/H is a connected undirected graph. By $k_{1,g-1} = 2$, Γ/H is an undirected circuit of length s . Suppose $s = 2$. Pick $y \in \Gamma_{1,g-1}(x_{0,1}) \setminus \{x_{1,0}\}$. Then $y = x_{i,0}$ for some $i \geq 2$, and $(x_{0,1}, y, x_{i+1,0}, \dots, x_{q-1,0}, x_{0,0})$ is a path of length $q-i+1$ from $x_{0,1}$ to $x_{0,0}$, contrary to the fact $\partial(x_{0,1}, x_{0,0}) = g-1$. Hence, $s \geq 3$.

Let $(H(x_{0,0}), H(x_{0,1}), \dots, H(x_{0,s-1}))$ be an undirected circuit. By (4), we may assume that $(x_{0,0}, x_{0,1}, \dots, x_{0,s-1})$ is a path with arcs of type $(1, g-1)$. By Lemma 19, $(x_{0,j}, x_{0,j+1}, x_{1,j}, x_{1,j+1}, x_{2,j}, \dots, x_{q-1,j}, x_{q-1,j+1})$ is a circuit with arcs of type $(1, g-1)$ for $j = 0, 1, \dots, s-2$. Hence, there exists $k \in \{1, 2, \dots, q\}$ such that $\tilde{\partial}(x_{0,s-1}, x_{q-k+1,0}) = (1, g-1)$, where the first subscription of x are taken modulo q . By Lemma 19 again, we obtain $\tilde{\partial}(x_{i,s-1}, x_{i-k+1,0}) = \tilde{\partial}(x_{i-k+1,0}, x_{i+1,s-1}) = (1, g-1)$ for each i . Since

$$(x_{0,0}, x_{0,1}, \dots, x_{0,s-1}, x_{q-k+1,0}, x_{q-k+2,0}, \dots, x_{q-1,0})$$

is a circuit of length $s+k-1$ with different types of arcs, by Lemma 14, we get $s+k-1 > q$. From Theorem 8, the desired result follows. \square

Combining Propositions 18 and 20, we complete the proof of Theorem 1.

In the forthcoming paper [10], we shall classify cubic commutative weakly distance-regular digraphs with one type of arcs.

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